### NORTHWESTERN UNIVERSITY

# Methods for Derivative-Free Optimization with Applications in Machine Learning

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Melody Qiming Xuan

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### ABSTRACT

Methods for Derivative-Free Optimization with Applications in Machine Learning

#### Melody Qiming Xuan

Derivative-free optimization (DFO) has received growing attention due to important problems arising in practice. Various research communities, ranging from machine learning to engineering design, have adopted distinct DFO methods. In this thesis, we present extensive studies as a meaningful step towards a comprehensive understanding of DFO methods. We study the relative merits of this disparate class of methods in a controlled setting, where the derivatives of the functions are unavailable and the functions values are subject to noise.

In Chapter 2, we provide numerical investigations of finite-difference-based approaches against interpolation-based methods for solving unconstrained, (nonlinear) least-square and constrained optimization problems. This approach relies on finite differences to perform gradient approximations and applies existing gradient-based nonlinear optimization solvers. The numerical results demonstrate the effectiveness of finite-difference-based methods for both deterministic and noisy problems, when a proper implementation of finite differences is employed. The empirical results suggest that finite-difference-based methods are competitive with IBO methods, which are among the most reliable and efficient DFO solvers.

We present in Chapter 3 an adaptive procedure to choosing a reliable finite differencing interval for noisy DFO problems, using a bisection approach. When noise is present in the objective function, the finite differencing interval needs to be chosen meticulously to balance the truncation error and noise. In particular, the optimal interval depends on the noise level of the function and a bound on the higher-order derivative. We show that the proposed adaptive procedure produces near-optimal differencing intervals without explicitly approximating higher-order derivatives. We demonstrate its robustness and accuracy on selected problems.

In Chapter 4, we investigate general constrained DFO, where the constraints are analytically available. We propose to extend unconstrained interpolation-based optimization methods by combining them with a general-purpose nonlinear programming solver. This is a feasible method that constructs quadratic models of the objective and generates feasible iterates at every iteration. We discuss its effectiveness and limitations in this chapter. Our numerical results suggest that, when the objective is expensive to evaluate and the cost for evaluating constraints is negligible, this seemingly expensive approach is competitive with a state-of-the-art constrained DFO solvers that employs finite differences.

Next, in Chapter 5 we consider a DFO problem that arises in Natural Language Processing. With the wide success of pre-trained large language models (LLMs) and their growing model size, it is imperative to consider strategies that efficiently adapt pre-trained general LLMs to individual language tasks. In particular, we consider prompt design, which makes use of different prompts to condition LLMs for different tasks. We propose to treat prompt design as a noisy DFO problem, wherein we search for a prompt that optimizes performance without accessing the parameters of the underlying LLM. We present competitive results for solving prompt design using DFO methods and identify the relative merits of NEWUOA, FD L-BFGS and CMA-ES.

The observations in Chapter 5 motivate us to further analyze the relative strengths of NEWUOA and CMA-ES in Chapter 6. In particular, we consider a broader range of unconstrained problems from the CUTEst problem set. We observe that, contrary to the noiseless problems, while NEWUOA still demonstrates greater efficiency in achieving a desired level of accuracy, it is adversely affected by noise in function evaluations and consequently achieves lower accuracy compared to CMA-ES. Therefore, we investigate the impact of noise on IBO methods and briefly discuss strategies to improve the final accuracy of IBO methods in the presence of noise.

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#### CHAPTER 1

#### Introduction

The problem of minimizing a smooth function when its derivatives are not available has been widely studied in the literature [29, 59, 69]. A variety of methods have been proposed to solve general unconstrained derivative-free optimization (DFO) problems, with some being extended to nonlinear least squares and general constrained problems. Various research communities have indeed adopted distinct DFO methods to solve important problems arising from machine learning to engineering design. The objective function in those applications is often expensive to evaluate and may be perturbed with noise, e.g., computational noise, while DFO methods are typically designed for noiseless problems. In this thesis, we study several disparate classes of DFO methods in a controlled setting for different classes of problems.

Let us consider a problem of the form

(1.0.1) 
$$\min_{x \in \Omega} \phi(x)$$

where  $\phi : \mathbb{R}^n \to \mathbb{R}$  is a smooth function and  $\Omega \subseteq \mathbb{R}^n$  denotes the feasible set. The feasible set  $\Omega$  is specified by a set of equality and inequality constraints:

(1.0.2) 
$$\Omega = \{ x \in \mathbb{R}^n | \psi_i(x) = 0, \forall i \in \mathcal{E}; \psi_j(x) \le 0, \forall j \in \mathcal{I} \},$$

where  $\psi : \mathbb{R}^n \to \mathbb{R}^m$  is smooth, and  $\mathcal{E}$  and  $\mathcal{I}$  are finite index sets that can be empty (the problem reduces to unconstrained optimization if  $\mathcal{E} = \mathcal{I} = \emptyset$ ). We assume that the objective values  $\phi(x)$  are only accessible through a (noisy) zero-order oracle such that the derivatives of  $\phi$  are not available. Following the taxonomy of constraints [35], we assume that the constraints can be analytically available with access to their derivatives (class A) or only accessible via a (noisy) zero-order oracle (class S).

Formally, we define the noisy zero-order oracles of the objective and constraints as:

(1.0.3) 
$$f(x) = \phi(x) + \epsilon(x)$$

(1.0.4) 
$$c_i(x) = \psi_i(x) + \epsilon_i(x)$$

where  $\epsilon(\cdot)$  and  $\epsilon_i(\cdot)$  model the noise, which can be either deterministic or stochastic. In practice, such evaluation errors can stem from finite precision arithmetics as well as stochastic optimization problems encountered in machine learning applications. When the noise is uniformly bounded, we refer to  $\epsilon_f$  and  $\epsilon_c$  as the noise levels of the functions respectively, where

(1.0.5) 
$$|\epsilon(x)| \le \epsilon_f, \quad |\epsilon_i| \le \epsilon_c.$$

In the case when the noise terms are stochastic, we denote the standard deviations  $\sigma_f$ ,  $\sigma_c$  of the noise terms as the noise levels of the objective and constraints respectively. The noise level can be estimated via sampling (stochastic) and/or the difference table (deterministic) [70]. Thus we assume throughout this thesis that the noise levels of the functions are known. Several approaches have been developed for solving problem (1.0.1), including gradient approximation based methods, interpolation-based optimization (IBO) methods and randomized algorithms. In Chapter 2, we consider unconstrained, nonlinear least-squares and constrained optimization problems that can be only accessed via a noisy zerothorder oracle with stochastic noise. We provide a sophisticated adaptive finite-difference differencing interval selection scheme in Chapter 3 for bounded noise. In Chapter 4, we study (1.0.1) where both  $\mathcal{E}$  and  $\mathcal{I}$  can be nonempty and the constraints are available in analytical forms. In Chapter 5, we present an application of noisy unconstrained DFO problem arising in large language models. In Chapter 6 we consider unconstrained DFO problems again with noisy evaluations to study the behavior of DFO methods and discuss potential strategies for improving the final accuracy of interpolation-based optimization methods.

In the remainder of this chapter, we provide an overview of several classes of methods that can be applied to solve (1.0.1) and discuss their advantages and disadvantages, respectively. For methods that are not covered in this thesis, such as direct search and deterministic global optimization methods, see [59].

#### 1.1. Gradient Approximations

One simple yet effective approach to solving (1.0.1) is to perform approximations to the derivatives and apply derivative-based optimization solvers. Since the idea remains the same for both the objective and constraints, we will describe gradient approximation based on f(x). Overall, such approach computes an estimate g(x) of the gradient  $\nabla \phi(x)$  via forward differences (FD)

(1.1.1) 
$$g(x) = \sum_{i=1}^{N} \frac{f(x+hu_i) - f(x)}{h} u_i$$

or central differences (CD)

(1.1.2) 
$$g(x) = \sum_{i=1}^{N} \frac{f(x+hu_i) - f(x-hu_i)}{2h} u_i$$

where  $\{u_1, u_2, ..., u_N\}$  is a set of directions and h is the differencing interval or sampling radius [8]. We describe two strategies of constructing gradient approximations below.

#### 1.1.1. Finite Differences

For standard finite differences, (1.1.1) and (1.1.2) reduce to

(1.1.3) 
$$[g(x)]_i = \frac{f(x+h_i e_i) - f(x)}{h_i} \quad (FD)$$

(1.1.4) 
$$[g(x)]_i = \frac{f(x+h_i e_i) - f(x-he_i)}{2h_i} \quad (CD)$$

respectively for i = 1, ..., n, with N = n and  $u_i = e_i$  where  $e_i$  denotes the unit vector along the *i*-th coordinate.

When no noise is present in the function  $\phi$  with  $\epsilon(x) \equiv 0$ , the estimated gradient g(x)of  $\nabla \phi(x)$  can be computed by either (1.1.3) or (1.1.4) using

(1.1.5) 
$$h_i = \max(1, |x_i|)\sqrt{\epsilon_M} \quad (FD)$$

(1.1.6) 
$$h_i = \max(1, |x_i|) \sqrt[3]{\epsilon_M} \quad (CD)$$

where  $\epsilon_M$  denotes the machine precision. Such choice of differencing is designed to handle roundoff errors and is common in practice [7]. The central differences scheme (1.1.4) requires more function evaluations but provides a more accurate estimate of the gradient than the forward differences (1.1.3) scheme.

However, in the presence of noise in the function, the choice of the finite differencing interval can be delicate. To see this, consider the case of applying forward differences to a function with bounded noise. By Taylor expansion and (1.0.3), we have

(1.1.7) 
$$\mathbb{E}[([g(x)]_i - [\nabla \phi(x)]_i)^2] \le \frac{L_i^2 h_i^2}{4} + \frac{2\epsilon_f^2}{h_i^2}$$

where  $L_i$  is an upper bound for  $|e_i^T \nabla^2 \phi(x) e_i|$  [71]. Therefore, the choice of the differencing interval  $h_i$  must carefully balance the truncation error and the noise. If  $h_i$  is chosen too large, the truncation error resulting from the first term on the right-hand side of (1.1.7) damages the accuracy of the gradient approximation. On the other hand, if  $h_i$  is too small, the noise term dominates the approximation error. A near-optimal choice of the forward differencing interval can be obtained by minimizing the right-hand side of (1.1.7). With the noise level  $\epsilon_f$  and bound of second derivative  $L_i$  in hand, we can obtain

(1.1.8) 
$$h_i = \sqrt[4]{8} \sqrt{\frac{\epsilon_f}{L_i}}$$

Similarly, one can compute the expected error on the central differences case and the resulting differencing interval is

(1.1.9) 
$$h_i = \sqrt[3]{\frac{3\epsilon_f}{M_i}}$$
where  $M_i$  is the bound on the third derivative along the *i*-th coordinate direction.

Since the bounds on the second and third derivative are not available, one can estimate them by employing the procedure proposed by Gill et. al or Moré and Wild [40, 71] in practice. A more robust procedure that avoids direct estimation of the bounds is proposed in [92].

Finite differences enjoy two appealing features as they can be incorporated into existing derivative-based optimization solvers and can be easily parallelized in the function evaluations. However, when noise is present, one must carefully choose the differencing interval, which requires knowledge of the noise level and bounds on the higher-order derivatives. Kelly developed an implicit filtering method with diminishing differencing intervals, assuming that the noise diminishes as the iterate approaches the solution [21, 55]. Berahas et al. proposed a finite-difference-based L-BFGS method for solving problems with non-diminishing noise, combining strategies for noise level and bound on higher-order derivatives [7]. Such requirements call for a sophisticated and robust procedure to further improve the performance of finite-difference-based methods.

#### 1.1.2. Gaussian Smoothing

An alternative for approximating the gradient is to employ samplings along random directions, with the most popular choice being Gaussian directions [73]. Such approach is referred to as Gaussian Smoothing and has been proved efficient for applications in reinforcement learning [22, 88]. In particular, it considers (1.1.3) or (1.1.4) by sampling standard Gaussian directions  $u_i \sim \mathcal{N}(0, I)$  for i = 1, ..., N. To motivate such choice, consider the Gaussian smoothed version of function f:

$$f_h(x) = \mathbb{E}_{u \sim \mathcal{N}(0,I)}[f(x+hu)] = \int_{\mathbb{R}^n} f(x+hu)\pi(u|0,I)du$$

where  $\pi(u|0, I)$  denotes the density function of the standard Gaussian distribution evaluated at u. It can be shown that, the gradient of smoothed function F(x) can be computed as

$$\nabla f_h(x) = \frac{1}{h} \mathbb{E}_{u \sim \mathcal{N}(0,I)} [f(x+hu)u] = \frac{1}{h} \mathbb{E}_{u \sim \mathcal{N}(0,I)} [(f(x+hu) - f(x))u]$$

Therefore, by employing the idea of sample average approximations, one can obtain (1.1.3) for the forward difference case. We can verify that, when noise is not present, it follows that for  $h_i \to 0$ ,

$$\mathbb{E}_{u \sim \mathcal{N}(0,I)}[g(x)] = \nabla \phi(x).$$

When the function evaluation is noisy, similar as in the finite differences case, there are two sources of error. Because of the Gaussian distributed directions,  $[g(x)]_i$  can be arbitrarily large. Additionally, the performance of this approach is sensitive to the choice of the sampling radius h. As illustrated in [9], the randomized gradient approximation can perform significantly worse than the deterministic methods. See [8, 9] for further discussions on gradient approximations.

#### 1.2. Interpolation-Based Optimization (IBO) Methods

Interpolation-based optimization (IBO) methods have been extensively studied for unconstrained DFO problems due to their robustness and efficiency; see [29] for a thorough description. IBO methods typically construct a quadratic model of the objective via linear interpolation and generate the next iterate using a trust-region framework. While finite differences require at least n + 1 function evaluations at every iteration, IBO methods can make progress with only one function evaluation. However, it demands higher algebra cost due to interpolation. We provide a basic description of a general unconstrained IBO method in Algorithm 1.1, conceding that the details may vary for different IBO methods.

At each iteration k, IBO methods construct a quadratic model of the objective around the current iterate  $x_k$ 

(1.2.1) 
$$m_k(x_k + s) = f(x_k) + g_k^T s + \frac{1}{2} s^T H_k s$$

where  $g_k$  and  $H_k$  are computed by interpolating points in the poised interpolation set  $Y_k = \{y_0, y_1, y_2, ..., y_m\} \subseteq \mathbb{R}^n$ :

$$f(y_i) = m(y_i), \quad i = 0, 1, ..., m$$

This can be done uniquely with a cost of  $\mathcal{O}(n^6)$  when m = (n+1)(n+2)/2. The cost can be reduced to  $O(m^2)$  by employing a minimum Frobenius norm update of the Hessian approximation [78].

IBO methods aim to approximate the objective within a neighborhood of the current iterate  $x_k$ 

(1.2.2) 
$$B(x_k, \Delta_k) = \{ x \in \mathbb{R}^n | || x - x_k ||_2 \le \Delta_k \},\$$

where  $\Delta_k$  is referred to as the trust region radius. Note that other norms could also be used in (1.2.2). For unconstrained optimization, a trust region subproblem is solved at each iteration to obtain a candidate step  $s_k$  for the next iterate

(1.2.3) 
$$\min_{s \in B(0,\Delta_k)} m_k(x_k + s)$$

By evaluating the ratio test

(1.2.4) 
$$\rho_k = \frac{f(x_k) - f(x_k + s_k)}{m_k(x_k) - m_k(x_k + s_k)},$$

the algorithm accepts the step  $s_k$  if  $\rho_k$  is greater than or equal to a nonnegative user-defined value. The next iterate is then defined as  $x_k + s_k$  and the trust region  $\Delta_k$  may be increased. Otherwise, the method rejects the step and checks whether the interpolation set  $Y_k$  is sufficiently poised. The interpolation set  $Y_k$  is maintained and updated at every iteration to reflect the geometry of the objective around the current iterate  $x_k$ . It is required to satisfy certain geometry conditions to ensure quality of the interpolated. If  $Y_k$  is not well-poised, the algorithm improves the interpolation set for better model quality and leaves  $\Delta_k$  unchanged to avoid premature shrinkage of the trust region. If the model is sufficiently accurate, the algorithm follows classical trust region methods by reducing the trust region [29].

## Algorithm 1.1. IBO algorithm framework for solving unconstrained DFO problems.

1: Parameters:  $0 \le \eta < 1$ , and  $\Delta_0 > 0$ .

- 2: Choose  $x_0$ , a starting point, and construct an initial poised interpolation set  $Y_0$ .
- 3: Let k = 0.
- 4: while no convergence test is satisfied do
- 5: Build a local (quadratic) model (1.2.1) using interpolation set  $Y_k$ .
- 6: Compute a step  $s_k$  by solving the trust region subproblem (1.2.3).
- 7: Compute the ratio defined by (1.2.4).

8: if 
$$\rho_k \ge \eta$$
 then

- 9: Set  $x_{k+1} = x_k + s_k$ .
- 10: Choose  $\Delta_{k+1} \ge \Delta_k$ .
- 11: Update  $Y_k$  to include  $x_{k+1}$ .
- 12: else if  $Y_k$  needs to be improved then

13: Set 
$$x_{k+1} = x_k$$

- 14: Set  $\Delta_{k+1} = \Delta_k$ .
- 15: Improve  $Y_k$  using a geometry-improving procedure.
- 16: else
- 17: Set  $x_{k+1} = x_k$ .
- 18: Choose  $\Delta_{k+1} < \Delta_k$ .
- 19: Update  $Y_k$  to include  $x_k + s_k$ .
- 20: end if

 $21: \qquad Y_{k+1} = Y_k$ 

22: k = k + 1

## 23: end while

Unconstrained IBO methods have been extended by Powell to solve bound constrained [79], linear constrained [76, 80] and inequality constrained problems [77]. When the constraints are convex and the projection operator is known, [24, 51] propose to enforce the constraints in the trust region subproblem (1.2.3) and establish convergence analysis. A few works have been proposed to handle general (potentially nonconvex) constraints [11, 23, 26, 83]. However, no convergence result is known for general constrained IBO methods.

Surprisingly, as demonstrated in [93], IBO methods can be directly applied to noisy functions without modifications to their noiseless counterparts and achieve competent performance. Some noisy variants of IBO methods handle unconstrained stochastic DFO problems by maintaining model accuracy with high probability [17, 20]. Other works explore different strategies for extending unconstrained IBO to noisy problems, including [18, 99], in which the ratio test is modified or restarts are employed.

#### 1.3. Covariance Matrix Adaptation - Evolution Strategy (CMA-ES)

Covariance Matrix Adaptation Evolution Strategy (CMA-ES) is a randomized DFO method that conducts iterative sampling based on a multivariate Gaussian distribution with adaptive parameters. Many variants of this approach have been proposed to adopt different sampling and parameter update strategies, although the main idea remains the same [46].

To elaborate, at each iteration k, the algorithm CMA-ES samples new candidate solutions from a multivariate Gaussian distribution  $\mathcal{N}(m_k, \sigma_k^2 C_k)$ , where  $\sigma_k$  is the stepsize,  $m_k$  is the mean vector and  $C_k$  is the covariance matrix. In contrast to Gaussian Smoothing, in which the sampling distribution is fixed throughout the optimization procedure, this algorithm updates the parameters of the Gaussian distribution using information from past iterations: 1) the mean vector is computed as a weighted average of the best samples in the current iteration 2) the covariance matrix is modified via low-rank updates such that the variance along promising directions (i.e., directions that are more likely to reduce the objective) is increased 3) the stepsize  $\sigma_k$  is updated based on the cumulative path from past iterations to facilitate convergence. We provide a formal algorithm statement for basic CMA-ES in Algorithm 1.2.

After sampling  $\lambda$  new candidates  $\{x_1, ..., x_\lambda\}$  from the given Gaussian distribution  $\mathcal{N}(m_k, \sigma_k^2 C_k)$  at iteration k, the algorithm evaluate at those points and rank them such that  $f(x_{(1)}) \leq f(x_{(2)}) \leq ... \leq f(x_{(\lambda)})$ . The new mean  $m_{k+1}$  is then moved to a weighted average of the best  $\mu$  sampled points at the current iteration

(1.3.1) 
$$m_{k+1} = \sum_{i=1}^{\lambda} w_i x_{(i)} = m_{(k)} + \sum_{i=1}^{\lambda} w_i (x_{(i)} - m^{(k)}).$$

where  $w_i$  are pre-specified weights such that  $\sum_{i=1}^{\mu} w_i = 1, w_1 \ge w_2 \ge \dots \ge w_{\mu} > 0.$ 

Since reestimating the covariance matrix from scratch using  $\{x_1, ..., x_\lambda\}$  can be unstable, the algorithm performs a rank- $\mu$  update that exploits information from all past samples, and a rank-one update, which captures the correlation between consecutive steps in the past. The scale of the distribution depends on the stepsize  $\sigma_k$ , which depends on the cumulative path in the past. Intuitively, the length of the cumulative path is shorter if the past steps cancel out each and longer if they point at similar directions.

1: Parameters:  $c_{\sigma}, d_{\sigma}, c_c, c_1, c_{\mu}$  (parameters for updates);  $\lambda \geq \mu > 0$  (parameters for sampling);  $w_i : \sum_{i=1}^{\mu} w_i = 1, w_1 \ge w_2 \ge \dots \ge w_{\mu} > 0$  (weight parameters) 2: Initialization:  $m_0 \in \mathbb{R}^n, C_0 = I$  ,  $\sigma_0 > 0; \, p_{\sigma,0} = 0, p_{c,0} = 0; \, k = 0$ 3: while Not Terminated do Sampling: 4: Eigendecomposition:  $C_k = BD^2B^T$ 5: Sample  $\lambda$  points  $z_j \sim \mathcal{N}(0, I)$  for  $j = 1, ..., \lambda$ 6:  $y_i = BDz_k \sim \mathcal{N}(0, C_k)$ 7:  $x_i = m_k + \sigma_k y_i \sim \mathcal{N}(m_k, \sigma_k^2 C_k)$ 8: 9: **Rank and Recombination:** Evaluate  $f(x_1), \ldots, f(x_{\lambda})$ 10:Select the best  $\mu$  points:  $x_{(1)}, ..., x_{(\mu)}$  such that  $f(x_{(1)}) \leq f(x_{(2)}) \leq ... \leq f(x_{(\lambda)})$ 11: Let  $\bar{y}_w = \sum_{i=1}^{\mu} w_i y_{(i)}$  where  $y_{(i)}$  corresponds to  $x_{(i)}$ 12:Update parameters 13:Update *m*: 14: $m_{k+1} = m_k + \sum_{i=1}^{\lambda} w_i(x_{(i)} - m^{(k)})$ 15:Update stepsize: 16: $p_{\sigma,k+1} = (1 - c_{\sigma})p_{\sigma,k} + \sqrt{c_{\sigma}(2 - c_{\sigma})\mu_{\text{eff}}}(C_k)^{\frac{1}{2}}\bar{y}_w$ 17: $\sigma_{k+1} = \sigma_k \exp\left\{\frac{c_{\sigma}}{d_{\sigma}} \left(\frac{\|p_{\sigma,k+1}\|}{\mathbb{E}[\|\mathcal{N}(0,I)\|]} - 1\right)\right\}$ 18:Update Covariance Matrix: 19: $p_{c,k+1} = (1 - c_c)p_{\sigma,k} + h_{\sigma}\sqrt{c_c(2 - c_c)\mu_{\text{eff}}}\bar{y}_w$ 20:  $C_{k+1} = (1 - c_1 - c_\mu)C_k + c_1 p_{c,k+1} p_{c,k+1}^T + c_\mu \sum_{i=1}^{\mu} w_i y_{(i)} y_{(i)}^T$ 21:k = k + 122:23: end while

The covariance matrix  $C_k$  is designed to learn second order information of the objective, although no analysis has been established [45]. This method has been widely applied to practical applications such as Reinforcement Learning [98] and hyperparameter tuning [64]. The method does not involve any approximations to the gradient of the objective function and relies solely on rankings to perform updates. CMA-ES enjoys the benefit of having few algorithm parameters and is believed to work well on difficult problems that are ill-conditioned, require many function evaluations to solve and have relatively high dimensionality. It has been found to outperform 31 DFO methods on a collection of black-box functions [47], including IBO methods. However, when only a small number of function evaluations are available, CMA-ES can be inferior to other methods such as IBO. Convergence analysis has been established on a modified version of CMA-ES [36].

#### 1.4. Bayesian Optimization

Bayesian Optimization (BO) methods have been proposed to solve DFO problems with expensive objective evaluations and simple constraints. It has been popular for hyperparameter tuning in the machine learning community. See [39] for a complete overview for BO methods.

The basic idea of BO consists of two steps: i) modeling the objective function and ii) optimizing the acquisition function to obtain the next iterate. BO methods employ Gaussian Process (GP) regression to build a surrogate model for the objective function and quantify the uncertainty in the function. Given a set of evaluated points, GP assumes that the objective values are drawn from some multivariate Gaussian distribution with a mean vector and covariance matrix. Utilizing the properties of Gaussian distribution, one can compute a conditional distribution, or in the nomenclature of BO, posterior probability distribution, based on all available data for a new point that hasn't been evaluated. The mean vector and kernel for the covariance matrix are hyperparameters for BO and remains flexible as long as several properties are satisfied; see [39]. The acquisition function is then defined based on the poster distribution, with careful balance between exploration and exploitation. Common choices of acquisition functions include expected improvement, upper confidence bound and knowledge gradient [39]. The general idea indicates that the next iterate should either have a better expected objective (Exploitation) or high variance (Exploration). This is important for BO as the method is designed for global optimization.

When the objective contains stochastic Gaussian noise, BO can be easily adapted by incorporating such information in the kernel function. However, BO methods are only efficient for problems with small number of variables (up to 20) as it scales poorly with the dimension due to the GP component. The performance of the methods also relies on proper choice of the GP component, e.g., picking a good kernel, and can be volatile.

## 1.5. An Application: Prompting Large Language Models

DFO problems arise in many applications from engineering to machine learning, especially when the objective results from black-box simulations or model evaluation. Examples include hyperparameter tuning [58], adversarial attacks [19], traffic model calibration [4], parameter estimation [103] to name a few.

One recent application comes from the field of Natural Language Processing, where a general pre-trained large language model (LLM) needs to be adapted to address individual language tasks. Given the massive size of the LLMs, one approach that has gained increasing attention is to append a prompt, e.g., task description, to every input sequence, while keeping the language model parameters frozen. This approach is appealing as it eliminates the need for storing separate copies of the model parameters for each task. Unfortunately, this prompt design process is often done in a manual trial-and-error manner, yielding results that may not be as effective as model tuning. To address this, one strategy is to append a continuous prompt p to every input sequences  $x_i$  and perform optimization on this continuous vector without changing parameters of the massive language model.

Suppose we wish to classify the correct label  $y_i$  for each input sequence  $x_i$  using loss function  $\mathcal{L}$ . The prompt optimization problem can thus be formulated as

(1.5.1) 
$$\min_{p \in \mathbb{R}^D} \frac{1}{N} \sum_{i=1}^N \mathcal{L}(f(p; x_i), y_i)$$

where f is the pre-trained language model and N the number of training samples. With the massive size of the LLMs, computing the derivatives of the objective function via backpropagation can be prohibitively memory intensive. In addition, some LLMs are released as API services that only allows input and output manipulations, the users cannot perform backpropagation to obtain derivatives of the objective function in (1.5.1). As a result, this problem must be considered as a DFO problem that has no access to the derivatives of the functions. Additionally, the problem is a noisy DFO problem with deterministic errors as LLMs are often implemented in half or single precision. In Chapter 5, we will present our numerical experiments to improve the performance of LLMs on solving a collection of language tasks in a DFO setting.

## CHAPTER 2

# On the Numerical Performance of Finite-Difference Based Methods for Derivative-Free Optimization

#### 2.1. Introduction

The problem of minimizing a nonlinear objective function when gradient information is not available has received much attention in the literature; see e.g., [29, 59] and the references therein. A variety of methods have been developed for unconstrained optimization, and some of these methods have been extended to deal with constraints. The important benchmarking paper by Moré and Wild [69] showed that traditional methods, such as the Nelder-Mead simplex method [72] and a leading pattern-search method [44], are not competitive with the interpolation-based trust-region approach pioneered by Powell [78] and developed concurrently by several other authors [29]. The advantages of Powell's approach reported in [69] were observed for both smooth and nonsmooth problems, as well as noisy and noiseless objective functions. Rios and Sahinidis [85] confirmed the findings of Moré and Wild concerning the inefficiency of traditional methods based on Nelder-Mead or pattern searches. Based on these studies, we regard the interpolation-based approach of Powell as a leading method for derivative-free optimization (DFO).

There is, however, an alternative approach for DFO that has been largely neglected in the nonlinear optimization literature. It consists of approximating derivatives using finite differences, and using them within standard derivative-based nonlinear optimization algorithms. Many of the papers in the DFO literature dismiss this approach at the outset as being too expensive in terms of functions evaluations or as ineffective in the presence of noise. As a result of this prevalent view, the vast majority of papers on DFO methods do not present comparisons against a finite-difference approach. We believe that if such comparisons had been made, particularly in the noiseless setting, research in the field would have followed a different trajectory. Ironically, as pointed out by Berahas et al. [7], the finite-difference approach is widely used by practitioners for noiseless problems, often unwittingly, as many established optimization codes invoke a finite-difference option when derivatives are not provided. A disconnect occurred between research and practice, and no systematic effort was made to bridge this gap.

This paper builds upon Berahas et al. [7] and Nesterov and Spokoiny [73], and aims to bring the gradient approximation approach to the forefront of DFO research by illustrating its performance on a variety of unconstrained and constrained problems, with and without noise. Gradient approximations can be computed using finite differences or random sampling techniques [73]. In this paper, we study only the former approach.

As is well known, the use of finite differences is delicate in the presence of noise, and unless carefully implemented, can lead to inefficiencies or outright failure. Nevertheless, when the standard deviation of the noise is known or can be estimated, the approximation of derivatives can be placed on a solid theoretical footing. One can view this task as the computation of derivatives of a smoothed function [73], or as the computation of an estimator that minimizes a mean squared error [71].

Finite-difference-based methods for DFO enjoy two appealing features. They can easily exploit parallelism in the evaluation of functions during finite differencing, and they can be built using existing software for constrained and unconstrained optimization, sometimes rather easily. This obviates the need to redesign existing unconstrained DFO methods to handle more general problems, an effort that has taken two decades for interpolationbased trust-region methods [25–29, 76, 79, 80], and is yet incomplete. The strategy often suggested of simply applying an unconstrained DFO method to a penalty or augmented Lagrangian reformulation will not yield an effective general purpose method, as is known from modern research in nonlinear programming; see e.g. [75, chapter 17].

To gauge the efficiency of the finite-difference approach for DFO, we implemented it within three derivative-based codes: L-BFGS [75] for unconstrained optimization; LMDER [68] for nonlinear least squares; and KNITRO [16] for inequality constrained optimization. In L-BFGS we employ a relaxed line search for which convergence results in the presence of (errors) or noise are established in [7, 53, 104]. LMDER and KNITRO were not modified in any way, and thus our tests provide baseline performance as noise-tolerant variations of these methods should yield improved performance. (Convergence results have not been established for LMDER and KNITRO in the noisy setting, but [99] analyzes a modification of trust region methods that is relevant to LMDER.) We tested the aforementioned codes against three established codes designed specifically for DFO: NEWUOA [78] for general unconstrained optimization; DFO-LS [18] for nonlinear least squares; and COBYLA [77] for inequality constrained optimization. A large number of experiments were performed in this study. In sections 2.2, 2.3, and 2.4, we display graphs or tables that attempt to summarize the findings as well as possible. The complete set of results, as well as fine details of implementation, are presented in a companion technical report [94].

#### 2.1.1. Literature Review

Gill et al. [40] proposed an adaptive approach for computing the difference interval h, assuming that a bound on the errors in the objective function is known. This work, inspired by earlier research by Lyness [65], pays careful attention to the estimation of bounds on the second (or third) derivatives, since these bounds are needed to obtain an accurate estimate of h. Examples are presented illustrating cases when the approach may fail. To our knowledge, [40] is the first in-depth study (in the context of optimization) of finite-difference gradient approximations in the presence of errors. There was, however, no follow-up on the application of these techniques for derivative-free optimization.

The paper that influenced our work the most is by Moré and Wild [71], who discuss how to choose the differencing interval as a function of the noise level and a bound on the second (or third) derivative, so as to obtain nearly optimal gradient estimates. The noise level, defined as the standard deviation of the noise, can be estimated by sampling (in the case of stochastic noise) or using a table of differences [70] (in the case of computational noise). The authors propose ECnoise, a practical procedure for estimating stochastic or computational noise [70]. They also give some attention to the practical estimation of the second derivative.

Nesterov and Spokoiny [73] analyze first-order methods with randomized gradient approximations. They establish worst case complexity bounds for nonsmooth and smooth problems. Numerical experiments with these types of methods are reported in [8, 10]. Kelley et al. [21, 55] propose a finite-difference BFGS method, called implicit filtering, designed for the case when noise can be diminished at any iteration, as needed. In that approach, the finite-difference interval decreases monotonically. The book by Conn, Scheinberg and Vicente [29] gives a thorough treatment of the interpolation-based trust-region approach for DFO. It presents foundational theoretical results as well as detailed algorithmic descriptions. Larson, Menickelly and Wild [59] give a comprehensive review of DFO methods as of 2018. Their survey covers deterministic and randomized methods, noisy and noiseless objective functions, problem structures such as least squares and empirical risk minimization, and various types of constraints.

Audet and Hare [2] review direct-search and pattern-search methods, and describe a variety of practical applications solved with MADS, and NOMAD. Neumaier [74] reviews methods endowed with convergence guarantees, with an emphasis on global optimization. Kimiaei [57] proposes a randomized method, called VSBBON that implements a noisy line search and employs quadratic models in subspaces determined adaptively.

#### 2.1.2. Contributions of this Paper

To our knowledge, this is the first systematic investigation into the practical performance of finite-difference-based DFO methods relative to established techniques, across a range of problems, with and without noise in the functions. The main contributions of this paper can be summarized as follows. We use the acronym "FD-DFO method" for a method that employs some form of finite differences to approximate the gradient of the objective function and (possibly) the constraints.

(1) For noiseless functions, we found that the FD-DFO methods are at least as efficient, if not superior, to established methods, across all three categories of problems, without the use of sophisticated procedures for determining the finite-difference interval.

- (2) For noisy functions, we observed that NEWUOA is more efficient and accurate than the finite-difference L-BFGS method for unconstrained optimization, but not by a wide margin. DFO-LS is comparable to the finite-difference version of LMDER, for least-squares problems. A simple finite-difference version of KNITRO has comparable performance with COBYLA, for inequality-constrained problems.
- (3) The differencing formulas used in our experiments performed well on most, but not all noisy problems. More sophisticated techniques for estimating bounds on the second (or third) derivatives will improve the accuracy of the differencing interval, as well as the overall performance of FD-DFO methods, on noisy problems.

Conclusions (1) and (2) are based on sequential execution. If function evaluations can be parallelized, we expect FD-based methods to have a clear advantage over existing DFO methods, for all classes of problems. On the other hand, we should note that the interpolation-based trust-region methods we tested (NEWUOA and DFO-LS) proved to be more robust in the presence of noise than we expected. These methods do not require knowledge of the noise level in the objective function and yet performed reliably for most levels of noise, suggesting that the internal logic of the algorithm normally reacts correctly to the noise inherent in the problem (although rare failures were observed).

We also comment on some limitations of this work. For each problem class, we employed only one established DFO code to benchmark the efficiency of the corresponding FD-DFO method. We found it essential to work with a small number of codes that we could understand well, but benchmarking other software is desirable. Another limitation of this study is that it considers only one model of noise: additive uniformly-distributed bounded noise. Although we experimented with normal noise and observed that the FD-based codes worked well, we do not know how robust these methods are for heavy tail distributions or anisotropic noise. We focused on just one simple model of noise because this already raised some important algorithmic questions that need to be resolved to make FD-DFO methods highly reliable in practice.

As our focus was on scalable local optimization methods, we did not consider global optimization methods, such as Bayesian optimization, surrogate optimization, evolutionary methods [37, 39, 45], or other global optimization methods that employ restarts, grid searches or surrogate models [74, 87].

#### 2.2. Unconstrained Optimization

Let us consider the solution of unconstrained optimization problems of the form

(2.2.1) 
$$\min_{x \in \mathbb{R}^n} \phi(x),$$

given only noisy evaluations,

$$f(x) = \phi(x) + \epsilon(x).$$

Here,  $\epsilon(x)$  denotes (deterministic) computational error or the realization of a random variable at x representing noise. The function  $\phi$  is assumed to be smooth. We compare the performance of NEWUOA and L-BFGS with finite-difference gradients on 73 unconstrained **CUTEst** problems [42], varying the dimension of each problem up to  $n \leq 300$  whenever possible. All experiments were run in double precision. The methods tested in our experiments are as follows.

- NEWUOA: A model-based trust-region derivative-free algorithm that forms quadratic models using interpolation of function values, combined with a minimum Frobeniusnorm update of the Hessian approximation; see Powell [78]. We called the code in Python 3.7 through PDFO developed by Ragonneau and Zhang [84]. We used the default settings in NEWUOA, which in particular, set the number of interpolation points to 2n + 1.
- FD-L-BFGS: A finite-difference implementation of the limited-memory BFGS algorithm [75] with a bisection Armijo-Wolfe line search. We use a memory of size 10. Forward- and central-difference options are tested. At intermediate trial points generated during the line search, the directional derivative is computed via forward- or central-differences along the direction of interest.

Instead of NEWUOA we could have used DFOTR [5], which in our experience is often competitive with NEWUOA in terms of function evaluations, but has much higher periteration cost.

We first consider the case when noise is not present and later study the effect of noise on the performance of the algorithms.

#### 2.2.1. Experiments on Noiseless Functions

In the first set of experiments, we let  $\epsilon(x) \equiv 0$ , so that the only errors in the function evaluations are due to machine roundoff, and we let  $\epsilon_M$  denote unit roundoff. The approximate gradient  $g(x) \in \mathbb{R}^n$  of the objective function  $\phi(x)$ , computed by finite differencing, is given as:

$$(2.2.3)$$

$$[g(x)]_i = \frac{f(x+h_i e_i) - f(x)}{h_i}, \qquad h_i = \max\{1, |[x]_i|\}\sqrt{\epsilon_M}; \text{ (forward differencing)}$$

(2.2.4)

$$[g(x)]_{i} = \frac{f(x+h_{i}e_{i}) - f(x-h_{i}e_{i})}{2h_{i}}, \quad h_{i} = \max\{1, |[x]_{i}|\}\sqrt[3]{\epsilon_{M}}. \quad (\text{central differencing})$$

The directional derivative  $D_p \phi(x) = \nabla \phi(x)^T p$  of  $\phi$  along a direction p is required within the Armijo-Wolfe line search employed by L-BFGS. It is approximated as:

(2.2.5) 
$$g(x;p) = \frac{f(x+hp_u) - f(x)}{h} ||p||, \qquad h = \sqrt{\epsilon_M}; \quad \text{(forward differencing)}$$

(2.2.6) 
$$g(x;p) = \frac{f(x+hp_u) - f(x-hp_u)}{2h} ||p||, \quad h = \sqrt[3]{\epsilon_M}; \quad \text{(central differencing)}$$

where  $p_u = p/||p||$  is the normalized direction.

These choices of h can be improved by including contributions of the second and third derivatives, respectively, as discussed in the next subsection. However, we found that in the noiseless setting, and for our test functions, such a refinement is not needed to make the finite-difference L-BFGS approach competitive.

The algorithms are terminated when either

(2.2.7) 
$$\phi(x_k) - \phi^* \le \tau \cdot \max\{1, |\phi^*|\},\$$

with  $\tau = 10^{-6}$ , or when the limit of  $500 \times n$  function evaluations is reached. The optimal value  $\phi^*$  is determined by running BFGS with exact gradients (provided by CUTEst [42])

until no more progress can be made on the function. Complete numerical results are given in the technical report [94].

In Figure 2.1, we summarize the results using *log-ratio profiles* proposed by Morales [67], which in this case report the quantity

(2.2.8) 
$$\log_2\left(\frac{\text{evals}_{\text{LBFGS}}}{\text{evals}_{\text{NEWUOA}}}\right),$$

where  $evals_{LBFGS}$  and  $evals_{NEWUOA}$  denote the total number of function evaluations for FD-L-BFGS and NEWUOA to satisfy (2.2.7) or reach the maximum number of function evaluations. In the figures, the ratios (2.2.8) are plotted in increasing order. Thus, the larger the shaded region, the more successful a method.

Overall, we observe in these tests (with  $\tau = 10^{-6}$ ) that forward-difference L-BFGS outperforms NEWUOA on the majority of problems in terms of function evaluations. This is perhaps surprising because NEWUOA was designed to be parsimonious in terms of function evaluations, whereas finite-difference L-BFGS requires *n* function evaluations per iteration. It is also notable that this is achieved without any additional information (for example, the Lipschitz constant of the gradient for forward differencing) to squeeze out the best possible accuracy of the gradient in the finite-difference L-BFGS approach. As expected, central-difference L-BFGS requires significantly more function evaluations than the forward-difference option, and does not provide significant benefit in terms of solution accuracy for the majority of the problems. In particular, when run in double precision, forward-difference L-BFGS is able to converge to the same tolerance that one would expect with analytical gradients.



Figure 2.1. Efficiency, Noiseless Case. Log-ratio profiles for the total number of function evaluations to achieve (2.2.7) with  $\epsilon(x) = 0$ . The left figure compares forward-difference L-BFGS with NEWUOA, and the right figure compares central-difference L-BFGS with NEWUOA.

In terms of CPU time, NEWUOA's execution time grows much faster with the problem dimension than L-BFGS' because one iteration of NEWUOA requires  $O(n^2)$  flops, whereas the iteration cost of L-BFGS is O(n) flops and all test functions are inexpensive to evaluate. Across all the problems, we observe that when  $n \approx 100$ , NEWUOA can take at least 5-10 times longer than L-BFGS in terms of wall-clock time.

While NEWUOA is an inherently sequential algorithm, finite-difference L-BFGS offers ample opportunities for parallelism when computing the finite-difference approximation to the gradient. This is an often overlooked benefit of finite-difference-based methods in the DFO literature, as we are not aware of implementations of model-based trust-region methods that benefit significantly from parallelism.

These results also indicate that finite-difference L-BFGS is likely to be much more efficient than direct search methods (Nelder-Mead or pattern search) on noiseless problems since the latter are not competitive with the interpolation-based trust region approach, as mentioned above.

#### 2.2.2. Experiments on Noisy Functions

In this set of experiments, we synthetically inject uniform stochastic noise into the objective function. In particular, we sample  $\epsilon(x) \sim \sigma_f U(-\sqrt{3}, \sqrt{3})$  i.i.d. independent of x, where  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ . By construction of  $\epsilon(x)$ , we have that  $\sigma_f^2 = \mathbb{E}[\epsilon(x)^2]$ . We refer to standard deviation of the noise  $\sigma_f$  as the *noise level*.

We employ a more precise formula for the finite-difference interval than in the noiseless setting — one that depends both on the noise level  $\sigma_f$  and on the curvature of the function. Specifically, we compute a different  $h_i$  for each coordinate direction based on the well known results in [71]. We define

$$(2.2.9) \quad [g(x)]_i = \frac{f(x+h_ie_i) - f(x)}{h_i}, \qquad h_i = \sqrt[4]{8}\sqrt{\frac{\sigma_f}{L_i}}; \quad \text{(forward differencing)}$$

(2.2.10) 
$$[g(x)]_i = \frac{f(x+h_ie_i) - f(x-h_ie_i)}{2h_i}, \quad h_i = \sqrt[3]{\frac{3\sigma_f}{M_i}}, \quad \text{(central differencing)}$$

where  $L_i$  and  $M_i$  are bounds on the second and third derivative along the *i*-th coordinate direction  $e_i$ .

We estimate  $L_i$  using a second-order difference. Given a direction  $p \in \mathbb{R}^n$  where ||p|| = 1, we define

$$\Delta(t) = f(x+tp) - 2f(x) + f(x-tp),$$

where t is the second-order differencing interval. The Lipschitz constant L along the direction p can thus be approximated as  $L \approx \Delta(t)/t^2$ . However, as with the choice of h, we need to be careful in the selection of t, and for this purpose we employ an iterative

technique proposed by Moré and Wild [71], whose goal is to find an interval t that satisfies

(2.2.11) 
$$|\Delta(t)| \ge \tau_1 \epsilon_f, \qquad \tau_1 \gg 1$$

(2.2.12) 
$$|f(x \pm tp) - f(x)| \le \tau_2 \max\{|f(x)|, |f(x \pm tp)|\}, \quad \tau_2 \in (0, 1).$$

These heuristic conditions aim to ensure that t is neither too small nor too large; see [71] for a full description of the Moré-Wild (MW) technique. This technique cannot, however, be guaranteed to find a t that satisfies (2.2.11), (2.2.12). To account for this, we developed the following procedure for estimating the constants  $L_i$  for forward differencing.

## **Procedure I.** Adaptive Estimation of L

1. At the first iteration of the L-BFGS method, invoke the MW procedure to compute  $t_i$ , for i = 1, ..., n. If such a  $t_i$  can be found to satisfy (2.2.11), (2.2.12), set  $L_i = \max\{10^{-1}, |\Delta(t_i)|/t_i^2\}$  in (2.2.9). Otherwise, set it to  $L_i = 10^{-1}$ . Store the  $L_i$  in a vector **L**. To calculate the directional derivative in the Armijo-Wolfe line search, set the Lipschitz constant to  $L = ||\mathbf{L}||/\sqrt{n}$ .

2. If at any iteration of the L-BFGS algorithm the line search returns a steplength  $\alpha_k < 0.5$ , then re-estimate the vector **L** by calling the MW procedure as in step 1 at the current iterate  $x_k$ .

We modify the line search in FD-L-BFGS to better handle noise; see Shi et al. [91] for details. No other changes were made to the L-BFGS method. It is well known, that when the gradient is sufficiently accurate, L-BFGS generates well-scaled directions so that  $\alpha_k = 1$  is acceptable. Thus, the occurrence of  $\alpha_k < 0.5$  is viewed as an indication that the curvature of the problem may have changed, and should be re-estimated. The function evaluations performed in the estimation of  $\mathbf{L}$  will be accounted for in the numerical results presented below.

The technique for computing differencing interval for central differences is described in [94].

**2.2.2.1.** Accuracy. We first compare the accuracy achieved by each algorithm, as measured by the optimality gap  $\phi(x_k) - \phi^*$ . We do so by running NEWUOA until  $\rho = \rho_{\text{end}} = 10^{-6}$ , and running the FD-L-BFGS method until the objective function could not be improved over 5 consecutive iterations. In Figure 2.2, we report the log-ratio profile

(2.2.13) 
$$\log_2\left(\frac{\phi_{\text{LBFGS}} - \phi^*}{\phi_{\text{NEWUOA}} - \phi^*}\right)$$

for  $\sigma_f = 10^{-5}$ , where  $\phi_{\text{LBFGS}}$ ,  $\phi_{\text{NEWUOA}}$  denote the lowest objective achieved by each method. (The results are representative of those obtained for  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-7}\}$ .) Since we include only runs where both solvers converged to the same local minimizer, the differences in both the numerator and denominator are nonnegative.



Figure 2.2. Accuracy, Noisy Case for  $\sigma_f = 10^{-5}$ . Log-ratio optimality gap profiles comparing NEWUOA against forward-difference L-BFGS (left) and central-difference L-BFGS (right).

As seen in Figure 2.2, NEWUOA achieves higher accuracy in the solution than forwarddifference L-BFGS, while central-difference L-BFGS yields far better accuracy than both. It is not surprising that central differencing yields much higher accuracy than forward differencing since the noise levels of their gradient approximations are, respectively,  $O(\sigma_f^{2/3})$ and  $O(\sigma_f^{1/2})$ . On the other hand, it is not straightforward to analyze the error contained in the gradient approximation constructed by NEWUOA, and in turn its final accuracy. To try to shed some light into this question, we tested NEWUOA using only n + 1 interpolation points and observed that it now lags behind FD-L-BFGS in terms of accuracy. More generally, our tests suggest that the choice p = 2n + 1 recommended by Powell strikes the right balance between accuracy in the solution and the speed of algorithm.

**2.2.2.2. Efficiency.** We now report the number of function evaluations required to achieve

(2.2.14) 
$$\phi(x_k) - \overline{\phi}^* \le \tau \cdot (\phi(x_0) - \overline{\phi}^*),$$

for varied  $\tau$ . Here,  $\tilde{\phi}^*$  denotes the best solution obtained by NEWUOA. In Figure 2.3, we report log-ratio profiles based on

(2.2.15) 
$$\log_2\left(\frac{\text{evals}_{\text{LBFGS}}}{\text{evals}_{\text{NEWUOA}}}\right).$$

We observe from this figure that NEWUOA is more efficient than forward-difference L-BFGS. The advantage is less significant for  $\tau = 10^{-2}$  but becomes pronounced for  $\tau = 10^{-6}$ , which is consistent with our earlier observation about the ability of NEWUOA to achieve higher accuracy. Central-difference L-BFGS is also less efficient overall than NEWUOA, but becomes more competitive as  $\tau$  is decreased. It is notable that NEWUOA is able to deliver such



Figure 2.3. Efficiency, Noisy Case for  $\sigma_f = 10^{-5}$ . Log-ratio profiles for the number of function evaluations to achieve (2.2.14) for  $\tau = 10^{-2}$  (left) and  $10^{-6}$  (right), comparing NEWUOA against forward-difference L-BFGS (top) and central-difference L-BFGS (bottom). These plots are representative of the other noise levels  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ .

strong performance without knowledge of the noise level of the function. The adjustment of the two trust-region radii in NEWUOA seems to be quite effective: shrinking the radii fast enough to ensure steady progress, but not so fast as to create models dominated by noise. To our knowledge, there has been no in-depth study of the *practical* behavior of NEWUOA (or similar codes) in the presence of noise; we regard this as an interesting research topic.

The results reported in this section highlight the importance of developing more effective finite difference estimation techniques than those employed in our experiments. Recent work along these lines is described in [92].

#### 2.3. Nonlinear Least Squares

In many unconstrained optimization problems, the objective function has a nonlinear least-squares form. Therefore, it is important to pay particular attention to this problem structure in the derivative-free setting. We write the problem as

$$\min_{x \in \mathbb{R}^n} \phi(x) = \frac{1}{2} \|\gamma(x)\|^2 = \frac{1}{2} \sum_{i=1}^m \gamma_i^2(x),$$

where  $\gamma : \mathbb{R}^n \to \mathbb{R}^m$  is a smooth function. We assume that the Jacobian matrix  $[J(x)]_{ij} = \frac{\partial \gamma_i(x)}{\partial x_j}$  is not available but that the individual residual functions  $\gamma_i(x)$  can be computed. More generally, the evaluation of the  $\gamma_i$  may contain noise so that the observed residuals are given by

$$r_i(x) = \gamma_i(x) + \epsilon_i(x), \quad i = 1, \dots, m,$$

where  $\epsilon_i(x)$  models noise as in (2.2.1). Thus, the minimization of the true objective function  $\phi$  must be performed based on noisy observations  $r_i(x)$  that define the observed objective function

(2.3.1) 
$$f(x) = \frac{1}{2} ||r(x)||^2 = \frac{1}{2} \sum_{i=1}^m r_i^2(x).$$

Since noise is incorporated into each residual function, the model of noise is different from the additive noise model in the general unconstrained case discussed in the previous section. In particular, the function evaluation  $f(x) = \frac{1}{2} \sum_{i=1}^{m} \gamma_i^2(x) + 2\epsilon_i(x)\gamma_i(x) + \epsilon_i^2(x)$ contains both multiplicative and additive components of noise.

Our goal is to study the viability of methods based on finite-difference approximations to the Jacobian. To this end, we employ a classical Levenberg-Marquardt trust-region method where the Jacobian is approximated by differencing, and perform tests comparing it against a state-of-the-art DFO code designed for nonlinear least-squares problems. The rationale behind the selection of codes used in our experiments is discussed next.

## Interpolation Based Trust Region Methods.

Cartis and Roberts [86] proposed a Gauss-Newton type approach, referred to as DFO-GN, in which an approximation of the Jacobian is computed by linear interpolation using n + 1 function values at recently generated points. An improved version of DFO-GN is DFO-LS [18], which provides a variety of options and heuristics to accelerate convergence and promote a more accurate solution. The numerical results reported by Roberts et al. [18] indicate that DFO-LS is a state-of-the-art code for DFO least squares, and therefore will be used in our benchmarking.

#### Finite-Difference Gauss-Newton Method.

One can employ finite differencing to estimate the Jacobian matrix J(x) within any method for nonlinear least squares, and since this is a mature area, there are a number established solvers. We chose LMDER for our experiments, which is part of the MINPACK package [68] and is also available in the scipy library. We did not employ LMDIF, the finite-difference version of LMDER, because it does not allow the use of different differencing intervals for each of the residual functions  $r_i(x)$ ; we elaborate on this point below. Another code available in scipy is TRF [13], but our tests show that LMDER is slightly more efficient in terms of function evaluations, and tends to give higher accuracy in the solution. The code NLS, recently added to the Galahad library [43] would provide an interesting alternative. That method, however, includes a tensor to enhance the Gauss-Newton model, and since this may give it an advantage over DFO-LS, we decided to employ the more traditional code LMDER.

In summary, the solvers used in our tests are:

- LMDER: A derivative-based Levenberg-Marquardt trust-region algorithm from the MINPACK software library [68], where the finite-difference module is supplied by us. We call the code in Python 3.7 through scipy version 1.5.3, using the default parameter settings.
- DFO-LS: The most recent DFO software developed by Cartis et al. [18] for nonlinear least squares. This method uses linear interpolation to construct an approximation to the Jacobian matrix, which is then used in a Gauss-Newton-type method. We used version 1.0.2 in our experiments, with default settings except that the model.abs\_tol parameter is set to 0 to avoid early termination.

The test problems in our experiments are those used by Moré and Wild [69], which have also been employed by Roberts and Cartis [86] and Zhang et al. [105]. The 53 unconstrained problems in this test set include both zero and nonzero residual problems, with various starting points, and are all small dimensional, with  $n \leq 12$ . To measure efficiency, we regard m evaluations of individual residual components  $r_i(\cdot)$  as one function evaluation. These m evaluations of the individual residual components are not necessarily performed at the same point. We terminate the algorithms when either: i) the maximum number of function evaluations ( $500 \times n$ ) is reached, ii) an optimality gap stopping condition, specified below, is triggered; or iii) the default termination criterion of the two codes is satisfied with tolerance of  $10^{-8}$  (this controls the minimum allowed trust region radius).

#### 2.3.1. Experiments on Noiseless Functions

We first consider the noiseless case corresponding to  $\epsilon_i(x) \equiv 0$ ,  $\forall i$ . For LMDER, we estimate the Jacobian J(x) using forward differences. As before, let  $\epsilon_M$  denote machine precision. We first evaluate  $\{r(x + h_j e_j)\}_{j=1}^n$ , where  $h_j$  is defined as in (2.2.3), and compute the Jacobian estimate as

(2.3.2) 
$$[\hat{J}(x)]_{ij} = \frac{r_i(x+h_je_j) - r_i(x)}{h_j}.$$

As in the previous section, we consider log-ratio profiles to compare the efficiency and accuracy of LMDER and DFO-LS. We record

$$\log_2\left(\frac{\text{evals}_{\text{LMDER}}}{\text{evals}_{\text{DFOLS}}}\right) \quad \text{and} \quad \log_2\left(\frac{\phi_{\text{LMDER}} - \phi^*}{\phi_{\text{DFOLS}} - \phi^*}\right),$$

where  $\phi^*$  is obtained from [86],  $\phi_{\text{DFOLS}}$ ,  $\phi_{\text{LMDER}}$  denote the best function values achieved by the respective methods, and evals<sub>DFOLS</sub>, evals<sub>LMDER</sub> denote the number of function evaluations needed to satisfy the termination test (2.2.7) for various values of  $\tau$ . If either of the solvers reach the maximum number of function evaluations before finding a solution that satisfies (2.2.7), we will set the ratio to a very large number (or negated). If both solvers fail, then the log-ratio is zero. The results comparing DFO-LS and LMDER are summarized in Figures 2.4 and 2.5. The complete table of results is given in the technical report [94].

The performance of the two methods appears to be comparable since the area of the shaded regions is similar. This may be surprising since the Gauss-Newton approach seems to be a particularly effective way of designing an interpolation-based trust-region method, requiring only linear interpolation to yield useful second-order information. Cartis



Figure 2.4. Accuracy, Noiseless Case. Log-ratio optimality gap profiles comparing DFO-LS and LMDER for  $\epsilon(x) = 0$ .



Figure 2.5. Efficiency, Noiseless Case. Log-ratio profiles comparing DFO-LS and LMDER for  $\epsilon(x) = 0$ . The figures measure number of function evaluations to satisfy (2.2.7) for  $\tau = 10^{-1}$  (left),  $10^{-3}$  (middle), and  $10^{-6}$  (right).

and Roberts [86] dismiss finite-difference methods at the outset, and do not provide numerical comparisons with them. However, the tradeoffs of the two methods merit careful consideration. DFO-LS requires only one function evaluation per iteration, but its gradient approximation,  $\hat{J}(x_k)^T r(x_k)$ , is inaccurate until the iterates approach to the solution and the trust region has shrunk. In contrast, the finite-difference Levenberg-Marquardt method in LMDER computes quite accurate gradients in this noiseless setting, requiring a much smaller number of iterations, but at a much higher cost per iteration in terms of function evaluations. The tradeoffs of the two methods appear to yield, in the end, similar performance, but we should note that DFO-LS is typically more efficient in the early stages of the optimization, as illustrated in Figure 2.5 for the low tolerance level  $\tau = 10^{-1}$ . On the other hand, the finite-difference approach is more amenable to parallel execution, as mentioned in the previous section.

Let us now consider the linear algebra cost of the two methods. A typical iteration of DFO-LS solves the interpolation system with LU factorization, for n different right hand sides, at a per-iteration cost of  $O(mn^2)$  flops, assuming that m > n. (DFO-LS offers a number of options, such as regression, which may involve a higher cost, but we did not invoke those options.) The linear algebra cost of the finite-difference version of LMDER described above is O(mn) flops. Therefore, in terms of CPU time, LMDER is faster than DFO-LS whenever the cost of function evaluations does not dominate the iteration cost.

## 2.3.2. Experiments on Noisy Functions

Let us assume that the noise model is the same across all residual functions, i.e.,  $\epsilon_i(x)$  are i.i.d. for all *i*. As in Section 2.2.2, we generate noise independently of *x*, following a uniform distribution,  $\epsilon_i(x) \sim \sigma_f U(-\sqrt{3}, \sqrt{3})$ , with noise levels  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ .

Differencing will be performed more precisely than in the noiseless case. Following [70], the forward-difference approximation of the Jacobian is defined as

(2.3.3) 
$$[\hat{J}(x_k)]_{ij} = \frac{r_i(x+h_{ij}e_j) - r_i(x)}{h_{ij}}, \text{ where } h_{ij} = 8^{1/4} \left(\frac{\sigma_f}{L_{ij}}\right)^{1/2}$$

where  $L_{ij}$  is a bound on  $|e_j^T \nabla^2 \gamma_i(x) e_j|$  within the interval  $[x, x + h_{ij} e_j]$ . To estimate  $L_{ij}$  for every pair (i, j), and at every iteration, would be impractical, and normally unnecessary. Several strategies can be designed to provide useful information at an acceptable cost. For concreteness, we estimate  $L_{ij}$  once at the beginning of the run of LMDER. More concretely, we compute a different  $h_{ij}$  for every residual function  $\gamma_i$  and each coordinate direction  $e_j$  at the starting point, and keep  $h_{ij}$  constant throughout the run of LMDER. The  $\{L_{ij}\}$  for i = 1, ..., m, j = 1, ..., n are obtained by applying the Moré-Wild (MW) procedure [71] described in the previous section to estimate the Lipschitz constant for function  $\gamma_i$  along coordinate directions  $e_j$ . If the MW procedure fails, we set  $L_{ij} = 1$ . The cost of computing the  $L_{ij}$ , in terms of function evaluations, is accounted for in the results reported below.

In DFO-LS, we did not employ restarts, and set the obj\_has\_noise option to its default value False, which also changes some trust-region-related parameters. We did so for two reasons. First, restarts introduce randomness, and as Cartis et al. [18] observed, can lead the algorithm to a different minimizer, making comparisons difficult. In addition, restarts are designed to allow the algorithm to make further progress as it reaches the noise level of the function.

Accuracy. We compare the best optimality gap achieved by LMDER and DFO-LS. We run the algorithms using their default parameters (except that we set  $model.abs_tol = 0$  for DFO-LS) until no further progress could be made. In Figure 2.6, we plot the log-ratio profiles

(2.3.4) 
$$\log_2\left(\frac{\phi_{\text{LMDER}} - \phi^*}{\phi_{\text{DFOLS}} - \phi^*}\right),$$

for noise levels  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}.$ 

We observe that DFO-LS is more accurate than LMDER, which points to the strengths of DFO-LS, since it does not require knowledge of the noise level of the function in its



Figure 2.6. Accuracy, Noisy Case. Log-ratio optimality gap profiles comparing DFO-LS and LMDER for  $\sigma_f = 10^{-1}$  (upper left),  $10^{-3}$  (upper right),  $10^{-5}$ (bottom left),  $10^{-7}$  (bottom right). The bounds on the second derivatives are kept constant over the optimization process.

internal logic. However, our implementation of LMDER is not sophisticated, as fixing the Lipschitz constant at the start of the finite-difference method is not always a good strategy. An interesting open question is whether the adaptive strategies recently proposed in [92] would close the accuracy gap.

Efficiency. To measure the efficiency of the algorithms in the noisy case, we record the number of function evaluations required to satisfy the termination condition (2.2.14), where  $\tilde{\phi}^*$  denotes the best objective value achieved by the two solvers, for a given noise level. To do so, both solvers were run until they could not make more progress. The differencing interval in LMDER was computed as in the experiments measuring accuracy for the noisy setting, i.e., by employing only the Moré-Wild procedure at the first iteration.

In Figure 2.7, we plot

(2.3.5) 
$$\log_2\left(\frac{\text{evals}_{\text{LMDER}}}{\text{evals}_{\text{DFOLS}}}\right)$$

for two values of the tolerance parameter  $\tau$  and for two levels of noise. We omit the plots for other values of  $\tau$  and  $\sigma$ , as they demonstrate similar performance. (When (2.2.14) cannot be satisfied for a solver within the given budget of  $500 \times n$  function evaluations, we set the corresponding value in (2.3.5) to a very large number.)



Figure 2.7. Efficiency, Noisy Case. Log-ratio profiles comparing DFO-LS and LMDER for  $\sigma_f = 10^{-1}$  (top row), and  $10^{-3}$  (bottom row). The figure measures the number of function evaluations to satisfy (2.2.14) for  $\tau = 10^{-1}$  (left column), and  $10^{-6}$  (right column). Lipschitz constants were estimated only at the start of the LMDER run.

There is no clear winner among the two codes used in the experiments reported in Figure 2.7. For low accuracy ( $\tau = 0.1$ ), LMDER appears to be more efficient, whereas the opposite is true for high accuracy ( $\tau = 10^{-6}$ ). We note again that DFO-LS is able to
handle different noise levels efficiently and reliably, without knowledge of the noise level or Lipschitz constants. On the other hand, the finite-difference approach is competitive even with a fairly coarse Lipschitz estimation procedure, and perhaps more important, it can be incorporated into existing codes (doing so in LMDER required little effort). In other words, in the finite-difference approach to derivative-free optimization, algorithms do not need to be constructed from scratch but can be built as adaptations of existing codes.

**2.3.2.1.** Commentary. Our finite-difference approach is tailored to each individual residual component and each coordinate direction. However, it is natural to question the necessity of estimating the Lipschitz constant for each component of each individual residual, as we did in our experiments, since this is affordable only if one can evaluate residual functions individually. One can envision problems for which a much simpler Lipschitz estimation suffices. For example, in data-fitting applications all individual residual functions  $\gamma_i$  may be similar in nature. In this setting, one could use a single Lipschitz constant, say L, across all components and residual functions, especially when the variables are scaled prior to optimization. L could be updated a few times in the course of the optimization process. On the other hand, if the scale of the variables varies significantly, one can compute Lipschitz constants  $L_i$  for each component across all residual functions, requiring the estimation of n Lipschitz constants.

## 2.4. Constrained Optimization

We now consider inequality-constrained nonlinear optimization problems of the form

(2.4.1) 
$$\min_{x \in \mathbb{R}^n} \phi(x) \qquad \text{s.t.} \quad \psi(x) \le 0, \quad l \le x \le u,$$

where  $\psi : \mathbb{R}^n \to \mathbb{R}^m$  represents a set of *m* linear or nonlinear constraints,  $l, u \in \mathbb{R}^n$ , and  $\phi$ and  $\psi$  are twice continuously differentiable. We assume that the derivatives of  $\phi$  and  $\psi$  are not available, and more generally that we have access only to noisy function evaluations:

(2.4.2) 
$$f(x) = \phi(x) + \epsilon(x), \qquad c_j(x) = \psi_j(x) + \epsilon_j(x),$$

We assume the same noise model for the objective and each of the constraint functions, for simplicity. We do not present comparisons for general problems involving both equality and inequality constraints because we were not able to find an established DFO code of such generality that was sufficiently robust in our experiments (COBYLA accepts only inequality constraints).

To our knowledge, there have been very few comparative studies of DFO methods for constrained optimization as judged by the very few references in the comprehensive review by Larson et al. [59]. The best known interpolation-based DFO software available for solving problem (2.4.1) is COBYLA, developed by Powell [77]. The method implemented in that code constructs linear approximations to  $\phi$  and  $\psi$  at every iteration using function interpolation at points placed on a simplex in  $\mathbb{R}^n$ ; it is designed to handle only inequality constraints. A more recent method by Powell, in the spirit of NEWUOA, is LINCOA [76]. It implements an interpolation-based trust-region approach but it can handle only linear constraints.

There are many production-quality software packages for deterministic constrained optimization where we could implement the finite difference DFO approach. We chose KNITRO because one of the algorithms it offers is a simple sequential quadratic programming (SQP) method that is close in spirit to COBYLA. We did not to employ the interior-point methods offered by KNITRO, which are known to be very powerful techniques for handling inequality constraints, because they may put COBYLA at an algorithmic disadvantage. In the same vein, we did not employ SNOPT because it implements a sophisticated SQP method with many advanced features to improve efficiency and reliability. In short, we selected a simple nonlinear optimization method to more easily identify the strengths and weaknesses of the finite difference approach.

The two codes are tested under the following settings.

- COBYLA. We ran the version of COBYLA maintained in the PDFO package [84]. We set the final trust region radius to  $10^{-8}$  (rhoend=1e-8) to observe its asymptotic behavior, particularly in the noiseless case. We ran PDFO version 1.0, and called COBYLA via its Python interface (Python 3.7.7).
- KNITRO. We ran Artelys Knitro 12.2 with alg=4 (an SQP algorithm), gradopt=2 (forward differencing), and hessopt=6 (L-BFGS). The choice of the finite difference interval h is described below. In order to make the algorithm as close as possible to COBYLA, we set the memory size of L-BFGS updating to its minimum value, t = 1 (lmsize=1). For consistency with COBYLA, we disabled the termination test based on the optimality error by setting opttol=1e-16, and findiff\_terminate=0, and instead terminate when the computed step is less than 10<sup>-8</sup> (by setting xtol=1e-8 and xtol\_iters=1). We called KNITRO via its Python interface.

As in the previous sections, we used test problems from the CUTEst set [42], which were called through the Python interface, PyCUTEst version 1.0. We recall from (2.4.1) that n and m refer to the number of variables and constraints, respectively (excluding bound constraints). We first selected fixed-size problems that have at least one general nonlinear inequality and have no equality constraints, and for which  $n \leq 100$  and  $m \leq 100$ . The characteristics of the resulting 49 problems are listed in the technical report [94].

## 2.4.1. Experiments on Noiseless Functions

In the first set of experiments, we applied COBYLA and KNITRO to solve the 49 small-scale, fixed-size CUTEst problems with exact function evaluations, i.e., with  $\epsilon(x) = \epsilon_j(x) \equiv 0$  for all *i*. As in prior sections, the approximate gradient  $g(x) \in \mathbb{R}^n$  of the objective function and Jacobian  $\hat{J}(x) \in \mathbb{R}^{m \times n}$  of the constraints are evaluated by simple forward differences, where  $h_i$  is set as in (2.2.3). Both algorithms were stopped when the number of function evaluations exceed 500 max (n, m), or when the trust-region radius or steplength reaches its lower bound for COBYLA or KNITRO, respectively. Both solvers report feasibility error as the max-norm of constraint violations.

We sorted the results into four groups, according to the following outcomes:

(i) Both solvers converged to a feasible point with approximately the same objective function values.

- (ii) The solvers converged to feasible points with different objective function values.
- (iii) One of the solvers terminated at an infeasible point.
- (iv) Both solvers terminated at infeasible points.

There were 32 problems associated with outcome (i). We can use them to safely compare the performance of the two solvers in terms of accuracy and efficiency. (We comment on outcomes (ii)-(iv) in the companion technical report [94].) Given that the two codes achieved feasibility in these 32 problems, we measure accuracy by comparing the best objective function value obtained by each solver with the value  $\phi^*$ obtained by running KNITRO with exact gradients until it could not make further progress. (In all the runs for determining  $\phi^*$ , feasibility error was less than  $10^{-10}$ .) To measure accuracy, we report the ratios

(2.4.3) 
$$\log_2\left(\frac{\max\{\phi_{\text{KNITRO}} - \phi^*, 10^{-8}\}}{\max\{\phi_{\text{COBYLA}} - \phi^*, 10^{-8}\}}\right).$$

We compare accuracy up to eight digits because reporting, say, the ratio  $10^{-12}/10^{-15}$  would be misleading, given that the accuracy in the constraint violation could have the inverse ratio. The results are presented in Figure 2.8, which shows that for most problems both solvers were able to achieve eight digits of accuracy; for the remaining problems, KNITRO gave higher accuracy.



Figure 2.8. Accuracy, Noiseless Case. Log-ratio profiles comparing KNITRO and COBYLA for  $\epsilon(x) = \epsilon_i(x) = 0$ . The figure plots the ratios (2.4.3) for problems for which the two solvers yielded the same solution.

To measure efficiency, we conducted a series of experiments using the same subset of 32 problems corresponding to outcome (i). We record the number of function evaluations,  $evals_{KNITRO}$  and  $evals_{COBYLA}$ , required by the two codes to satisfy condition (2.2.7) for various values of  $\tau$ . (If a solver fails to satisfy this test, we set the number of evaluations



Figure 2.9. Efficiency, Noiseless Case. Log-ratio profiles comparing KNITRO and COBYLA for  $\epsilon(x) = \epsilon_i(x) = 0$ . The figures measure number of function evaluations to satisfy (2.2.7) for  $\tau = 10^{-1}$  (left),  $10^{-3}$  (middle), and  $10^{-6}$  (right).

to a large value.) Figure 2.9 plots the ratios

(2.4.4) 
$$\log_2\left(\frac{\text{evals}_{\text{KNITRO}}}{\text{evals}_{\text{COBYLA}}}\right)$$

for  $\tau \in \{10^{-1}, 10^{-3}, 10^{-6}\}$ . The technical report [94] contains the complete set of results. We observe that for low accuracy, COBYLA is slightly more efficient, while KNITRO becomes significantly more efficient when high accuracy is required.

We should note that employing memory size t = 1 in L-BFGS updating yields a very weak quadratic model in the SQP method of KNITRO. We experimented with a memory of size t = 10 and observed that the performance of KNITRO improved, particularly in the early iterations of the runs, but not dramatically, which was surprising.

We also conducted experiments with CUTEst test problems of variable dimension and noted that the advantage of the FD approach becomes greater as the dimension of the problem increases; see [94]. This stands in stark contrast with a common perception in the DFO literature that finite differences require too many function evaluations compared to methods specifically designed for DFO problems.

## 2.4.2. Experiments on Noisy Functions

We now inject artificial noise in the evaluation of the objective and constraint functions, using the same noise model as in the unconstrained setting. We employ forward differences to approximate the gradient and Jacobian, with finite-difference interval

$$h_i = \max\{1, |[x]_i|\} \sqrt{\sigma_f}, \quad i = 1, ..., n.$$

We do not include Lipschitz constant estimates because this simple formula suffices for our purposes.

In the first experiment, we ran the two codes with their least stringent termination tests to observe the quality of the final solutions. As in the noiseless case, we set **rhoend=1e-8** for COBYLA and **xtol=1e-8** for KNITRO, and impose a limit of 500 max{n, m} function evaluations. We tested the 32 CUTEst problems for which both solvers converged to the same feasible solution in the noiseless case (outcome (i) above). In Figure 2.10, we compare the accuracy,  $\phi(x_k) - \phi^*$ , in the true objective given by the code codes, up to eight digits, as in (2.4.3). If the feasibility violation is large compared to the noise level, i.e.  $\|\max\{\psi(x_k), 0\}\|_{\infty} \ge \sqrt{3}\sqrt{\sigma_f}$ , we mark the corresponding run as a failure. We observe from Figure 2.10 that the performance of the two solvers is comparable, with KNITRO slightly more efficient for low accuracy. Interestingly, for the highest accuracy,  $\sigma_f = 10^{-7}$ , the performance of the codes is remarkably close (note that the log ratio is nearly zero for about half of the problems).



Figure 2.10. Accuracy, Noisy Case. Log-ratio profiles comparing accuracy in the objective by KNITRO and COBYLA, for  $\sigma_f = 10^{-1}$  (upper left),  $10^{-3}$ (upper right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).

Next, we compare the efficiency of the two solvers for two levels of accuracy in the objective. Specifically, we record the number of function evaluations required to satisfy

(2.4.5) 
$$\phi(x_k) - \tilde{\phi}_* \le \tau(\phi(x_0) - \tilde{\phi}_*) \quad \text{and} \quad \|\max\{\psi(x_k), 0\}\|_{\infty} \le \sqrt{3\sqrt{\sigma_f}},$$

for  $\tau \in \{10^{-2}, 10^{-6}\}$ , where  $\tilde{\phi}_*$  is the minimum objective value obtained by the two codes. Figure 2.11 plots the ratios (2.4.4), and suggests that it is difficult to choose between the two codes. Therefore, in our tests on noisy problems, an unsophisticated code (KNITRO/SQP) for deterministic nonlinear optimization, using a simple strategy for choosing the finite difference interval, is competitive with COBYLA, a method specifically designed for derivative-free optimization.



Figure 2.11. Efficiency, Noisy Case. Log-ratio profiles (2.4.4) comparing KNITRO and COBYLA for  $\sigma_f = 10^{-3}$  (top row) and  $10^{-7}$  (bottom row). The figure measures the number of function evaluations to satisfy (2.4.5) for  $\tau = 10^{-2}$  (left) and  $10^{-6}$  (right).

# 2.5. Final Remarks

Finite-difference approximations are widely employed in numerical analysis, particularly for solving differential equations. The limitations of finite-difference methods are welldocumented, particularly in the noisy case. However, optimization provides a more benign setting as errors do not necessarily accumulate; if a poor gradient estimation yields a bad step, it may be corrected at a later step. Two attractive features of finite-difference methods for derivative-free optimization are the simplicity of building them upon existing nonlinear (gradient-based) optimization solvers, and their ease of parallelization.

Our numerical study indicates that finite-difference-based optimization methods can be made competitive, in most cases, against state-of-the-art methods designed specifically for DFO. Results to that effect were reported by Berahas et al. [7] in the context of unconstrained optimization. Our study tests a simpler version of L-BFGS than that explored in [7], and more important, considers also least squares and constrained optimization—settings for which comparative studies do not exist, to the best of our knowledge.

Our experiments show that the finite-difference approach can be further improved by employing more sophisticated (adaptive) procedures for computing the finite-difference interval. The key is in better automatic estimation of local Lipschitz constants. Investigation along this important line of inquiry should lead to more robust methods for derivative-free optimization.

# CHAPTER 3

# Adaptive Finite-Difference Interval Estimation for Noisy Derivative-Free Optimization

## 3.1. Introduction

A powerful approach for derivative-free optimization is to utilize finite differences. This is done by computing a finite-difference approximation to the gradient, and substituting the exact gradient with the approximation within a known nonlinear optimization method; see [93]. These methods operate by spending at least n + 1 function evaluations at each iteration to take a meaningful step, where n is the total number of variables. This lies in contrast to interpolation-based methods, which utilize prior function evaluations with only one new evaluation at each iteration; see [29, 59]. Therefore, in order for the finitedifference approach to be effective, one must ensure that the quality of the gradient is satisfactory and significant progress is being made at each iteration of the algorithm (as opposed to n steps of the interpolation-based approach).

Often, black-box functions to be optimized are contaminated by stochastic or computational noise. This noise could arise naturally from modeling randomness within a simulation, or as a bi-product of an adaptive computation, for example through the early termination of an iterative solver. The presence of noise has largely prevented finitedifference methods from gaining more popularity within the derivative-free optimization community, as the precise choice of the finite-difference interval becomes increasingly critical as the noise level increases.

In particular, the finite-difference interval requires knowledge of both the noise level and higher-order derivative of the function. While the former may be known *a priori* or can be estimated by sampling or computing difference tables [70], the latter quantity is not normally available to the user. In order to make finite-difference methods a viable alternative in the presence of noise, a robust procedure is needed for estimating higher-order derivatives, either implicitly or explicitly.

To put this more precisely, let us consider the problem of estimating the *d*-th order derivative of a smooth univariate function  $\phi : \mathbb{R} \to \mathbb{R}$ . We will assume that we are only provided noisy function evaluations of the form

(3.1.1) 
$$f(t) = \phi(t) + \epsilon(t)$$

where  $\epsilon : \mathbb{R} \to \mathbb{R}$  models the error, and that the error is bounded, i.e., there exists  $\epsilon_f \ge 0$ such that  $|\epsilon(t)| \le \epsilon_f$ . We call  $\epsilon_f$  the *noise level* of the function. We focus on the univariate case, although this can be easily extended to the multivariate setting for computing the gradient by applying the procedure to each component.

The simplest and cheapest finite-difference approximation to the first derivative is the forward-difference approximation. If  $\phi^{(d)}$  denotes the *d*-th order derivative of  $\phi$ , then the forward-difference approximation is computed by

(3.1.2) 
$$\phi^{(1)}(t) \approx \frac{f(t+h) - f(t)}{h} \triangleq f^{(1)}(t;h)$$

where h > 0 is the finite-difference interval. Note the slight abuse of notation by denoting  $f^{(d)}$  as the finite-difference approximation to the *d*-th order derivative. With no noise, excluding round-off error, one would ideally choose *h* as small as possible, the common practical choice being  $h = \max\{1, |x|\}\sqrt{\epsilon_M}$ , where  $\epsilon_M$  is machine precision, to handle rounding errors. However, this choice of the finite-difference interval may be poor under the presence of large errors, as is well-known.

To see this, consider the following decomposition of the error in the forward-difference approximation:

(3.1.3) 
$$\left| f^{(1)}(t;h) - \phi^{(1)}(t) \right| \le \left| \frac{\phi(t+h) - \phi(t)}{h} - \phi^{(1)}(t) \right| + \left| \frac{\epsilon(t+h) - \epsilon(t)}{h} \right|$$

We will call the error induced by the first term *truncation error* since it arises from truncation of the Taylor series, and the error induced by the second term *measurement error* due to error in the function evaluations.

Note that if h is small, then the truncation error is small but the measurement error may be large. On the other hand, if h is too large, the measurement error may be small but the truncation error may be too high. Therefore, the optimal h trades off these two terms by making the error from each of these two sources equal. In this paper, we propose an adaptive procedure for estimating the finite-difference interval in the presence of noise that properly balances these two different sources of error. The procedure must be: (1) reliable, that is, applicable to most, if not all, practical problems of interest; (2) accurate, producing near-optimal estimates of the finite-difference interval; and (3) efficient, employing the least number of function evaluations possible. We argue that our procedure achieves these goals in many practical situations. This paper is organized into five sections. We present the notation and literature review in the rest of this section. In Section 2, we introduce our finite-difference interval estimation procedure for the forward-difference case. In Section 3, we present the generalized procedure for arbitrary finite-difference schemes and provide theoretical guarantees for the termination of our procedure. Extensive numerical results on synthetic problems with injected noise are provided in Section 4, and concluding remarks are made in Section 5.

### 3.1.1. Literature Review

The problem of estimating derivatives, particularly in the presence of rounding errors, is a fundamental question within numerical analysis and scientific computing. Fornberg proposed a stable algorithm for generating finite-difference formulas on arbitrarily spaced grids [38]. Lyness and Moler observed that the Cauchy integral theorem allows one to evaluate the *d*-th derivative of a complex function as a closed complex integral via numerical integration techniques [66]. This was simplified and extended by Squire and Trapp who observed that one could avoid cancellation error by using complex perturbations in the Taylor expansion, called complex step differentiation [96]. This has more recently led to extensions of the complex step to evaluating the Hessian by Hare and Srivastava [50]. Brekelmans, et al. compared design of experiments schemes against standard finite-difference schemes within the stochastic noise regime [14].

To handle rounding errors, Curtis and Reid describe a heuristic that estimates the truncation and rounding errors using central and forward-difference estimates. The ratio between the two estimates of these errors are used to determine the finite-difference interval [30]. Stepleman and Winarsky use a set of decreasing central-difference intervals. The optimal interval is obtained by the smallest interval that does not violate monotonic decrease in the absolute difference between consecutive central-difference estimates [97]. Gill, Murray, Saunders, and Wright introduced an adaptive procedure for computing forward-difference intervals by utilizing a ratio to determine the second derivative [40, 41]. Their procedure has some similarities with our approach, which we discuss in Section 3.2.1. Barton proposed an adaptive procedure for handling rounding or multiplicative errors by interpreting the function values as correct up to a fixed number of significant digits and ensuring that at least a certain number of significant digits change from the resulting difference interval [6]. Most recently, Moré and Wild proposed a heuristic for estimating the second derivative for determining the forward-difference interval that checks two conditions: (1) if the noise dominates the second-order derivative; and (2) if the forward and backward difference is too large relative to the function values [71]. A comparison of the resulting errors between finite-difference and simplex gradients were analyzed in [9, 49].

Incorporating finite differences into optimization methods have also had a long history. Kiefer and Wolfowitz first applied finite differences to stochastic approximation [56]. Kelley developed an implicit filtering BFGS method that utilizes finite differences in the case where noise decays as the iterates converge to the solution [21, 55]. Berahas, et al. proposed a finite-difference L-BFGS method that incorporates ECNoise and a heuristic for estimating the second derivative by Moré and Wild into L-BFGS [7, 70, 71]. Most recently, Shi, et al. tested finite-difference methods within the unconstrained, least squares, and constrained settings assuming knowledge of the noise level [93].

# 3.1.2. Notation

In the following sections, we will use Bachmann-Landau notation liberally. Suppose  $f,g : \mathbb{R} \to \mathbb{R}^{\geq 0}$ . We will write  $g(h) = \mathcal{O}(f(h))$  if there exists a  $C \in \mathbb{R}$  such that g(h) = Cf(h). If g(h) = o(f(h)), then for every  $\epsilon > 0$  there exists a constant N such that  $|g(h)| \leq \epsilon |f(h)|$  for all  $h \leq N$ . Similarly, if g(h) = O(f(h)), then there exists constants  $\epsilon > 0$  and N such that  $|g(h)| \leq \epsilon |f(h)|$  for all  $h \leq N$ .

We will use  $\phi^{(d)} : \mathbb{R} \to \mathbb{R}$  to denote the *d*-th order derivative of  $\phi$ . For a given vector  $x \in \mathbb{R}^n$ ,  $[x]_i$  denotes the *i*-th component of *x*. Similarly, for a given matrix  $A \in \mathbb{R}^n$ ,  $[A]_{ij}$  denotes the (i, j)-th entry of *A*. We will use  $\|\cdot\|$  to denote the standard Euclidean norm unless otherwise specified.

#### 3.2. An Adaptive Forward-Difference Interval Estimation Procedure

Suppose we are interested in determining the finite-difference interval for the forwarddifference approximation of the first derivative of  $\phi$ . Since the Taylor expansion of the function  $\phi$  is given by

$$\phi(t+h) = \phi(t) + \phi^{(1)}(t)h + \frac{\phi^{(2)}(t)}{2}h^2 + o(h^2),$$

the total error can be bounded by

$$|\phi^{(1)}(t) - f^{(1)}(t;h)| \le \underbrace{\frac{|\phi^{(2)}(t)|h}{2}}_{T_1} + \underbrace{\frac{2\epsilon_f}{h}}_{T_2} + o(h).$$

By ignoring the higher-order term, this yields an optimal interval (with respect to the upper bound) of

(3.2.1) 
$$h^* \approx 2\sqrt{\frac{\epsilon_f}{|\phi^{(2)}(t)|}}.$$

This formula requires an estimate of the second derivative  $|\phi^{(2)}(t)|$ . We now propose a procedure that yields an interval  $h = \mathcal{O}\left(\sqrt{\frac{\epsilon_f}{|\phi^{(2)}(t)|}}\right)$  without estimating  $|\phi^{(2)}(t)|$  separately.

Our procedure balances the truncation  $T_1$  and measurement error  $T_2$ . To do so, it estimates the ratio between these two errors directly and attempts to find an interval hfor which this ratio is close to some constant value. We claim that the ratio  $T_1/T_2$  of the truncation over measurement error can be approximated, for example, by the *testing ratio* 

(3.2.2) 
$$r(h; f, t, \epsilon_f) = \frac{|f(t+4h) - 4f(t+h) + 3f(t)|}{8\epsilon_f}$$

Given  $r_l > 1$  and  $r_u > r_l + 2$ , we perform a bisection search to find an interval h > 0 that satisfies

$$(3.2.3) r(h; f, t, \epsilon_f) \in [r_l, r_u].$$

In particular, if  $r(h; f, t, \epsilon_f) < r_l$ , then the numerator is dominated by noise, indicating that h is too small. On the other hand, if  $r(h; f, t, \epsilon_f) > r_u$ , then the numerator significantly dominates the noise, which implies that h is too large. Our procedure for forward differences is summarized in Algorithm 3.1.

1: Input: One-dimensional noisy function  $f : \mathbb{R} \to \mathbb{R}$ ; noise level  $\epsilon_f > 0$ ; lower- and upper-bound  $(r_l, r_u) = (1.5, 6);$ 2: **Output:** Finite-difference interval h such that (3.2.3) holds. 3:  $h \leftarrow \frac{2}{\sqrt{3}}\sqrt{\epsilon_f};$ 4:  $l \leftarrow 0, u \leftarrow +\infty;$ 5: while True do Evaluate  $r(h; f, t, \epsilon_f) = \frac{|f(t+4h)-4f(t+h)+3f(t)|}{8\epsilon_f};$ 6: if  $r(h; f, t, \epsilon_f) < r_l$  then 7: $l \leftarrow h$ : 8: else if  $r(h; f, t, \epsilon_f) > r_u$  then 9:  $u \leftarrow h;$ 10: else 11: 12:break; end if 13:if  $u = +\infty$  then 14: 15: $h \leftarrow 4h;$ else if l = 0 then 16: $h \leftarrow h/4;$ 17:else 18:  $h \leftarrow (l+u)/2;$ 19:end if 20: 21: end while 22: return h

To see why this procedure works to give us a near-optimal h, note that

(3.2.4) 
$$\phi(t+4h) - 4\phi(t+h) + 3\phi(t) = 6\phi^{(2)}(t)h^2 + o(h^2).$$

Therefore, if we expand (3.2.2), we obtain:

(3.2.5) 
$$r(h; f, t, \epsilon_f) = \left| \frac{3\phi^{(2)}(t)h^2}{4\epsilon_f} + \frac{\epsilon(t+4h) - 4\epsilon(t+h) + 3\epsilon(t)}{8\epsilon_f} + o(h^2) \right|.$$

Since  $\left|\frac{\epsilon(t+4h)-4\epsilon(t+h)+3\epsilon(t)|}{8\epsilon_f}\right| \leq 1$  by the fact that  $|\epsilon(t)| \leq \epsilon_f$  for all  $t \in \mathbb{R}$ , by imposing (3.2.3) and ignoring the  $o(h^2)$  term, we approximately have

(3.2.6) 
$$\frac{3|\phi^{(2)}(t)|h^2}{4\epsilon_f} \in [r_l - 1, r_u + 1] \iff h \in \frac{2}{\sqrt{3}}\sqrt{\frac{\epsilon_f}{|\phi^{(2)}(t)|}} \cdot [\sqrt{r_l - 1}, \sqrt{r_u + 1}].$$

Therefore, if  $r_l$  and  $r_u$  are chosen properly, such as  $r_l = 1.5$  and  $r_u = 6$ , we obtain  $h \in [\sqrt{0.5}, \sqrt{7}] \cdot \sqrt{\frac{\epsilon_f}{|\phi^{(2)}(t)|}}$ , which is the same order as the optimal finite-difference interval (3.2.1), differing only by a small constant factor.

By scaling h by a factor of 4 in Algorithm 3.1 when  $u = \infty$  or l = 0, the new trial h only requires a single new function evaluation to check the testing ratio when h is updated monotonically.

In addition, the testing ratio is affine-invariant with respect to the function of interest in the sense that  $r(h; f, t, \epsilon_f)$  remains unchanged if applied to a modified function  $\tilde{f}(t) = af(t) + b$  for  $a \neq 0$  and  $b \in \mathbb{R}$  with noise level  $|a|\epsilon_f$ , i.e.  $r(h; \tilde{f}, t, |a|\epsilon_f) = r(h; f, t, \epsilon_f)$ . Therefore, the finite-difference interval h will correctly remain unchanged under this transformation.

### 3.2.1. Comparison to Prior Methods

Although the definition of the testing ratio appears similar to the ratio in Gill, et al. [40], which is defined as the inverse ratio

(3.2.7) 
$$\frac{4\epsilon_f}{|f(t+\tilde{h}) - 2f(t) + f(t-\tilde{h})|},$$

our approach markedly differs from theirs in three aspects: (1) we utilize the h derived from Algorithm 3.1 directly as the chosen difference interval, whereas Gill, et al. use the difference interval  $\tilde{h}$  to estimate the second derivative; (2) our approach utilizes a bisection search to find h rather than monotonically increasing or decreasing h; and (3) it relies on second-order forward-difference instead of central-difference estimates.

The first aspect follows from the observation made in (3.2.6) that  $h = \mathcal{O}\left(\sqrt{\frac{\epsilon_f}{|\phi^{(2)}(t)|}}\right)$ . Hence, the difference interval h obtained by our bisection procedure is near-optimal in that it is only a small constant factor away from the optimal interval (except for certain cases discussed below). Two observations can be made from this derivation. The first is that since f(t+h) and f(t) are used in the evaluation of the testing ratio, one can reuse prior function evaluations from the bisection procedure to estimate  $f^{(1)}(t;h)$ . The second observation is that the derivation motivates an initial choice of  $h = \mathcal{O}(\sqrt{\epsilon_f})$  as opposed to  $h = \mathcal{O}(\sqrt{\epsilon_f})$ , as used in Moré and Wild [71]. This is corroborated by our experiments in Section 3.4.

The second aspect follows from different tradeoffs between cost and accuracy in the finite-difference interval estimate. In particular, whereas Gill, et al.'s procedure is cheaper but yields a less accurate estimate, our procedure is able to guarantee a sufficiently accurate estimate at higher cost.

Lastly, the third aspect is designed to avoid cancellation due to symmetry in the numerator of the testing ratio. In particular, Gill, et al. [40] rely on a testing ratio using the central difference:

(3.2.8) 
$$\frac{4\epsilon_f}{|f(t+\tilde{h}) - 2f(t) + f(t-\tilde{h})|} \in [0.001, 0.1].$$

However, we can obtain poor estimates of the derivative due to cancellation under symmetry of the function, for example, on a quartic function  $\phi(t) = t^4$  near (but not at) t = 0 with sufficiently large noise.

Our procedure also differs from Moré and Wild's procedure [71]. Their procedure estimates the second derivative by

(3.2.9) 
$$\phi^{(2)}(t) \approx \frac{f(t+\tilde{h}) - 2f(t) + f(t-\tilde{h})}{\tilde{h}^2} = f^{(2)}(t;\tilde{h}),$$

with interval  $\tilde{h} > 0$ , then inserts this estimate into the optimal formula (3.2.6). The difference interval  $\tilde{h}$  is required to satisfy

$$(3.2.10) \qquad |f(t+\tilde{h}) - 2f(t) + f(t-\tilde{h})| \ge \tau_1 \epsilon_f$$

(3.2.11) 
$$|f(t \pm \tilde{h}) - f(t)| \le \tau_2 \max\{|f(t)|, |f(t \pm \tilde{h})|\}$$

with  $\tau_1 \gg 1$  and  $\tau_2 \in (0, 1)$ . Their method attempts to satisfy this within two trials as follows:

- (1) Set  $\tilde{h}_1 = \sqrt[4]{\epsilon_f}$  and compute  $\mu_1 = |f^{(2)}(t; \tilde{h}_1)|$ . If conditions (3.2.10) and (3.2.11) are satisfied for  $\tilde{h}_1$ , return  $\mu_1$ .
- (2) Set  $\tilde{h}_2 = \sqrt[4]{\epsilon_f/\mu_1}$  and compute  $\mu_1 = |f^{(2)}(t; \tilde{h}_2)|$ . If conditions (3.2.10) and (3.2.11) are satisfied for  $\tilde{h}_2$ , return  $\mu_2$ .

(3) If 
$$|\mu_1 - \mu_2| \le \frac{1}{2}\mu_2$$
, return  $\mu_2$ .

If the heuristic is unable to return an estimate  $\mu$  after two trials, this is considered as a failure.

Similar to Gill's procedure, this method differs from ours in that it estimates the second derivative rather than utilizing the h derived from the testing ratio directly. As a result, it initializes  $\tilde{h} = \mathcal{O}(\sqrt[4]{\epsilon_f})$  rather than  $\mathcal{O}(\sqrt{\epsilon_f})$ . We have found this to be a poor initial choice of  $\tilde{h}$ , particularly when  $\epsilon_f$  is large, as we will see in Section 3.4.

While (3.2.10) appears similar to the testing ratio, it is better interpreted as ensuring that noise does not dominate the second-derivative estimation due to the large choice of  $\tau_1 = 100$ . The second condition (3.2.11) is not affine-invariant in the sense that adding a sufficiently large *b* can force the condition to be satisfied. This is undesirable as perturbations of the function should not change the overall behavior of the method.

## 3.3. Generalized Finite-Difference Interval Estimation

Typically, finite-difference interval estimation procedures for numerical optimization focus on forward differences [6, 40, 71]. However, in the noisy regime, higher-order finite-difference approximations, such as central differences, can yield more accurate approximations; see [93]. As a result, in order to attain the highest possible accuracy, one must design methods that efficiently find a near-optimal difference interval for more general finite-difference schemes. To handle this, we propose a generalization of the forward-difference case, Algorithm 3.1, for d-th order derivatives.

Consider a finite-difference approximation scheme S = (w, s) defined over m points, where we approximate  $\phi^{(d)}(t)$  using the equation

(3.3.1) 
$$f_S^{(d)}(t;h) = \frac{\sum_{j=1}^m w_j \cdot f(t+hs_j)}{h^d} \approx \phi^{(d)}(t)$$

where  $w \in \mathbb{R}^m$  and  $s \in \mathbb{R}^m$  are the associated weights and shifts of the finite-difference scheme. As in the forward-difference setting, we will use a slight abuse of notation by denoting the finite-difference approximation as  $f_S^{(d)}(t;h)$ . Forward-difference and centraldifference schemes for approximating the first derivative (i.e., d = 1) are obtained by defining s and w as  $s = (0,1)^T$  and  $w = (-1,1)^T$  and  $s = (-1,1)^T$  and  $w = (-\frac{1}{2},\frac{1}{2})^T$ , respectively. The standard second-order central-difference scheme is defined as s = $(-1,0,1)^T$  and  $w = (1,-2,1)^T$ .

In order for the finite-difference scheme to be valid, the coefficients w and shifts s must be chosen such that the Taylor expansion of the finite-difference approximation over the function  $\phi$  satisfies

(3.3.2) 
$$\sum_{j=1}^{m} w_j \cdot \phi(t+hs_j) = \phi^{(d)}(t)h^d + c_q \phi^{(q)}(t)h^q + o(h^q),$$

where  $q \ge d+1$  denotes the order of the remainder term<sup>1</sup>. This ensures that  $f_S^{(d)}(t;h) \approx \phi^{(d)}(t)$ . In order to guarantee this, the finite-difference scheme S must satisfy

$$\frac{1}{l!} \sum_{j=1}^{m} w_j s_j^l = \begin{cases} 1 & \text{if } l = d \\ 0 & \text{if } l < q, \ l \neq d \end{cases}$$

and as a result

$$c_q = \frac{1}{q!} \sum_{j=1}^m w_j s_j^q.$$

See Appendix B.1 for more detail on how generic finite-difference schemes are derived.

Therefore, in the presence of noise, the worst-case error for the finite-difference scheme of interest can be bounded by

$$|f^{(d)}(t;h) - \phi^{(d)}(t)| \le |c_q| \left| \phi^{(q)}(t) \right| h^{q-d} + ||w||_1 \epsilon_f h^{-d} + o(h^{q-d}).$$

One can define an approximately optimal choice of h:

(3.3.3) 
$$h^* \approx \left| \frac{d}{q-d} \cdot \frac{\|w\|_1 \epsilon_f}{c_q \phi^{(q)}(t)} \right|^{1/q}$$

While  $\epsilon_f$  is assumed to be known and d, q, w and  $c_q$  are available, the q-th order derivative  $\phi^{(q)}(x)$  is unknown and often difficult to estimate. Following the idea from the forward-difference case, we propose a procedure for estimating (3.3.3) directly. We first construct a testing ratio  $r_S$  associated with scheme S:

$$r_S(h; f, t, \epsilon_f) = \frac{\left|\sum_{j=1}^{\tilde{m}} \tilde{w}_j \cdot f(t+h\tilde{s}_j)\right|}{\epsilon_f}$$

<sup>&</sup>lt;sup>1</sup>Note that the order of accuracy can be higher than d + 1 for certain schemes, such as central-difference approximations.

where  $\tilde{w}, \tilde{s} \in \mathbb{R}^{\tilde{m}}$  where  $\tilde{m} \ge q+1$ ,  $\tilde{s}_j \ne \tilde{s}_k$  for all  $j \ne k$ , and  $\tilde{w}$  and  $\tilde{s}$  satisfies

(3.3.4) 
$$\sum_{j=1}^{\tilde{m}} \tilde{w}_j \cdot \phi(t+h\tilde{s}_j) = c_r \phi^{(q)}(t)h^q + o(h^q), \ c_r = \frac{1}{q!} \sum_{j=1}^{\tilde{m}} \tilde{w}_j \tilde{s}_j^q \neq 0.$$

Without loss of generality, we require  $\tilde{w}$  to satisfy

$$\|\tilde{w}\|_1 = 1,$$

the reason for which will be evident below. We then perform a bisection search to find an interval h that satisfies

$$(3.3.5) r_S(h; f, t, \epsilon_f) \in [r_l, r_u]$$

for some  $r_l > 1$  and  $r_u > r_l + 2$ . The procedure is summarized in Algorithm 3.2.

1:  $h \leftarrow h_0$ ; 2:  $l \leftarrow 0, u \leftarrow +\infty;$ 3: while True do Evaluate  $r_S(h; f, t, \epsilon_f)$ ; 4: if  $r_S(h; f, t, \epsilon_f) < r_l$  then 5: $l \leftarrow h;$ 6: else if  $r_S(h; f, t, \epsilon_f) > r_u$  then 7:  $u \leftarrow h;$ 8: 9: elsebreak; 10: end if 11: if  $u = +\infty$  then 12: $h \leftarrow \eta h;$ 13: else if l = 0 then 14: $h \leftarrow h/\eta;$ 15:16:else  $h \leftarrow (l+u)/2;$ 17:18: end if 19: end while 20: return h

By (3.3.4), we can see that

(3.3.6) 
$$r_S(h; f, t, \epsilon_f) = \left| \frac{c_r \phi^{(q)}(t) h^q}{\epsilon_f} + \frac{\sum_{j=1}^{\tilde{m}} \tilde{w}_j \cdot \epsilon(t+h\tilde{s}_j)}{\epsilon_f} + o\left(h^q\right) \right|$$

(3.3.7) 
$$= \left| \frac{c_r \phi^{(q)}(t) h^q}{\epsilon_f} + \Delta + o(h^q) \right|, \ \Delta \triangleq \frac{\sum_{j=1}^{\tilde{m}} \tilde{w}_j \cdot \epsilon(t + h\tilde{s}_j)}{\epsilon_f}.$$

Note that by definition of  $\Delta$ ,  $|\Delta| \leq 1$ . This is a consequence of the requirement that  $\|\tilde{w}\|_1 = 1$  and that  $|\epsilon(t)| \leq \epsilon_f$  for all  $t \in \mathbb{R}$ . Therefore, if we have  $r_S(h; f, t, \epsilon_f) \in [r_l, r_u]$  and if we the ignore  $o(h^q)$  term, then we (approximately) have

(3.3.8) 
$$\left|\frac{c_r\phi^{(q)}(t)h^q}{\epsilon_f}\right| \in [r_l - 1, r_u + 1],$$

i.e.,

(3.3.9) 
$$h \in \left[ \left( \frac{r_l - 1}{|c_r|} \frac{\epsilon_f}{|\phi^{(q)}(t)|} \right)^{1/q}, \left( \frac{r_u + 1}{|c_r|} \frac{\epsilon_f}{|\phi^{(q)}(t)|} \right)^{1/q} \right].$$

Note from (3.3.3) that h has the same dependence on  $\epsilon_f$  and  $\phi^{(q)}(t)$  as  $h^*$ . As in the forward-difference case, our algorithm is invariant to affine transformations with respect to the function.

*Example 1 (First-Order Central Difference).* Consider the first-order central-difference scheme for approximating the first derivative:

(3.3.10) 
$$f_S^{(1)}(t;h) = \frac{f(t+h) - f(t-h)}{2h},$$

where  $s = (-1, 1)^T$  and  $w = (-\frac{1}{2}, \frac{1}{2})^T$ . The Taylor expansion of the numerator is given as:

$$\frac{\phi(t+h) - \phi(t-h)}{2h} = \phi^{(1)}(t) + \frac{\phi^{(3)}(t)h^2}{6} + o(h^2).$$

The full error of the derivative approximation and the approximate optimal choice of h are:

$$\left| f_S^{(1)}(t;h) - \phi^{(1)}(t) \right| \le \frac{\left| \phi^{(3)}(t) \right| h^2}{6} + \frac{\epsilon_f}{h} + o(h^2), \ h^* \approx \sqrt[3]{\frac{3\epsilon_f}{|\phi^{(3)}(t)|}}.$$

One example of a valid testing ratio is:

(3.3.11) 
$$r_S(h; f, t, \epsilon_f) = \frac{|f(t+3h) - 3f(t+h) + 3f(t-h) - f(t-3h)|}{8\epsilon_f}$$

*Example 2 (Second-Order Central Difference).* Consider the second-order central-difference scheme for approximating the second derivative:

(3.3.12) 
$$f_S^{(2)}(t;h) = \frac{f(t+h) - 2f(t) + f(t-h)}{h^2},$$

where  $s = (-1, 0, 1)^T$  and  $w = (-\frac{1}{2}, \frac{1}{2})^T$ . The Taylor expansion of the numerator is given as:

$$\frac{\phi(t+h) - 2\phi(t) + \phi(t-h)}{h^2} = \phi^{(2)}(t) + \frac{\phi^{(4)}(t)h^2}{24} + o(h^2).$$

The full error of the derivative approximation and the approximate optimal choice of h are:

$$\left| f_S^{(2)}(t;h) - \phi^{(2)}(t) \right| \le \frac{\left| \phi^{(4)}(t) \right| h^2}{24} + \frac{4\epsilon_f}{h^2} + o(h^2), \ h^* \approx 2\sqrt[4]{\frac{6\epsilon_f}{|\phi^{(4)}(t)|}}.$$

One example of a valid testing ratio is:

(3.3.13) 
$$r_S(h; f, t, \epsilon_f) = \frac{|f(t+2h) - 4f(t+h) + 6f(t) - 4f(t-h) + f(t-2h)|}{16\epsilon_f}.$$

# 3.3.1. Practical Considerations

In order to make the procedure both efficient and robust, we discuss a number of practical considerations below.

I. Generation of Testing Ratio. Although many choices of  $r_S$  are possible for any finitedifference scheme S, it would be useful to have a method for automatically generating valid testing ratios that efficiently utilize function values. A simple yet useful way to construct  $r_S$  is through the formula

(3.3.14) 
$$r_{S}^{\alpha}(h; f_{S}, t, \epsilon_{f}) = \frac{\left| \left( f_{S}^{(d)}(t; h) - \alpha^{-d} f_{S}^{(d)}(t; \alpha h) \right) h^{d} \right|}{A \epsilon_{f}},$$

where  $\alpha \neq 1$  and A is computed by normalizing the coefficients such that  $\|\tilde{w}\|_1 = 1$  is satisfied.

This approach is guaranteed to generate a valid testing ratio  $r_S$  for any  $\alpha \neq 1$  since it cancels out the  $\phi^{(d)}(t)$  term in the Taylor expansion, leaving only the relevant higher-order term of order q of interest. In particular, since

$$\sum_{j=1}^{m} w_j \cdot \phi(t+hs_j) = \phi^{(d)}(t)h^d + c_q \phi^{(q)}(t)h^q + o(h^q)$$
$$\sum_{j=1}^{m} w_j \cdot \phi(t+\alpha hs_j) = \phi^{(d)}(t)(\alpha h)^d + c_q \phi^{(q)}(t)(\alpha h)^q + o(h^q),$$

we obtain that

$$\left( f_S^{(d)}(t;h) - \alpha^{-d} f_S^{(d)}(t;\alpha h) \right) h^d = \sum_{j=1}^m \left( w_j \cdot \phi(t+hs_j) - \alpha^{-d} w_j \cdot \phi(t+\alpha hs_j) \right)$$
$$= c_q (1 - \alpha^{q-d}) \phi^{(q)}(t) h^q + o(h^q),$$

which satisfies (3.3.4) with an effective  $c_r = c_q (1 - \alpha^{q-d})/A$  as desired.

For finite-difference schemes with equidistant points, a small modification to the bisection search allows us to reuse 1 - 2 function evaluations at each iteration, depending on the original scheme S. To do this, we can multiply or divide by the same factor  $\eta = \alpha$  when monotonically increasing or decreasing h within the bisection search, as done with  $\eta = 4$  in the forward-difference algorithm (Algorithm 3.1).

In addition, with this choice of the testing ratio, the function evaluations needed within the finite-difference scheme is implicitly evaluated within the testing ratio. We can therefore obtain the finite-difference approximation using previously computed function values at no additional cost.

II. Choice of  $r_l$  and  $r_u$ . Ideally, one should choose  $r_l$  and  $r_u$  such that they are close to the optimal ratio

$$r^* = \frac{d}{q-d} \cdot \left| \frac{c_r}{c_q} \right| \cdot \|w\|_1$$

in order to yield an h that is close to  $h^*$  in (3.3.3). However, this is not directly possible in the presence of noise, which requires that  $1 < r_l < r_u - 2$  in order to ensure finitetermination; see Section 3.3.2. We therefore select  $(r_l, r_u)$  sufficiently large such that  $1 < r_l < r_u - 2$  and, if possible, such that  $r^*$  is logarithmically centered within the interval  $[r_l, r_u]$ :

(3.3.15) 
$$r_l = \max\left\{1+\eta, \frac{r^*}{\beta}\right\}, \ r_u = \max\left\{3(1+\eta), \beta r^*\right\}$$

for some  $\eta > 0$  and  $\beta > 1$ . (In our experiments, we set  $\eta = 0.1$  and  $\beta = 2$ .)

Note that when  $r_l, r_u > r^*$ , the algorithm may overestimate  $|\phi^{(q)}(t)|$  and hence underestimate h. In order to avoid this in practice, we have found that it is preferable to choose a testing ratio such that the optimal ratio  $r^* \ge \beta(1 + \eta)$ . This could be done by choosing a different testing ratio, such as by choosing a larger  $\alpha$  in (3.3.14).

III. Initialization of  $h_0$ . Since the difference interval h that satisfies the procedure is approximately of the form (3.3.9), it is preferable to initialize  $h_0 = \mathcal{O}(\epsilon_f^{1/q})$ . Two possible choices are  $h_0 = \epsilon_f^{1/q}$  or  $\left(\frac{d}{q-d} \cdot \frac{\|w\|_1}{|c_q|} \cdot \epsilon_f\right)^{1/q}$ . The latter is based on the assumption that  $|\phi^{(q)}(t)| \approx 1$ . If instead the finite-difference interval is re-estimated within an optimization algorithm, we can initialize  $h_0$  as the difference interval h used at the prior outer iteration of the algorithm.

We observe that on rare occasions a poor initial choice of  $h_0$  can result in large error in the derivative approximation. This occurs when the initial choice of  $h_0$  is too large to capture the local behavior of the function. Reducing the initial interval  $h_0$  resolves this issue.

*IV. Handling of Special Cases.* The Taylor expansion analysis elucidates two possible failure cases for our procedure. In particular, observe that

$$r_S(h; f, t, \epsilon_f) = \left| \frac{c_r \phi^{(q)}(t) h^q}{\epsilon_f} + \Delta + o(h^q) \right|.$$

If h is large (for example, when the noise level  $\epsilon_f$  is high), the higher-order terms  $o(h^q)$  can dominate the other terms in the Taylor series expansion. This can yield poor estimates of h even if the condition  $r_S(h; f, t, \epsilon_f) \in [r_l, r_u]$  is satisfied. In practice, we have not found this to be a common issue.

The more common case is when  $\phi^{(q)}(t) \approx 0$ . In this case,  $r_S$  will be dominated by  $\Delta$ . In this case,  $r_S(h; f, t, \epsilon_f) < r_l$  for all h and h will thus monotonically increase until the maximum number of iterations is reached (which we set max\_iter to 20). This occurs, for example, with any (q - 1)-th degree polynomial. In this case, the method provides a warning but does not flag this as a failure. Note that in this case, h is a good choice because sending  $h^* \to \infty$  would allow for indefinite reduction in the noise.

V. Extension to Standard Deviation. In some settings, we only have access to the standard deviation of the noise, where  $\epsilon(x)$  is modeled as a random variable. Assuming  $\mathbb{E}[\epsilon(x)] = 0$ , one can extend this procedure to the stochastic setting by replacing  $\epsilon_f$  with  $\sigma_f = \sqrt{\mathbb{E}[\epsilon(x)^2]}$ . While finite termination (see next subsection) is not guaranteed, if the procedure succeeds, it will yield an h that has the same dependence on  $\sigma_f$  and  $|\phi^{(2)}(t)|$  as the optimal finite-difference interval with respect to its mean-squared error.

VI. Error Bound Estimate on Gradient. Using the h we obtain from our procedure, we can approximate the error bound. Ignoring the  $o(h^{q-d})$  term, the error is approximately given by:

$$\epsilon_g(t;h) \approx |c_q| |\phi^{(q)}(t)| h^{q-d} + ||w||_1 \epsilon_f h^{-d}.$$

As we have shown in (3.3.8), we can bound  $\left|\phi^{(q)}(t)\right|$  by

$$\left|\phi^{(q)}(t)\right| \le \frac{\epsilon_f(r_u+1)}{|c_r| h^q}.$$

Therefore, we can approximately bound the error by

$$\epsilon_g(t;h) \lesssim \left(\frac{|c_q|}{|c_r|}(r_u+1) + \|w\|_1\right) \epsilon_f h^{-d}.$$

In practice, we have found that this error bound is able to obtain an order-of-magnitude of the actual error, but may underestimate the error due to the  $o(h^{q-d})$  term.

# 3.3.2. Finite Termination

Next, we prove a finite termination theorem for Algorithm 3.2. We start by making the following assumptions:

Assumption 3.3.1. The testing ratio  $r_S$  satisfies:

$$|r_S(h;\phi,t,\epsilon_f) - r_S(h;f,t,\epsilon_f)| \le 1, \ \forall t \in \mathbb{R}, \ h > 0$$

This assumption is satisfied by our requirement that  $\|\tilde{w}\|_1 = 1$  and that  $|\epsilon(t)| \leq \epsilon_f$ . Note that this can be easily satisfied by simply rescaling the numerator to ensure that the total noise accumulated in the numerator is bounded by  $\epsilon_f$ .

Assumption 3.3.2. As a function of h,  $r_S(h; \phi, t, \epsilon_f)$  is continuous with  $r_S(0; \phi, t, \epsilon_f) = 0$ , and there exists an integer  $K \in \mathbb{N}$  such that

$$r_S(2^K h_0; \phi, t, \epsilon_f) \ge r_u - 1$$

Assuming that  $\epsilon_f > 0$ , the requirement that  $r_S(0; \phi, t, \epsilon_f) = 0$  is automatically satisfied by validity of the testing ratio (3.3.4). The second part of Assumption 3.3.2, while technical, is satisfied, for example, when  $|\phi^{(q)}(\xi)| \ge \eta > 0$  for all  $\xi \in [\min_j \{t + h\tilde{s}_j\}, \max_j \{t + h\tilde{s}_j\}]$ . With these assumptions, we can now show finite termination.

**Theorem 3.3.3.** Suppose Assumptions 3.3.1 and 3.3.2 are satisfied. In addition, suppose  $r_u$  and  $r_l$  are chosen such that  $0 < r_l < r_u - 2$ . Then, Algorithm 3.2 will terminate successfully in a finite number of iterations.

**Proof.** We assume that  $h \ge 0$ . Assume by contradiction that Algorithm 3.2 does not terminate finitely. We denote the variables l, u, h used at the beginning of the k-th iteration of Algorithm 3.2 as  $l_k, u_k, h_k$ , respectively. Obviously, we have

$$0 \le l_k \le h_k \le u_k, \ \forall k \in \mathbb{N},$$

and

$$l_k \le l_{k+1} < u_{k+1} \le u_k, \ \forall k \in \mathbb{N}.$$

First, we show that  $r_S(l_k; \phi, t, \epsilon_f) < r_l + 1$  for all  $k \in \mathbb{N}$ , by induction on k. Clearly this is true for k = 0 since  $l_0 = 0$ , and we have  $r_S(0; \phi, t, \epsilon_f) = 0$  by Assumption 3.3.2. Suppose the statement holds for  $k \leq K$ . We have two cases: (1)  $r_S(h_K; f, t, \epsilon_f) < r_l$ , which by Assumption 3.3.1 implies  $r_S(h_K; \phi, t, \epsilon_f) \leq r_S(h_K; f, t, \epsilon_f) + 1 < r_l + 1$ . In this case  $l_{K+1} = h_K$ , so  $r_S(l_{K+1}; \phi, t, \epsilon_f) = r_S(h_K; \phi, t, \epsilon_f) < r_l + 1$ . (2)  $r_S(h_K; f, t, \epsilon_f) > r_u$ , in which case  $l_{K+1} = l_K$  so by the induction hypothesis  $r_S(l_{K+1}; \phi, t, \epsilon_f) = r_S(l_K; \phi, t, \epsilon_f) < r_l + 1$ . Therefore the induction hypothesis holds for (K + 1)-th iteration. By a similar argument, we can show that either  $u_k = +\infty$ , or  $u_k < +\infty$  and  $r_S(u_k; \phi, t, \epsilon_f) > r_u - 1$  for all  $k \in \mathbb{N}$ .

In summary, we can show that for all  $k \in \mathbb{N}$ , we have

(3.3.16) either 
$$r_S(l_k; \phi, t, \epsilon_f) < r_l + 1 < r_u - 1 < r_S(u_k; \phi, t, \epsilon_f)$$

(3.3.17) or 
$$r_S(l_k; \phi, t, \epsilon_f) < r_l + 1$$
 and  $u_k = +\infty$ .

Next, we claim that there exists  $K_1 \in \mathbb{N}$  such that  $u_k < +\infty$  for  $k \geq K_1$ . Suppose this is not the case, then we have  $r_S(h_k; f, t, \epsilon_f) < r_l$ ,  $\forall k \in \mathbb{N}$ . In this case, we have  $h_{k+1} = 2l_{k+1} = 2h_k$ , so  $h_k = 2^k h_0$  for all  $k \in \mathbb{N}$ . By Assumption 3.3.2, there exists  $K \in \mathbb{N}$ such that  $r_S(h_K; \phi, t, \epsilon_f) \geq r_u - 1$ , and since  $r_S(h_K; f, t, \epsilon_f) \geq r_S(h_K; \phi, t, \epsilon_f) - 1$ , we have  $r_S(h_K; f, t, \epsilon_f) \geq r_u - 2 > r_l$ , contradicting the inequality  $r_S(h_k; f, t, \epsilon_f) < r_l$ ,  $\forall k \in \mathbb{N}$ . This proves the existence of  $K_1$ .

We are now ready to present the contradiction. For  $k \ge K_1$ , since  $u_k < \infty$ , we have

$$u_{k+1} - l_{k+1} = \frac{1}{2} \left( u_k - l_k \right)$$

This implies that  $u_k - l_k \to 0$ . Since  $r_S(h; \phi, t, \epsilon_f)$  (as a function of h) is continuous and  $u_{K_1} < +\infty$ ,  $[0, u_{K_1}]$  is compact so  $r_S(h; \phi, t, \epsilon_f)$  (as a function of h) is uniformly continuous on  $[0, u_{K_1}]$ . Note that  $l_k, u_k \in [0, u_{K_1}]$  for  $k \ge K_1$ , therefore we have

$$r_S(u_k; \phi, t, \epsilon_f) - r_S(l_k; \phi, t, \epsilon_f) \to 0$$

This contradicts the fact that

$$r_S(l_k; \phi, t, \epsilon_f) < r_l + 1 < r_u - 1 < r_S(u_k; \phi, t, \epsilon_f), \ \forall k \in \mathbb{N}, \ k \ge K_1$$

Therefore, Algorithm 3.2 must terminate finitely. Clearly, whenever it terminates, the output  $h_R$  must satisfy

$$r_S(h_R; f, t, \epsilon_f) \in [r_l, r_u].$$

## **3.4.** Numerical Experiments

In this section, we present numerical results demonstrating the reliability of our finitedifference interval estimation procedure. We first utilize the method for computing first and second derivatives of commonly tested functions, with added noise. We then insert our procedure into a standard L-BFGS implementation and demonstrate its usefulness on a subset of synthetic noisy CUTEst problems [42]. All methods were implemented in Python 3.

# 3.4.1. Finite-Difference Interval Estimation

We first test our proposed procedure on several univariate functions. We focus on the case where d = 1 and 2 as this is most relevant to optimization. We test Algorithm 3.2 using 6 different estimating schemes, shown in Table 3.1. The testing ratios are generated using formula (3.3.14) with different choices of  $\alpha$ . The  $\alpha$  for each scheme is chosen as the smallest integer such that  $r^* > \beta = 2$ .
labeldswq $\alpha$ $r^*$ CommentFD1 $(0,1)$ $(-1,1)$ 243forward differenceCD1 $(-1,1)$ $(-1/2,1/2)$ 333central differenceFD_3P1 $(0,1,2)$ $(-3/2,2,-1/2)$ 333.69forward difference w/ 3 pointFD_4P1 $(0,1,2,3)$ $(-11/6,3,-3/2,1/1)$ 438.25forward difference w/ 4 pointCD_4P1 $(-2,-1,1,2)$ $(1/12,-2/3,2/3,-1/12)$ 522.5central difference w/ 4 point								
FD1 $(0,1)$ $(-1,1)$ 243forward differenceCD1 $(-1,1)$ $(-1/2,1/2)$ 333central differenceFD_3P1 $(0,1,2)$ $(-3/2,2,-1/2)$ 333.69forward difference w/ 3 pointFD_4P1 $(0,1,2,3)$ $(-11/6,3,-3/2,1/1)$ 438.25forward difference w/ 4 pointCD_4P1 $(-2,-1,1,2)$ $(1/12,-2/3,2/3,-1/12)$ 522.5central difference w/ 4 point	label	d	S	w	q	$\alpha$	$r^*$	Comment
CD 1 $(-1,1)$ $(-1/2,1/2)$ 3 3 3 central difference FD_3P 1 $(0,1,2)$ $(-3/2,2,-1/2)$ 3 3 3.69 forward difference w/ 3 por FD_4P 1 $(0,1,2,3)$ $(-11/6,3,-3/2,1/1)$ 4 3 8.25 forward difference w/ 4 por CD_4P 1 $(-2,-1,1,2)$ $(1/12,-2/3,2/3,-1/12)$ 5 2 2.5 central difference w/ 4 points CD_4P 2 - (-1,0,1) - (1/2,-2/3,2/3,-1/12) - (-2,-2/3,2/3,-1/12)	FD	1	(0,1)	(-1, 1)	2	4	3	forward difference
FD_3P1 $(0,1,2)$ $(-3/2,2,-1/2)$ 333.69forward difference w/ 3 poFD_4P1 $(0,1,2,3)$ $(-11/6,3,-3/2,1/1)$ 438.25forward difference w/ 4 pointCD_4P1 $(-2,-1,1,2)$ $(1/12,-2/3,2/3,-1/12)$ 522.5central difference w/ 4 point	CD	1	(-1, 1)	(-1/2, 1/2)	3	3	3	central difference
FD_4P 1 $(0,1,2,3)$ $(-11/6,3,-3/2,1/1)$ 4 3 8.25 forward difference w/ 4 point CD_4P 1 $(-2,-1,1,2)$ $(1/12,-2/3,2/3,-1/12)$ 5 2 2.5 central difference w/ 4 point CD_4P 2 (1,0,1)	FD_3P	1	(0, 1, 2)	(-3/2, 2, -1/2)	3	3	3.69	forward difference w/ 3 points
CD_4P 1 $(-2, -1, 1, 2)$ $(1/12, -2/3, 2/3, -1/12)$ 5 2 2.5 central difference w/ 4 point of the second difference w/ 4 point of the seco	FD_4P	1	(0, 1, 2, 3)	(-11/6, 3, -3/2, 1/1)	4	3	8.25	forward difference w/ 4 points
	$CD_4P$	1	(-2, -1, 1, 2)	(1/12, -2/3, 2/3, -1/12)	5	2	2.5	central difference w/ 4 points
$L2\_CD$ 2 $(-1,0,1)$ $(1,-2,1)$ 4 2 3 2nd-order central difference	$L2_CD$	2	(-1, 0, 1)	(1, -2, 1)	4	2	3	2nd-order central difference

Table 3.1. Schemes for approximating the *d*-th order derivative used in the experiments. The scheme is defined by S = (w, s) as in (3.3.1); *q* is defined in (3.3.2).

For a specific testing function  $\phi$  at point t with noise  $\epsilon_f$  and scheme S = (w, s), we plot the worst case relative error, as a function of differencing interval h, defined as:

$$\delta_S(h;\phi,t,\epsilon_f) = \frac{1}{|\phi^{(d)}(t)|} \left[ \left| \frac{\sum_{j=1}^p w_j \phi(t+s_j h)}{h^d} - \phi^{(d)}(t) \right| + \|w\|_1 \frac{\epsilon_f}{h^d} \right].$$

This function captures the relative error of the estimation scheme S on the noisy function f at t, in the worst case. The differencing interval h that minimizes  $\delta_S(h; \phi, t, \epsilon_f)$  is the optimal h. Notice that  $\delta_S(h; \phi, t, \epsilon_f)$  is a deterministic function that does not rely on the realization of actual noise in f(t).

We manually inject uniformly distributed, stochastic noise into  $\phi$ ,

$$f(t) = \phi(t) + \epsilon(t), \ \epsilon(t) \sim \text{Uniform}(-\epsilon_f, \epsilon_f),$$

independent of all other quantities. We then apply Algorithm 3.2 to obtain  $h_{\dagger}$ . We plot  $h_{\dagger}$ and observe how far it is from the minimizer of  $\delta_S(h; \phi, t, \epsilon_f)$ . In Appendix B.2, we report the minimizer of the function  $\delta_S(h; \phi, t, \epsilon_f)$  obtained by scipy.optimize.minimize\_scalar.



Figure 3.1. Worst case relative error  $\delta_S(h; \phi, t, \epsilon_f)$  for forward and central differences against h on function  $\phi(t) = \cos(t)$  with different noise levels; the vertical dashed line represents the  $h_{\dagger}$  output by Algorithm 3.2.

**3.4.1.1.** Robustness to Different Noise Levels. To demonstrate that our method is reliable across a range of noise levels, we test our adaptive procedure for both forward and central differences (FD, CD) on the simple function  $\phi(t) = \cos(t)$  for different noise levels. The plot of the worst case relative error and obtained interval  $h_{\dagger}$  are illustrated in Figure 3.1. Our results demonstrate that our method performs consistently well across a range of noise levels. For all figures on all finite-difference schemes listed in Table 3.1 and complete numerical results, see Appendix B.2.

**3.4.1.2.** Difficult and Special Examples. In this subsection, we consider examples of difficult functions given in [40] and [93]. These examples include:

(1)  $\phi(t) = (e^t - 1)^2$ , at t = -8. This function has extremely small first and secondorder derivative at t = -8, but quickly increases as t increases beyond t = 0; a naive choice of  $h = \sqrt{\epsilon_f / |\phi^{(2)}(t)|}$  for forward differences can result in an extremely large h and lead to huge error.

(2) 
$$\phi(t) = e^{100t}$$
, at  $t = 0.01$ .

- (3)  $\phi(t) = t^4 + 3t^2 10t$ , at t = 0.99999. This function is considered difficult because  $\phi'(1) = 0$ , and represents a case where the estimated derivative is very close to 0. In addition, this function is a fourth-order polynomial, so the optimal h for CD\_4P is  $+\infty$ .
- (4)  $\phi(t) = 10000t^3 + 0.01t^2 + 5t$ , at  $t = 10^{-9}$ . This example is difficult in that it has approximate central symmetry at t = 0, which can lead to issues for adaptive procedures such as those proposed in [40].

For each example, we again fix  $\epsilon_f = 10^{-3}$ , and perform our estimation procedure for different schemes, and plot the worst case relative error. The results can be found in Figure 3.2.

For the first two examples, we see that our procedure is able to estimate the derivative well even when the function increases rapidly. These are cases where using our adaptive procedure is significantly more effective than computing an interval based on higher-order derivative information at the point of interest, as observed in [93].

It is also interesting to observe the results for the two polynomials. For  $\phi(t) = t^4 + 3t^2 - 10t$  and scheme CD\_4P, our procedure generates a large  $h_{\dagger}$ ; this is consistent with the fact that scheme CD\_4P has q = 5, and  $\phi^{(5)}(\xi) = 0$  for all  $\xi$  on this example, which implies that we should choose h to be as large as possible. This similarly holds true for the schemes FD\_4P and CD\_4P on the function  $\phi(t) = 10000t^3 + 0.01t^2 + 5t$ .

While theoretically speaking we should choose  $h = \infty$  in such cases, we can observe in Figure 3.2 that this is not the case. When plotting the worst-case relative error, we see that there exists a large h such that  $\delta_S(h; \phi, t, \epsilon_f)$  is minimized, beyond which the relative error begins to sharply increase. This phenomenon is due to round-off error. When h becomes



Figure 3.2. Worst case relative error  $\delta_S(h; \phi, t, \epsilon_f)$  against h on several special cases; the vertical dashed line represents the  $h_{\dagger}$  output by Algorithm 3.2.

too large, round-off error (which is multiplicative) will dominate  $\epsilon_f$ ; this approximately happens when  $\max_j |\phi(t+s_jh)| \epsilon_M$  becomes comparable to  $\epsilon_f$ .

# 3.4.2. Finite-Difference L-BFGS

In order to show the robustness of our procedure, we apply it within the L-BFGS method. We now let  $\phi$  denote a smooth multivariate function,  $\phi : \mathbb{R}^n \to \mathbb{R}$ , and consider the problem

(3.4.1)  $\min_{x \in \mathbb{R}^n} \phi(x),$ 

while only provided noisy function evaluations of the form

(3.4.2) 
$$f(x) = \phi(x) + \epsilon(x), \ \epsilon(x) \sim \text{Uniform}(-\epsilon_f, \epsilon_f),$$

where  $\epsilon_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ . We perform our tests on a subset of synthetically generated noisy CUTEst problems [42] detailed in Table 3.2.

Problem	Dim $(n)$						
AIRCRFTB	5	CRAGGLVY	100	FREUROTH	100	PFIT4LS	3
ALLINITU	4	CUBE	2	GENROSE	100	QUARTC	100
ARWHEAD	100	DENSCHND	3	GULF	3	SINEVAL	2
BARD	3	DENSCHNE	3	HAIRY	2	SINQUAD	100
BDQRTIC	100	DIXMAANH	90	HELIX	3	SISSER	2
BIGGS3	3	DQRTIC	100	NCB20B	100	SPARSQUR	100
BIGGS5	5	EDENSCH	36	NONDIA	100	TOINTGSS	100
BIGGS6	6	EIGENALS	110	NONDQUAR	100	TQUARTIC	100
BOX2	2	EIGENBLS	110	OSBORNEA	5	TRIDIA	100
BOX3	3	EIGENCLS	30	OSBORNEB	11	WATSON	31
BRKMCC	2	ENGVAL1	100	PENALTY1	100	WOODS	100
BROWNAL	100	EXPFIT	2	PFIT1LS	3	ZANGWIL2	2
BROWNDEN	4	FLETCBV3	100	PFIT2LS	3		
CLIFF	2	FLETCHBV	100	PFIT3LS	3		

Table 3.2. Subset of unconstrained CUTEst problems and their problem dimensions [42].

The L-BFGS method has the form

(3.4.3) 
$$x_{k+1} = x_k - \alpha_k H_k g(x_k),$$

where  $g(x_k)$  is a finite-difference approximation to the gradient,  $H_k$  is the L-BFGS matrix with memory of 10 (see [75]), and  $\alpha_k$  is a steplength selected by a relaxed Armijo-Wolfe line search designed to handle noise. To describe the line search, let  $\alpha_k^j$  denote the *j*th trial steplength at iteration *k*. Similar to Shi et al.[90], the Armijo condition is relaxed as follows:

(3.4.4)

$$f(x_{k} + \alpha_{k}^{j}p_{k}) \begin{cases} \leq f(x_{k}) + c_{1}\alpha_{k}^{j}g(x_{k})^{T}p_{k} & \text{if } j = 0, \ g(x_{k})^{T}p_{k} < -\epsilon_{g}(x_{k})\|p_{k}\| \\ \leq f(x_{k}) + c_{1}\alpha_{k}^{j}g(x_{k})^{T}p_{k} + 2\epsilon_{f} & \text{if } j \geq 1, \ g(x_{k})^{T}p_{k} < -\epsilon_{g}(x_{k})\|p_{k}\| \\ < f(x_{k}) & \text{if } g(x_{k})^{T}p_{k} \geq -\epsilon_{g}(x_{k})\|p_{k}\|, \end{cases}$$

where  $c_1 = 10^{-4}$ ,  $c_2 = 0.9$ , and  $\epsilon_g(x_k)$  is the estimated gradient error described below. Thus, we relax the line search only when the gradient is reliable; otherwise, we enforce simple decrease. We check the Wolfe condition by evaluating the directional derivative  $\nabla \phi(x)^T p$  using finite differences along the direction p, as in [90].

**3.4.2.1.** Forward Differences. In the first set of experiments, the gradient approximation  $g(x_k)$  is obtained by forward differences,

$$[g(x_k)]_i = \frac{f(x + h_i e_i) - f(x)}{h_i}, \quad i = 1, \dots n_i$$

where the differencing interval  $h_i$  is determined by one of the following three strategies.

1. Fixed. The interval h is fixed across all components i and for the entire iteration. This strategy tries to emulate the common practice of hand-tuning  $h_i$  at the start using problem specific information. We simulate this using the formula

(3.4.5) 
$$h = 2\sqrt{\frac{\epsilon_f}{L_2}}, \text{ where } L_2 = \max\left\{10^{-1}, \sqrt{\frac{\sum_{i=1}^n [\nabla^2 \phi(x_0)]_{ii}^2}{n}}\right\},$$

which assumes that the diagonals of the Hessian are known. The gradient error is approximated assuming  $L_2$  is correct, that is,  $\epsilon_g(x) = 2\sqrt{nL_2\epsilon_f}$ . We created this option for benchmarking purposes only.

2. MW. The Moré-Wild heuristic for estimating and interval  $h_i$  for every component i of the gradient. We set  $L_2 = \max\{10^{-1}, \hat{L}_2\}$ , where  $\hat{L}_2$  is the estimate given by the Moré and Wild heuristic. If the heuristic fails, we set  $L_2 = 10^{-1}$ . The gradient error is estimated similar to Fixed but componentwise, i.e.,  $\epsilon_g(x) = 2\sqrt{\epsilon_f \sum_{i=1}^n L_{2,i}}$ .

3. Adaptive Our adaptive procedure for estimating  $h_i$  along each component using Algorithm 3.1.

For the MW and Adaptive strategies, we re-estimate the second derivative or finitedifference interval whenever a partial derivative needs to be approximated. For example, when computing the full gradient, we estimate the finite-difference interval along each coordinate direction separately. We chose not to compare against Gill et al. [40] as we regard the Moré-Wild heuristic to be an improvement over their approach.

We present results for a few representative problems in Figure 3.3. The L-BFGS method described above is terminated if no further progress is made on the objective function over 5 consecutive iterations. Figure 3.3 plots the optimality gap  $\phi(x_k) - \phi^*$  against the number of function evaluations. The optimal value  $\phi^*$  is obtained by solving the original problem to completion without noise with L-BFGS. While we found Moré and Wild's heuristic to work well for  $\epsilon_f < 10^{-1}$ , their heuristic fails frequently for the case where  $\epsilon_f = 10^{-1}$ . (This can be seen in the complete results presented in Appendix B.2.) For this reason, we report results for  $\epsilon_f = 10^{-1}$  and  $10^{-5}$  to demonstrate the robustness of our algorithm compared



Figure 3.3. Comparison of forward-difference L-BFGS methods with difference intervals determined using a fixed interval, the Moré and Wild heuristic, and our adaptive algorithm. Comparisons are made on representative problems with noise level  $\epsilon_f = 10^{-1}$  (top) and  $10^{-5}$  (bottom). The solid line plots the observed function value and the dashed line plots the true function value. The dashed black line shows the noise level  $\epsilon_f$  of the function.

to Moré and Wild for different noise levels. When Moré and Wild's heuristic succeeds, we observe that our algorithm (Adaptive) is able to more efficiently achieve comparable accuracy to the Moré and Wild heuristic, while attaining more accurate solutions than using a fixed interval for some problems. The lack of accuracy in the Fixed strategy can be explained by inability for a fixed interval to adapt to changes in the Hessian over the course of the iteration — an exception being the TRIDIA problem, which is very well scaled.

**3.4.2.2. Central Differences.** In the second set of experiments, we employ central differences,

$$[g(x;h)]_i = \frac{f(x+h_i e_i) - f(x-h_i e_i)}{2h_i}.$$

The differencing interval is determined via a Fixed strategy or the Adaptive procedure described in Algorithm 3.2. (The Moré and Wild's heuristic does not apply to this case.)



Figure 3.4. Comparison of central-difference L-BFGS methods with difference intervals determined using a fixed interval and our adaptive algorithm. Comparisons are made on representative problems with noise level  $\epsilon_f = 10^{-5}$ . The solid line plots the observed function value and the dashed line plots the true function value. The dashed black line shows the noise level  $\epsilon_f$  of the function.

For the Fixed strategy, we choose

(3.4.6)

$$h = \sqrt[3]{\frac{3\epsilon_f}{L_3}}, \text{ where } L_3 = \max\left\{10^{-1}, \sqrt{\frac{1}{n}\sum_{i=1}^n \left(\frac{[\nabla^2 \phi(x_0 + \tilde{h}e_i)]_{ii} - [\nabla^2 \phi(x_0)]_{ii}}{\tilde{h}}\right)^2}\right\}$$

and  $\tilde{h} = \max\{1, |[x_0]_i|\}\sqrt{\epsilon_M}$ . Note that noiseless forward differences are applied to the true Hessian to estimate the third derivative along each coordinate direction at the initial point. This synthetic **Fixed** strategy is presented for benchmarking purposes; it is not generally viable in practice.

Representative results are shown in Figure 3.4. Similar to the forward-difference case, our algorithm is able to obtain higher accuracy in the solution compare to the Fixed strategy, but at higher cost as expected. Complete experimental results for all problems and noise levels are presented in Appendix B.2.

## 3.5. Final Remarks

We have developed a principled and robust procedure for determining the difference interval for estimating gradients in optimization methods, assuming that the noise level is known. Our procedure applies to any finite-difference scheme, including central- and higher-order difference schemes. It performs a bisection search on a ratio that balances the truncation and measurement errors such that one typically attains a near-optimal difference interval. Whereas some methods for estimating the difference interval prioritize efficiency, such as Moré and Wild [71], and others compromise cost and accuracy, such as Gill, et al. [40], our approach is designed to be as robust as possible so that finite-difference gradient approximations can be reliably used in established nonlinear optimization techniques for solving noisy problems.

As demonstrated in our experiments, reusing previous difference intervals from prior iterations allows us to reduce the cost of the estimation procedure. Additional savings can achieved by re-estimating the difference interval periodically; for simple problems only a few times during the course of the optimization will suffice. The ability to exploit parallelism by distributing the computation of the gradient is an advantage that should not be underestimated when comparing the finite-difference approach with other techniques for derivative-free optimization.

# CHAPTER 4

# A Feasible Nonlinear Programming Approach for Constrained Derivative-Free Optimization

### 4.1. Introduction

Interpolation-based trust-region optimization (IBO) methods are among the most efficient and reliable techniques for unconstrained derivative-free optimization (DFO) [69]. Yet extending IBO methods to general constrained optimization is not straightforward and has received limited attention in the literature [59]. In this paper, we investigate a method first proposed by Conn et al. [26] designed for problems where the constraints and their derivatives are available analytically. This is a feasible method that constructs a quadratic model of the objective using an IBO approach and computes a step by minimizing this model subject to the original constraints of the problem and a trust region. The step computation is in general a nonlinear optimization sub-problem that is solved using a general purpose nonlinear solver. We refer to this approach as FIBO, for feasible interpolation-based optimization method. The numerical results presented in [26], and subsequently [23], are inconclusive and the FIBO approach is rarely, if ever, used in practice. In this paper, we present numerical results that suggest that, for an important class of practical applications, our proposed implementation of FIBO is competitive with the best methods proposed to date. The problem under consideration is

 $(4.1.1) \qquad \qquad \min_{x} \quad f(x)$ 

$$(4.1.3) c_i(x) \le 0, \quad i \in \mathcal{I}$$

for some finite index sets  $\mathcal{E}$  and  $\mathcal{I}$ , and where  $f : \mathbb{R}^n \to \mathbb{R}$  and  $c_i(x) : \mathbb{R}^n \to \mathbb{R}$ ,  $i \in \mathcal{I} \cup \mathcal{E}$ , are smooth functions. We assume that the derivatives of f are not available but those of the constraints  $c_i(x)$  are.

Applications of this problem setting are often found in practical applications in which the objective is provided by a simulation model and the constraints are given analytically, involving say trigonometric functions, or linear constraints. To cite a concrete example, in constrained reinforcement learning a single reward objective may not be sufficient and various constraints must be enforced to ensure safety requirements or physical limitations of the task [54].

#### 4.1.1. Literature Review

Early methods for analytically constrained derivative-free optimization followed heuristic penalty and augmented Lagrangian approaches in which pattern search methods are employed to solve the sub-problems; see [59] for a survey of these methods. An important drawback of this approach is that derivative information of the constraints is not exploited. SID-PSM requires derivatives of the constraints but is difficult to scale to high dimensions [85]. In addition, as demonstrated by Moré and Wild [69] and Sahinidis [85], pattern search methods are inefficient compared to IBO methods, both on smooth and nonsmooth DFO problems.

Powell developed efficient IBO methods and software to deal with increasingly more complex sets of constraints. BOBYQA [79] is designed for problems with simple bounds; it builds upon the minimum-Frobenius norm update method first introduced in NEWUOA. Powell's last code, LINCOA, is able to deal with linear constraints [80]. These methods are quite sophisticated as they seek to minimize linear algebra costs. They are accessible through the Python and Matlab interface developed by Zaikun [84]. Powell did not develop an IBO methods to handle general (nonlinear) constraints other than COBYLA, which solves general inequality constrained DFO problems by constructing linear models of the objective and constraints. An extension of this approach, named COBYQA, was proposed by Ragonneau [83] to handle both equality and inequality constraints via quadratic models.

A variety of IBO methods have been proposed to enforce feasibility of the iterates. When the constraints are analytically available, this set of methods often impose the constraints in the trust region subproblem at every iteration. This line of work was first extended to convex constraints by [24], in which the authors provide convergence analysis assuming that the local model is always fully-linear. [51] extended the global convergence analysis by proposing a generalized fully-linear model in the general convex constrained case, relaxing this fully-linear model assumption. However, [51] only provides numerical results on least-square functions and neither work demonstrate the effectiveness of their approach on general objective functions. Limited work has been done when general constraints are present, as the resulting subproblem from the constraint enforcement is in general computationally expensive to solve. To our knowledge, this approach was first explored by Conn et. al [26]. Yet the numerical experiments do not provide convincing evidence regarding the effectiveness of the method due to the ill-considered experiment settings. In particular, the other DFO methods being considered do not exploit the availability of constraint derivatives. [23] revisited this approach but under the same reason, does not provide additional insights into this framework. No convergence proof has yet been established for this feasible approach with general constraints. A similar idea is again explored in CONDOR, which is designed to solve general inequality constraints by treating bound and linear constraints via active-set methods and the remaining with SQP [11]. For a thorough treatment of IBO methods, see [29].

A recent line of work in unconstrained DFO focuses on methods that employ finite difference approximations to the gradient [7, 73, 93]. This approach can be effective even for noisy functions, if the finite difference interval is chose appropriately. Given that state-of-the-art nonlinear programming solvers normally include a finite difference option for both functions and constraints, they should serve as a benchmark for testing constrained DFO methods.

#### 4.2. The FIBO Algorithm

Let us begin by sketching a basic IBO method for unconstrained derivative-free optimization, and refer the reader to [29, 78] for a discussion of many fine points of implementation. The algorithm starts by evaluating the objective function at a set  $Y_0$  of points on a stencil around the initial point  $x_0$ . The algorithm then sets  $k \leftarrow 0$  and constructs a quadratic model

(4.2.1) 
$$m_k(x_k+s) = f(x_k) + g_k^T s + \frac{1}{2} s^T H_k s.$$

This model can be defined in two different ways. Assuming that  $Y_k$  is *poised* and that  $|Y_0| = (n+1)(n+2)/2$ , this model is uniquely specified by requiring that it interpolate the function values at the current interpolation set. An alternative IBO approach that greatly reduces the linear algebra costs of the iteration employs only  $|Y_k| = 2n + 1$  interpolation points and defines a model through a minimum Frobenius norm update of the matrix  $H_k$  subject to  $|Y_k|$  interpolation conditions.

Regardless of how the model is constructed, the trial step  $s_k$  is given by the solution of the trust region subproblem

(4.2.2a) 
$$\min_{s} \quad m_k(x_k+s)$$

(4.2.2b) s.t. 
$$||s||_2 \le \Delta_k$$

As in any trust region method, acceptance of the step  $s_k$  is determined based on a ratio

(4.2.3) 
$$\rho_k = \frac{f(x_k) - f(x_k + s_k)}{m_k(x_k) - m_k(x_k + s_k)}$$

of actual vs predicted reduction in the objective. If sufficient decrease is obtained ( $\rho_k \ge \eta$ for some  $\eta \in (0, 1)$ ), the new iterate is defined as  $x_{k+1} = x_k + s_k$ ; otherwise the step is rejected and one sets  $x_{k+1} = x_k$ .

The trust region  $\Delta_k$  is increased for successful steps. Otherwise, the algorithm considers two cases. If the interpolation set is not poised, (i.e., if the interpolation points (nearly) lie in a linear subspace of  $\mathbb{R}^n$ ),  $\Delta_k$  remains unchanged to avoid premature shrinkage of the trust region and  $Y_k$  is improved via a geometry phase. If the step is rejected and  $Y_k$  is poised, then  $\Delta_k$  is decreased. The set  $Y_k$  is normally updated at every iteration by removing one point from  $Y_k$ (for example, the furthest from  $x_{k+1}$ ) and by adding  $x_{k+1}$  to this set. This trust region interpolation-based approach has proved to be a very efficient for solving unconstrained DFO problems.

We can now describe FIBO, the extension of this method to the constrained optimization problem (4.2.2). We assume—and this is important—that the constraints are both analytically available and inexpensive to evaluate compared to the objective function. Then, it is realistic to compute the step by

(4.2.4a) 
$$\min_{s} m_{k}(x_{k}+s) = f(x_{k}) + g_{k}^{T}s + \frac{1}{2}s^{T}H_{k}s$$

(4.2.4b) s.t. 
$$c_i(x_k + s) = 0, \quad i \in \mathcal{E}$$

$$(4.2.4c) c_i(x_k+s) \le 0, \quad i \in \mathcal{I}$$

$$(4.2.4d) ||s||_2 \le \Delta_k.$$

The model  $m_k$  in (4.2.4a) is constructed by interpolating function values at a set of interpolation points, and as in the unconstrained case, models the objective function f. Note that (4.2.4) includes the original constraints of the problem and is, in general, a nonlinear programming problem that must be solved using any general-purpose nonlinear programming solver (KNITRO in our experiments).

Let us assume for now that (4.2.4) can be solved to optimality. If the initial iterate is feasible, then all subsequent iterates will be feasible and it is therefore sufficient for  $m_k$ to model only the objective and disregard the contribution of the constraints. All other details of the algorithm, such as the update of the interpolation set and trust region, are defined as in the unconstrained IBO method. In fact, to develop the FIBO method one can start with any IBO method for unconstrained optimization and simply replace the step computation (4.2.2) by (4.2.4).

This framework therefore provides substantial flexibility in that any IBO method can be used to construct the model and handle the trust region logic, and any nonlinear programming code can be invoked to solve the trust region subproblem (4.2.4). We state the proposed method in Algorithm 4.1.

# Algorithm 4.1. FIBO

- 1: Parameters:  $0 \le \eta < 1$ , and  $\Delta_0 > 0$ .
- 2: Choose  $x_0$ , a **feasible** starting point, and construct an initial poised interpolation set  $Y_0$ .

3: Let k = 0.

- 4: while no convergence test is satisfied do
- 5: Build a local (quadratic) model (4.2.1) using interpolation set  $Y_k$ .
- 6: Compute a step  $s_k$  by solving the trust region subproblem (4.2.4).

7: Compute the ratio defined by 
$$(4.2.3)$$

8: if 
$$\rho_k \ge \eta$$
 then

9: Set  $x_{k+1} = x_k + s_k$ .

10: Choose  $\Delta_{k+1} \ge \Delta_k$ .

- 11: Update  $Y_k$  to include  $x_{k+1}$ .
- 12: else if  $Y_k$  needs to be improved then

13: Set 
$$x_{k+1} = x_k$$
.

- 14: Set  $\Delta_{k+1} = \Delta_k$ .
- 15: Improve  $Y_k$  using a geometry-improving procedure.

```
16: else
```

- 17: Set  $x_{k+1} = x_k$ .
- 18: Choose  $\Delta_{k+1} < \Delta_k$ .
- 19: Update  $Y_k$  to include  $x_k + s_k$ .
- 20: end if

21:  $Y_{k+1} = Y_k$ 

$$22: \qquad k = k + 1$$

# 23: end while

We now comment on the following salient properties of this approach.

Application to Noisy DFO. The FIBO method can be applied to noisy DFO problems without additional modifications. As noted in [93], IBO methods are normally efficient and robust in the presence of noise for reasons that are not well understood, and one can expect these benefits in the constrained case as well. An added advantage is that information regarding the noise level in the problem is not required, compared to DFO methods that employ finite differences.

Feasibility. Assuming that the nonlinear programming solver is successful at every iteration, all iterates of the FIBO algorithm are feasible. This is a desirable property for DFO problems in which the objective function is not defined outside the feasible set. (The nonlinear programming method used to solve the subproblem (4.2.4) typically generates infeasible iterates, but it is only the quadratic objective  $m_k$  of this subproblem that is evaluated at those points and not the original objective function.) Although the success of the nonlinear solver is not guaranteed since the subproblem (4.2.4) is not convex, we did not observe any failures in our experiments—and we should note that other approaches for constrained DFO [59] do not provide convergence guarantees in all situations, either.

Feasible Interpolation. As discussed in [59], sufficiently good model accuracy can be difficult or impossible to obtain when only feasible points are interpolated. A similar situation occurs when the objective cannot be evaluated at infeasible points. The initial interpolation set has to be constructed differently to include only feasible points. We do not explore this possibility in this work.

Per-Iteration Cost. The three main expenses of the FIBO iteration are:

- (1) Evaluation of the objective f. This cost is problem dependent but there are many applications where the function evaluation is much more costly than all other computations in the algorithm.
- (2) Construction of the quadratic model. Computing the model  $m_k$  from scratch at every iteration requires  $O(n^6)$  flops, which is quite expensive when the number of variables is large Nevertheless, Powell's approach implemented in NEWUOA requires only  $O(n^2)$  flops since it employs only O(n) interpolation points.
- (3) Solution of the Trust Region Problem. When the number of variables n is very small (say less than 100), the cost of solving (4.2.4) is of no concern even if the nonlinear programming algorithm performs many iterations to return an accurate solution, *provided* the constraints are inexpensive to evaluation. But as the number of variables becomes large the total linear algebra cost on the nonlinear programming is of concern. Note, in particular that since the Hessian of  $m_k$  is dense, each evaluation of  $m_k$  requires  $O(n^2)$  flops. Similarly, if the cost of evaluating the constraints and their derivatives is high, then the FIBO approach is not viable given the need for multiple evaluations for the solution of a single trust region step. In the next section, we quantify more precisely the class of problems for which FIBO is efficient.

In summary, the FIBO method is appropriate for problems with

- expensive function evaluation (at least as expensive as  $O(n^2)$  operations),
- inexpensive constraint evaluation,

such that the cost associated with solving the subproblem (4.2.4) and solving the interpolation system is dominated by computational overhead of the objective evaluation.

#### 4.3. Numerical Experiments

We adapt the Python code for DFOTR to incorporate general constraints. In particular, we replace the original module for solving trust region subproblem by calling KNITRO to solve the general constrained subproblem. Note that the derivative information of this subproblem (4.2.4) is available. Since DFOTR picks the interpolation point with the minimum objective value as the best iterate, we disable this option as such point may be infeasible.

We will compare FIBO with finite-difference based KNITRO (denoted by FD), which has been demonstrated to be effective for noiseless problems [93]. We do not present numerical results for solving noisy constrained problems as we believe the problems can be solved similarly as in the nonnoisy constrained case. The details of the solvers are described below.

- FIBO. We employ the code DFOTR to perform trust region updates, with final trust region radius  $10^{-8}$ . The maximum number of function evaluations allowed is set to be  $500 \times \max(m, n)$ , where n is the number of variables and m the number of constraints. We set the stop\_predict convergence test to 0 avoid early termination caused by the ratio test. All remaining parameters are set to default. We use KNITRO with *exact gradients* to solve the trust region subproblem and the parameters of KNITRO are set to default except alg = 4 (SQP) and hessopt = 6 (L-BFGS).
- FD. We ran KNITRO with alg = 4 (SQP) and hessopt = 6 (L-BFGS). The memory size is set to be 10 (default value for lmsize). We set xtol = 1e-8,

**xtol\_iters = 1**, and **findiff\_terminate=0** to obtain similar convergence criterion as DFOTR. Maximum number of function evaluations allowed is  $500 \times \max(m, n)$ . We perform *forward differencing* to the objective and provide exact derivative information of constraints to the algorithm.

We consider 38 general constrained problems and assume that the initial points  $x_0$  are feasible. Since the initial point provided by CUTEst is usually infeasible, we ran KNITRO to obtain a feasible point by replacing the objective by 0. The cost to obtain a feasible solution is omitted for now, as we assume the constraints are not expensive to evaluate. Note that the observations remain the same when we include the cost of obtaining a feasible starting point for FIBO and run FD from an infeasible initial point  $x_0$ . See C.2 for detailed numerical results.

We study the performance of FIBO by comparing its final objective accuracy and the associated cost. To study the accuracy of the solvers, we measure accuracy with log-ratios:

(4.3.1) 
$$\log_2(\frac{\max(f_{\text{FIBO}} - f^*, 10^{-8})}{\max(f_{\text{FD}} - f^*, 10^{-8})})$$

where  $f_{\text{FIBO}}$  and  $f_{\text{FD}}$  are the final accuracy obtained by the corresponding approach, and the optimal value  $f^*$  is obtained by running KNITRO with exact gradient. We compare the final accuracy of the two methods up to 8 digits and report the ratios using log-ratio plot in Figure 4.1 [67]. The ratios (4.3.1) are plotted in increasing order such that the area of the shaded region provides a rough idea of the relative performance of the methods. FD achieves higher accuracy than FIBO, but not by a wide margin. For full numerical results, see Appendix C.2.



Figure 4.1. Log-ratio Plot for Comparing the Final Accuracy (4.3.1) of FIBO and FD.

To understand the efficiency of the compared algorithms, we terminate the algorithms when

(4.3.2) 
$$f(x_k) \le f^* + \tau \cdot \max(1, |f^*|)$$

where  $f^*$  is the optimal solution obtained by running KNITRO with exact gradients on the test problems, and  $\tau$  is the desired accuracy level. In particular, we look at the number of objective evaluations and constraint evaluations to understand the associated cost to reach a given accuracy level. We denote the number of function evaluations required to satisfy (4.3.2) by evals<sub>FIBO</sub> and evals<sub>FD</sub> respectively for the corresponding method. Similarly, we use cevals<sub>FIBO</sub> and evals<sub>FD</sub> to denote the number of constraint evaluations required to satisfy (4.3.2). If the condition (4.3.2) cannot be satisfied by an algorithm, we simply set the corresponding quantity as a very large number. We specifically present results for  $\tau = 10^{-1}, 10^{-3}, 10^{-5}$  and  $10^{-7}$  to study the efficiency for reaching different accuracy levels. Again, we summarize the results using the log-ratio profile proposed by Morales [67] and plot the quantities

(4.3.3) 
$$\log_2(\frac{\text{evals}_{\text{FIBO}}}{\text{evals}_{\text{FD}}}), \quad \log_2(\frac{\text{cevals}_{\text{FIBO}}}{\text{cevals}_{\text{FD}}}).$$

The plots for comparing the number of objective evaluations and constraints evaluations are reported in Figure 4.2 and 4.3.



Figure 4.2. Log-ratio plot comparing FIBO and FD in terms of the number of function evaluations to satisfy (4.3.2) for  $\tau = 10^{-1}$  (upper left),  $10^{-3}$  (upper right),  $10^{-5}$  (bottom left),  $10^{-7}$  (bottom right).

As indicated by Figure 4.2, FIBO outperforms finite-difference based KNITRO in terms of objective evaluations across various accuracy levels. Yet FD outperforms FIBO in terms of constraint evaluations across all accuracy levels. Since each constraint evaluation is associated with an evaluation of the quadratic model (4.2.1), FIBO is very effective when each evaluation of the objective function is at least as expensive as  $O(n^2)$  and the cost of



Figure 4.3. Log-ratio plot comparing FIBO and FD in terms of the number of constraint evaluations to satisfy (4.3.2) for  $\tau = 10^{-1}$  (upper left),  $10^{-3}$  (upper right),  $10^{-5}$  (bottom left),  $10^{-7}$  (bottom right).

each constraint evaluation is negligible. Otherwise, when the cost of solving the subproblem (4.2.4) is formidable, FD is a viable approach.

# 4.4. Final Remarks

We have demonstrated the effectiveness of the FIBO approach based on empirical investigations. When only analytical constraints are present, our study has revealed that, with minor modifications, existing unconstrained IBO methods can be extended to solve general constrained DFO problems. Our numerical results indicate that FIBO is a competitive method for problems where the cost of the objective function dominates that of the constraints. Our framework applies to any IBO methods and gradient-based nonlinear programming solvers. Yet more sophisticated algorithmic designs are required when the constraints cannot be violated, or black-box constraints are present. It is non-trivial to compute interpolation points that are feasible, when model-improving points are necessary to improve model quality. In addition, when constraints are expensive to evaluate, further algorithmic development for IBO is required to improve efficiency of the proposed method. We have observed fair performance of FIBO even if the termination test for (4.2.4) is relaxed, but the relaxation strategy requires more careful investigations.

# CHAPTER 5

# Prompting Large Language Models with Derivative-Free Optimization

### 5.1. Introduction

Pre-trained large language models (LLMs) like GPT-4 have been widely successful in various language tasks. However, LLMs are general-purpose models and require additional adaptations, such as fine-tuning or prompting, to address specific language tasks like sentiment analysis. Fine-tuning requires updates of the model parameters on each new language task [52]. This process can be resource-intensive and challenging given the model complexity and the limited sample size of the subsequent task. Additionally, maintaining separate copies of the language model for different tasks poses additional challenge in terms of storage requirements, especially given the massive size of the LLMs. On the other hand, prompting offers a storage-efficient solution by appending a prompt, e.g., task description texts, to the input text while keeping the parameters of the pre-trained model frozen. Prompting has demonstrated its effectiveness on LLMs with the appealing feature of using the same pre-trained model on a variety of tasks [15]. However, effective design of the prompts still remains a challenge: they are often designed in a manual trial-and-error manner and can be sensitive to perturbations.

Several efforts have been proposed to tackle this challenge, including gradient-based approaches that optimize continuous prompts, which are continuous vectors instead of discrete texts [60, 61]. However, accessing the gradient of the loss function through backpropagation in LLMs like GPT-4 is often not available, as they typically only allow input and output manipulations. Furthermore, the massive size of LLMs makes computing derivatives of the objective function through backpropagation impractical, primarily due to memory constraints. Moreover, LLMs are usually implemented using finite precision arithmetic, with common choices in half and single precision. This choice of precision can introduce higher roundoff errors during model evaluation compared to computations performed in double precision.

In light of this limitation, we approach prompt design as a noisy derivative-free optimization (DFO) problem. Several classes of DFO methods have gained popularity within different research communities, each offering unique advantages. Randomized methods such as evolution strategies have been widely embraced within the machine learning communities. This class of methods employs iterative sampling based on a stochastic distribution, which may be updated to leverage iteration information. Their effectiveness has been demonstrated on applications arising from Reinforcement Learning [98] to Natural Language Processing [100]. Alternatively, interpolation based optimization (IBO) methods and finite-difference (FD) based methods have been recognized as the state-of-the-art DFO algorithms within the optimization community for their robustness and efficiency [93]. These methods rely on utilizing function values to construct approximations to the objective function and/or gradients.

Considering the diverse range of DFO methods available, each with its unique advantages, we examine these approaches in solving the noisy DFO problem arising from prompt design.

#### 5.1.1. Literature Review

**Transformer-based Language Models**. Transformer-based language models have have been in the spotlight in recent years due to their impressive performance on a wide range of language tasks. The original transformer architecture was proposed in [101], which introduced the self-attention mechanism within an encoder-decoder structure, allowing the model to extract contextual information from the input sequence. The text-to-text transformer (T5) model uses the encoder-decoder structure [82]. Examples of model usage include question-answering. BERT and its variants use an encoder structure in the seq-2-seq architecture to extract general text representations [33]. They can be extended to applications such as sentiment classification. The GPT-2 [81] and GPT-3 [15] are decoder-based transformer models, so that they are trained in an autoregressive manner. The model can be used in, for instance, automatic sentence completion and text generation.

In practice, LLMs often have billions of parameters and thus require significant computational resources and memory, making gradient-based training infeasible for many users and applications. Furthermore, the majority of performant LLMs, such as GPT-4, and BARD, are not open-sourced. They are exclusively accessed through blackbox APIs that typically consume and generate natural language sequences. Due to the limited access to the models, leveraging them to address specific language tasks can be challenging. This calls for effective strategies such as prompt design.

**Prompting**. Prompts are pieces of information to be prepended to the input X so that the language model can generate a desirable output Y using the extra information. Prompting is desirable for large-size language models, for its advantage of not requiring updates or storage of all model parameters as in the fine-tuning paradigm, which can be

prohibitively expensive. Various prompting approaches have been proposed and yet it still remains an open question to achieve competitive performance against fine-tuning. See [62] for a thorough treatment.

Prompts can be classified into hard prompts and soft prompts. Hard prompts are often composed of task descriptions texts that rely on manual efforts. For instance, "I felt bored watching the movie." may be more likely to be classified as a negative sentence by a language model if we append the prompt "It was \_\_\_' to the sentence and ask the model to generate the word at the corresponding position. Yet hard prompts can be sub-optimal and sensitive to prompt choices [60]. In contrast, soft prompts are continuous vectors to be injected to the input. They do not necessarily use the same embedding as in the pre-trained language models. The vectors can be tuned or they can be treated as learnable parameters and be optimized via gradient-based methods through backpropagation [60, 61], if the gradient information of the pretrained LLM is available.

However, as mentioned in the previous sections, LLMs may only be accessed as blackboxes so that their gradient information or parameters are not accessible. As a result, [100] proposed Black-Box Tuning, which optimizes continuous prompts by employing the Covariance Matrix Adaptation Evolution Strategy (CMA-ES) method. Other blackbox approaches for prompt tuning include [32, 34], which employ discrete prompt optimization.

# 5.2. Problem Formulation

We consider a general classification task, which predicts a label  $y \in \mathcal{Y}$  for a given input sequence  $x = (x_1, ..., x_L)$  of length L using a general-purpose pre-trained LLM. To simplify the setting, we assume access to the embedding matrix E of the LLM and we denote  $X \in \mathbb{R}^{L \times D}$  as the embedded representation of the original input sequence x, where D is the embedding dimension. The embedded input X can be viewed as a numerical representation of the original text sequence via a linear projection using the embedding matrix.

Our goal is to search for a continuous prompt  $p \in \mathbb{R}^d$  to be combined with every embedded input  $X_i$  in the training samples by solving the following optimization problem:

(5.2.1) 
$$\min_{p \in \mathbb{R}^d} \frac{1}{N} \sum_{i=1}^N \mathcal{L}(f(p; X_i), y_i),$$

where  $\mathcal{L}$  is the loss function (e.g., cross entropy loss function), f is the black-box pre-trained language model and N is the number of training samples. The size of the prompt p is typically chosen as  $d = L_p \times D$ , where  $L_p$  is a user-provided hyperparameter. Note that the prompt  $p \in \mathbb{R}^d$  and the *i*-th embedded input  $X_i \in \mathbb{R}^{L_i \times D}$  may be passed to the rest of the LLM in the form of either  $\begin{bmatrix} p \\ X_i \end{bmatrix} \in \mathbb{R}^{(L_i+L_p) \times D}$ , or  $\hat{p} + X_i \in \mathbb{R}^{L_i \times D}$ , where  $\hat{p} \in \mathbb{R}^{L_i \times D}$  is usually generated from p using zero padding. We consider the latter form of prompt-input combination in our experiments, as it in general results in shorter input sequences lengths and thus less training time. In addition, p = 0 can be used as a natural starting point. Note that the same prompt is appended to every training sample.

The key challenge in solving (5.2.1) lies in the dimensionality of the problem. Since the embedding dimension (D) is often very high, the size of prompt vector p is on a scale of thousands. Yet it has been empirically shown that LLMs have low intrinsic dimensionality. In other words, by tuning only hundreds of parameters and performing random projection onto the full parameter space, one make achieve reasonably good performance on downstream NLP tasks [1]. Therefore, (5.2.1) can be reformulated as

(5.2.2) 
$$\min_{z \in \mathbb{R}^n} \frac{1}{N} \sum_{i=1}^N \mathcal{L}(f(Az; X_i), y_i)$$

where  $A \in \mathbb{R}^{d \times n}$  is a random projection matrix such that  $n \ll d$ .

Recall that the gradient of the function f is not available or very expensive to compute. Thus (5.2.1) and (5.2.2) are unconstrained derivative-free optimization problems.

Furthermore, the problem formulation (5.2.2) is noisy in practice as LLMs are typically implemented in single precision, half precision or mixed precision due to GPU memory constraints and speed concerns. Therefore, DFO algorithms such as finite differences must be applied carefully when such deterministic noise is present in the function evaluations.

We remark that our approach assumes that the continuous projected prompt Az in (5.2.2) can be combined with the embedding representation X of the input sequence x directly. However, there are cases when a given LLM only accepts text inputs and does not provide access to its embeddings. In such scenarios, the prompt must be prepended to the input text sequence before passing through the embedding. This renders the prompt p discrete and (5.2.1) intractable to solve. We argue that it can be beneficial for LLM services to include a feature that enables seamless integration of continuous prompts into the embeddings. This would simplify the handling of hard prompts, ultimately benefiting a wide range of language tasks and workflows.

# 5.3. Numerical Experiments for Optimizing Prompts

We evaluate the effectiveness of our proposed approach based on transformer-based language models RoBERTa and T5, in addressing several common language classification tasks, including sentiment classification, natural language inference and topic classification. To tackle these tasks, we examine three state-of-the-art classes of DFO algorithms: IBO methods, finite differences and randomized algorithms for solving the problem (5.2.2). We describe below our experiments details.

#### 5.3.1. Experiment Setup

**Datasets**. We evaluate on 6 common language classification datasets: SST-2 [95], Yelp P. [106], SNLI [12], RTE [102], DBpedia [3], AG's News [106]. The datasets include tasks such as sentiment analysis, topic classification and natural language inferences.

**Few-Shot Learning**. It has been demonstrated that LLMs have great power for few-shot learning [15]. To simulate the scarce data setting, we consider a 16-shot learning setting in our experiments. In particular, we randomly choose 16 samples for each class to construct the training and validation set. As in [100], we use the original test set for consistency with existing benchmarks.

**Backbone Model**. In our experiments, we consider open-source transformer-based models RoBERTa-large [63] and T5-large [82] for solving the language tasks so that the prompts can be directly applied to the embedded representations. The chosen models are representative of bi-directional encoder and encoder-decoder models respectively.

Hyperparameters. We employ cross entropy loss as our loss function  $\mathcal{L}$  and use accuracy to measure performance. We consider the decision variable z of dimension 500 (n = 500) along with prompt length  $L_p = 50$ . We employ a projection matrix A that is generated from a Gaussian distribution with mean and covariate matrix computed using the mean and standard deviation of the language model embeddings. Note that the projection matrix A remains fixed during the experiments. We set a budget of 20,000 function evaluations in our numerical experiments.

**DFO Solvers**. We describe below the implementation details of the DFO algorithms.

- CMA-ES: A stochastic DFO algorithm developed by [45]. We call an Python implementation of CMA-ES via the pycma module with version 3.3.0. The parameters are set to default and no additional information is provided regarding noise. Specifically, the method samples  $4 + 3\sqrt{n}$  points at every iteration.
- FD L-BFGS: A finite-difference based implementation of the limited-memory BFGS algorithm with a bisection Armijo-Wolfe linesearch. The memory size is set to 10. We used forward differencing to compute a finite difference approximation of the gradient and the directional derivative along the search direction during linesearch. We employed differencing interval h of size 0.1 based on machine epsilon for single precision and an estimated bound on the second order derivative of the loss function.
- NEWUOA: A state-of-the-art unconstrained IBO method that is available through the PDFO package [84]. We set **rhoend** =  $10^{-8}$  as a termination criterion for this algorithm. All other parameters are set to default.

We use z = 0 as the starting point for all algorithms.

# 5.3.2. Results

We comment on the training loss of the prompt optimization problem (5.2.2) across different datasets. Our observations reveal interesting insights into the performance of different solvers. Overall we observe that CMA-ES quickly reduces the training loss within *n* function evaluations. On the other hand, FD LBFGS requires n + 1 function evaluations to approximate a gradient and NEWUOA requires 2n + 1 function evaluations for quadratic model construction during the first iteration. As the number of model evaluation increases, we notice that NEWUOA becomes more efficient in minimizing the training loss compared to the other solvers. When a sufficiently large number of model evaluations is allowed, CMA-ES may obtain the lowest training loss among the solvers we have compared. We also observe that FD L-BFGS tends to be the least efficient solvers among the solvers. To illustrate the relative performance across the solvers, we plot representative training loss and validation loss in Figure 5.1 and Figure 5.2.



Figure 5.1. Training and Validation Loss for Prompt Optimization of RoBERTa-large for Solving SST-2.

Dataset	SST-2	Yelp P.	AG's News	DBPedia	SNLI	RTE
CMA-ES	89.79	85.82	83.89	85.65	45.04	42.24
NEWUOA	88.30	91.62	83.11	95.78	44.51	42.24
FD L-BFGS	86.93	91.85	81.68	82.70	47.94	53.07
Manual Prompt	83.60	89.64	75.75	31.92	38.82	51.43

Table 5.1. Test accuracy: Prompt Optimization for RoBERTa-large.

Moreover, we present the final test accuracy based on the model that corresponds to the lowest validation loss. The results are summarized in Table 5.1 and Table 5.2



Figure 5.2. Training and Validation Loss for Prompt Optimization of T5large for Solving DBPedia.

Dataset	SST-2	Yelp P.	AG's News	DBPedia	SNLI	RTE
CMA-ES	86.70	95.91	83.51	93.25	46.62	46.93
NEWUOA	86.01	93.97	77.74	89.46	42.59	56.32
FD L-BFGS	91.28	94.92	77.64	90.26	40.17	53.79
Manual Prompt	69.04	82.24	40.07	59.40	40.18	54.51

Table 5.2. Test accuracy: Prompt Optimization for T5-large.

for RoBERTa-large and T5-large, respectively. We also include results for a manually tuned prompt, which in general yields lower accuracy compared to the final test accuracy obtained by the DFO solvers. Our numerical experiments indicate that NEWUOA exhibit comparatively lower test accuracy across the evaluated datasets despite its efficiency in reducing the training loss. For RoBERTa-large, neither FD L-BFGS nor CMA-ES emerges as a clear winner. In contrast, when T5-large is used as the backbone model, CMA-ES achieves the highest test accuracy across majority of the datasets.

# 5.4. Final Remarks

We have addressed an important problem of adapting large language models to various language classification tasks by leveraging DFO algorithms. We have demonstrated the
efficacy of several DFO algorithms, including CMA-ES, NEWUOA and FD L-BFGS in solving this DFO problem. Our numerical study has highlighted the efficiency of NEWUOA in effectively minimizing the training loss. Additionally, we have observed the potential of CMA-ES in achieving the lowest training loss when a sufficiently large number of model evaluations are employed. Nevertheless, further investigation is needed to assess the generalization of DFO methods on the testing set, as well as a more sophisticated formulation of the DFO problem. We also note that the unavailability of embedding information for some existing LLMs poses challenges for this prompt design approach. Additional investigations can be conducted to study the feasibility of utilizing external embeddings for this prompt design problem and to provide insights into the applicability of DFO methods in this setting.

# CHAPTER 6

# Analyzing the Performance of DFO Methods on a Wider Class of Problems

### 6.1. Introduction

Motivated by the observations in Chapter 5, we delve deeper into the comparative analysis of DFO methods by exploring more generalized problem settings. We are interested in whether the relative performance of the compared DFO solvers, concerning the objective presented in problem (5.2.2), extends to other functions in general. We aim to gain further insights into the relative merits and performance characteristics of these DFO methods across a broader range of problem scenarios, and identify potential avenues for improving their performance in practice. In particular, we explore strategies for improving the accuracy of IBO methods in the presence of noise.

Formally, we now consider general unconstrained DFO problems of the following form as in (1.0.1) (assuming  $\Omega = \mathbb{R}^n$ )

(6.1.1) 
$$\min_{x \in \mathbb{R}^n} \phi(x)$$

where  $\phi : \mathbb{R}^n \to \mathbb{R}$  is a smooth function. Recall from Chapter 1 that the derivatives of  $\phi$  are not available and only noisy evaluations of the function can be observed:

(6.1.2) 
$$f(x) = \phi(x) + \epsilon(x),$$

where  $\epsilon(x)$  models the evaluation error at x. We assume that the noise  $\epsilon(x)$  is stochastic and independent of x. See Chapter 1 for further discussions on the noise model.

As has been well-studied in [93], NEWUOA and FD-LBFGS are competitive for noiseless problems, with NEWUOA showcasing a slight advantage over FD L-BFGS for noisy problems. Consequently, we narrow down our empirical investigations to NEWUOA and CMA-ES for solving unconstrained DFO problems. Therefore, the solvers tested in this chapter are as follows.

CMA-ES. We call an Python implementation of CMA-ES via the pycma module with version 3.3.0. The parameters are set to default and no additional information is provided with respect to noise. To elaborate, the method samples  $4 + 3\sqrt{n}$  points at every iteration.

NEWUOA. We use NEWUOA through a Python wrapper PDFO [84]. We set rhoend =  $10^{-8}$  as a termination criterion for this algorithm. All other parameters are set to default, e.g., the number of interpolation point is set to 2n + 1.

To summarize the per-iteration cost of the methods, NEWUOA incurs  $O(n^2)$  flops of computation cost at every iteration due to interpolation, while CMA-ES requires  $O(n^3)$  flops due to eigendecomposition of the covariance matrix. The computation cost of CMA-ES can be reduced to  $O(n^2)$  when the decomposition is performed at every n/10 iterations [48].

We will investigate the above methods on both noiseless and noisy functions in Section 6.2 to study the effect of noise on their performance.

#### 6.2. Numerical Experiments

We selected 73 unconstrained CUTEst problems with varying dimensions up to 200 to test the performance of NEWUOA and CMA-ES. We compare them on problems with and without noise to simulate different settings of DFO problems.

Several metrics have been proposed in the literature to study the relative performance of different methods, including performance profiles, data profiles and the log-ratio profiles [67, 69]. We choose the log-ratio profile [67] in our study, as in contrast to the other two options, it provides both an overview of the general performance of the methods and their relative performance for each specific problem. Specifically, the height of the plotted bars and the area of the shaded region can offer insights into the performance of the algorithms with relative scale, although it works the best when only two methods are being compared.

We present our empirical findings based on accuracy and efficiency of the algorithms.

#### 6.2.1. Noiseless Functions

We first present the case when the functions do not contain noise, i.e., the true objective is always returned in double precision.

Accuracy. To measure the accuracy of the final solutions obtained by CMA-ES and NEWUOA, we vary the budget, i.e. total number of function evaluations allowed, from 0.5n to 500n and compare the optimality gap achieved by them. Formally, we compare the relative accuracy of CMA-ES and NEWUOA up to 8 digits by computing the log-ratio:

(6.2.1) 
$$\log_2 \frac{\max(\phi_{\text{CMA-ES}} - \phi^*, 10^{-8})}{\max(\phi_{\text{NEWUOA}} - \phi^*, 10^{-8})}$$

where  $\phi_{\text{CMA-ES}}$  and  $\phi_{\text{NEWUOA}}$  denote the best objectives achieved by the corresponding algorithms within the given budget, and  $\phi^*$  denotes the optimal objective value. We plot (6.2.1) in an increasing order so that the areas of the shaded regions give a general overview of the relative performance of the compared methods, with each bar demonstrating the relative accuracy for a given problem.



Figure 6.1. (Selected) Accuracy Log-Ratio Profiles for Noiseless Problems. Plots of (6.2.1) comparing CMA-ES and NEWUOA within a budget of 0.5n (upper left), 2n (upper right), 10n (bottom left) and 500n (bottom right).

As observed in Figure 6.1, CMA-ES can achieve higher accuracy than NEWUOA when the allowed number of function evaluations is limited to less than 2n. This behavior aligns with our expectations since NEWUOA requires 2n + 1 function evaluations to perform the first iteration. However, as the allowed budget for function evaluations increases, NEWUOA

outperforms CMA-ES on nearly all problems. Nevertheless, it is worth noting that when a significantly larger budget of 500n function evaluations is allowed, the performance gap between NEWUOA and CMA-ES is narrowed. This indicates that CMA-ES can require a considerate number of function evaluations to reach the solution.

Efficiency. We can further study the relative performance of CMA-ES and NEWUOA by checking the minimum number of function evaluations required by each method to satisfy

(6.2.2) 
$$\phi(x_k) - \phi^* \le \tau \cdot \max(1, |\phi^*|)$$

where  $\phi^*$  is the optimal objective value and  $\tau$  is a pre-specified accuracy level. We denote the required number of function evaluations as  $eval_{CMA-ES}$  and  $eval_{NEWUOA}$ . We fix the budget as 500*n* function evaluations and vary  $\tau$  from 0.1 to  $10^{-8}$  to study the performance with respect to various accuracy levels and plot the log-ratio

(6.2.3) 
$$\log_2 \frac{\operatorname{eval}_{\mathrm{CMA-ES}}}{\operatorname{eval}_{\mathrm{NEWUOA}}}$$

in an increasing order in Figure 6.2.

We observe that NEWUOA is much more efficient than CMA-ES in achieving (6.2.2) across nearly all problems, regardless of the accuracy level. Therefore, when dealing with noiseless DFO problems where attaining a specific accuracy level is deemed sufficient or when the budget is sufficiently large for model construction, NEWUOA is more preferable. Otherwise, if only fewer than 2n function evaluations are available, CMA-ES can be served as an alternative for potential higher accuracy.



Figure 6.2. (Selected) Efficiency Log-Ratio Profiles for Noiseless Problems. Plots of (6.2.3) comparing CMA-ES and NEWUOA with  $\tau = 0.1$  (upper left),  $10^{-4}$  (upper right),  $10^{-6}$  (bottom left) and  $10^{-8}$  (bottom right).

# 6.2.2. Noisy Functions

We now present results for problems with stochastic noise. To simulate a noisy setting, we inject i.i.d Gaussian noise with mean zero and variance  $\epsilon_f^2$  to the 73 noiseless CUTEst functions:

(6.2.4) 
$$f(x;\xi) = \phi(x) + \epsilon(x;\xi), \quad \epsilon(x;\xi) \sim \mathcal{N}(0,\epsilon_f^2),$$

where  $\epsilon_f$  is denoted as the noise level of the problem.

Accuracy. To investigate the robustness of both CMA-ES and NEWUOA against computation errors, we ran the solvers using the same parameter settings as in the noiseless case. Similar to Section 6.2.1, we measure the best true objective value obtained by both methods within budgets ranging from 0.5n to 500n function evaluations. We report the log-ratio profiles in Figure 6.3 for a noise level of  $10^{-3}$ , which is representative for other noise levels. The log-ratio profiles reveal that for very small budgets ( $\leq 2n$ ), CMA-ES is slightly more competitive than NEWUOA. However, as the budget increases to medium level, NEWUOA consistently outperforms CMA-ES in terms of final accuracy. Intriguingly, when more generous budgets of function evaluations are allowed, CMA-ES demonstrates its ability to achieve significantly higher accuracy than NEWUOA across a majority of the problems.



Figure 6.3. (Selected) Accuracy Log-Ratio Profiles for Noisy Problems ( $\epsilon_f = 10^{-3}$ ). Plots of (6.2.1) comparing CMA-ES and NEWUOA within a budget of 0.5n (upper left), 2n (upper right), 10n (bottom left) and 500n (bottom right).

Efficiency. We evaluate the efficiency of CMA-ES and NEWUOA in the noisy setting by comparing the number of function evaluations, denoted as  $eval_{CMA-ES} eval_{NEWUOA}$ , that are required to achieve

(6.2.5) 
$$\phi(x_k) - \phi^* \le \tau \cdot (\phi(x_0) - \phi^*),$$

where  $\tau$  varies from 0.1 to  $10^{-8}$  and  $\phi^*$  denotes the optimal objective value. We plot the log-ratio (6.2.3) in Figure 6.4 and 6.5 to report the efficiency of the methods corresponding to noise levels of  $10^{-1}$  and  $10^{-7}$  respectively. Our observations indicate that NEWUOA is more efficient than CMA-ES in obtaining a given accuracy level when the desired accuracy level is low or the noise level is low. For highly noisy problems, CMA-ES and NEWUOA are competitive in achieving high accuracy levels. This finding highlights the advantage of CMA-ES when an abundant number of function evaluations are available for solving noisy DFO problems, especially when high accuracy is desirable.

We note that, compared to the noiseless case, the efficiency of NEWUOA in solving functions with high noise levels and the final accuracy obtained by it under large budgets are both impacted by noise. Although NEWUOA is more competitive than CMA-ES in solving noiseless problems, with the presence of noise, CMA-ES is more capable of achieving high accuracy especially for large noise levels. This implies that the presence of noise poses additional challenges for NEWUOA, impacting its effectiveness in solving unconstrained DFO problems. In light of this finding, we briefly discuss in the next section the effect of noise on the performance of an IBO method and potential strategies to mitigate the effects of noise on the performance of NEWUOA and other IBO methods.



Figure 6.4. (Selected) Efficiency Log-Ratio Profiles for Noisy Problems ( $\epsilon_f = 0.1$ ). Plots of (6.2.3) comparing CMA-ES and NEWUOA with  $\tau = 0.1$  (upper left),  $10^{-4}$  (upper right),  $10^{-6}$  (bottom left) and  $10^{-8}$  (bottom right).

## 6.3. On the Accuracy of IBO methods in the Presence of Noise

Recall from Chapter 1 that IBO methods rely on model construction and ratio tests to perform updates within a trust region framework. With the presence of noise in the objective function, the quadratic model can be inaccurate due to the noisy function values of the interpolation points. Moreover, the noise can cause the ratio test to suggest misleading suggestions for step acceptance or rejection, as the noise can lead to arbitrarily wrong ratio values. Despite these challenges, our numerical results suggest that NEWUOA still demonstrates reasonable progress even when noise is present. To shed light on this observation, we present below a result that provides an upper bound for the approximation



Figure 6.5. (Selected) Efficiency Log-Ratio Profiles for Noisy Problems ( $\epsilon_f = 10^{-7}$ ). Plots of (6.2.3) comparing CMA-ES and NEWUOA with  $\tau = 0.1$  (upper left),  $10^{-2}$  (upper right),  $10^{-4}$  (bottom left) and  $10^{-6}$  (bottom right).

error of the underdetermined quadratic model m(x), when bounded noise is present. See Appendix D for a proof based on [29]. A similar bound can be obtained by considering the mean squared approximation error.

**Theorem 6.3.1.** Suppose  $|\epsilon(x)| \leq \epsilon_f, \forall x \in \mathbb{R}^n \text{ and } \phi : \mathbb{R}^n \to \mathbb{R}$  is continuously differentiable with  $\nabla \phi$  being  $L_2$ -Lipschitz continuous. We further assume that the interpolation set  $Y = \{y^0, y^1, ..., y^p\}$  is poised in  $B(y^0; \Delta)$ . Then for all  $x \in B(y^0; \Delta)$ , we have that

(6.3.1) 
$$\|\nabla\phi(x) - \nabla m(x)\| \le \frac{5\sqrt{p}}{2} \|\hat{M}^{\dagger}\| (L_2 + \|H\|)\Delta + \frac{2\sqrt{p} \|\hat{M}^{\dagger}\|\epsilon_f}{\Delta}$$

(6.3.2) 
$$|\phi(x) - m(x)| \le \frac{5\sqrt{p} \|M^{\dagger}\| + 1}{2} (\|H\| + L_2) \Delta^2 + (2\sqrt{p} \|\hat{M}^{\dagger}\| + 1)\epsilon_f$$

where  $\hat{M}^{\dagger} = (\hat{M}^T \hat{M})^{-1} \hat{M}^T$  and

$$\hat{M} = \begin{bmatrix} (y^1 - y^0)/\Delta & (y^2 - y^0)/\Delta & \dots & (y^p - y^0)/\Delta \end{bmatrix}$$

Note that in contrast to the noiseless case when  $\epsilon_f = 0$ , for which the upperbound in (6.3.1) and (6.3.2) shrink to zero, instead the upperbound for (6.3.1) tends to infinity due to the additional noise term. Therefore, when the trust region is sufficiently large and the noise term is dominated by the other term, one can expect IBO methods to make progress similarly as in the noiseless case. However, one cannot expect IBO methods to progress effectively once the trust region becomes too small. Indeed, by minimizing the right-hand-side of (6.3.1), the trust region that minimizes the approximation is  $\underline{\Delta} = 2\sqrt{\frac{\epsilon_f}{5(||H||+L_2)}}$ .

We present in Figure 6.6 an example of NEWUOA failing to make progress when the initial trust region is too small. We observe that when noise is not present in the function, NEWUOA is able to achieve similar accuracy level no matter what value the initial trust region is set to, although efficiency of the algorithm may differ. However, in the presence of noise, NEWUOA may stagnate when the trust region is too small with respect to the noise level (and potentially other characteristics of the function and interpolation set), with the algorithm keeps shrinking the trust region.

It has been found in [93] that using more interpolation points can be beneficial in increasing the final accuracy of NEWUOA. However, as there is an upper limit in the number of interpolations in quadratic model construction (we do not consider regression model), we briefly discuss three strategies that can improve the accuracy of an IBO method in the presence of noise.



Figure 6.6. Solving noiseless and noisy SROSENBR function with 10 variables using different initial trust region  $\Delta_0$ . Left: noiseless. Right: noise level = 0.1.

**Restarts**. A simple yet effective strategy of achieving higher accuracy for NEWUOA is to employ restarts. To elaborate, instead of terminate NEWUOA when the trust region converges to the minimum size  $(10^{-8} \text{ in our setting})$ , we restart the algorithm at the returned solution by considering it as the new initial point. Therefore, we abandon every point in the interpolation set and reset the trust region radius to the initial value. We repeat this process until we have reached the allowable budget, i.e., the maximum number of function evaluations. We have observed effectiveness of this approach in our empirical studies and we visualize the findings in the log-ratio plots in Figure 6.7 and Figure 6.8.

However, this naive approach can be deemed inefficient as in Figure 6.9. Note that by increasing the trust region to large size, the algorithm evaluates at a great number of points that have high function values and are far from the solution. We speculate that more effective restart procedure can be developed such that the function evaluations are used more efficiently. For instance, by using a smaller initial trust region for later runs or by setting the minimum trust region to be on the order of the noise level allows more



Figure 6.7. (Selected) Accuracy Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.1) comparing CMA-ES and NEWUOA with restarts within a budget of 500n.



Figure 6.8. (Selected) Efficiency Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.3) comparing CMA-ES and NEWUOA with restarts for  $\tau = 10^{-8}$ .

restarts and potentially higher accuracy. See [18] for a restart strategy for noisy least square problems.

**Ratio Test Relaxation**. As suggested in [99], applying relaxation to the ratio test can mitigate the effect of noise in the objective, particularly to prevent the trust region from becoming excessively small. Therefore, we modify NEWUOA such that the ratio test is



Figure 6.9. Solving DIXMAANJ with 90 variables using NEWUOA with restarts and CMA-ES. The noise level is  $10^{-3}$ .

relaxed as

(6.3.3) 
$$\frac{f(x_k) - f(x_k + s) + 2\epsilon_f}{m_k(x_k) - m_k(x_k + s) + 2\epsilon_f}$$

However, contrary to our expectations, our numerical experiments demonstrate that relaxing the ratio test for NEWUOA does not improve the final accuracy overall. We present the accuracy log-ratio plots in Figure 6.10. Interestingly, we have observed accuracy improvements in DFOTR, a different IBO algorithm, when performing the same ratio relaxation. This outcome may be attributed to the utilization of a trust region lower bound in NEWUOA, which prevents the algorithm from accepting steps that are too small with respect to this bounds. Indeed, as shown in Figure 6.11, NEWUOA is unable to recover from an initial trust region is too small even though the ratio test has been relaxed. This observation indicates that more sophisticated design may be required to further improve the robustness of NEWUOA for solving noisy DFO problems.



Figure 6.10. (Selected) Accuracy Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.1) comparing NEWUOA with and without restarts using a budget of 500n.



Figure 6.11. Solving noisy SROSENBR function with 10 variables and noise level of 0.1 using different initial trust region  $\Delta_0$ . The ratio test in NEWUOA is relaxed as in (6.3.3).

Adaptive Sampling. When the noise term  $\epsilon(x)$  in the problem (6.1.2) is stochastic, another viable approach for obtaining an higher accuracy is to employ sample averaging so that the function value at each point is computed as an average of multiple evaluations at the same point [20, 31, 89]. The noise level of the function is therefore reduced due to variance reduction. An adaptive sample averaging strategy can balance the tradeoff between the efficiency of the IBO method and the final accuracy. However, theoretical convergence of existing adaptive sample averaging based IBO methods often require a sample size of  $O(\Delta_k^{-4})$ . This is not desirable even though in practice the sample size is chosen as  $O(\Delta_k^{-1})$ , as it leads to inefficiency in early iterations. Since IBO methods can often make progress before the trust region become too small, a more sophisticated sampling strategy may be available. For instance, the sample size may only need to be increased when the trust region reaches a given level, say the trust region size that minimizes the upper bound of (6.3.1). Yet in our preliminary studies we have observed difficulties in the algorithm in making progress when the trust region is not adequately enlarged and the interpolation set is not adjusted. We speculate that an optimal adaptive sampling strategy requires careful handling of factors such as the sample size, trust region management and positioning of the interpolation sets.

#### 6.4. Final Remarks

We have performed a comparative study of NEWUOA and CMA-ES on a broad set of unconstrained problems in this chapter. We have observed that NEWUOA outperforms CMA-ES in terms of efficiency and accuracy when the function evaluations are exact. Notably, our findings indicate that, in the presence of noise, NEWUOA is still more efficient than CMA-ES for obtaining desirable accuracy levels within a moderate number of function evaluations. Yet as opposed to the noiseless case, it obtains lower final accuracy than CMA-ES when a large budget of function evaluations are available.

This performance gap motivates us to analyze the challenges associated with noisy evaluations for IBO methods and to understand potential strategies for improving their performance in the presence of noise. In particular, our preliminary empirical results suggest that utilizing naive restarts with NEWUOA can significantly narrow down the performance gap between CMA-ES in terms of final accuracy. However, relaxing the ratio test does not lead to improved accuracy of NEWUOA as opposed to other IBO methods such as DFO-TR. In addition, the effectiveness of adaptive sampling is contingent upon further investigation. These findings highlight the importance of designing more efficient restarts for improving the accuracy, the limited impact of relaxing the ratio tests, the need for more meticulous design of adaptive sampling, and the necessity of further redesign of the classical IBO methods in the presence of noise.

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# APPENDIX A

# On the Numerical Performance of Finite-Difference Based Methods for Derivative-Free Optimization

A.1. Numerical Investigation of Lipschitz Estimation

# A.1.1. Investigation of Theoretical Lipschitz Estimates

In Section 2.2.2, we approximated the bound on the second and third derivative Land M (or the Lipschitz constant of the first and second derivative) along the interval  $I = \{x \pm tp : t \in [0, h_0]\}$  for  $h_0 > 0$  every time finite-differencing is performed.

For forward-differencing, we employed the Moré and Wild heuristic [71]. The heuristic estimates the bound on the second derivative of a univariate function with noise. Assume  $\phi : \mathbb{R} \to \mathbb{R}$  is univariate. If we let  $\Delta(t) = f(x+t) - 2f(x) + f(x-t)$ , then t > 0 must satisfy two conditions:

(A.1.1) 
$$|\Delta(t)| \ge \tau_1 \epsilon_f, \qquad \tau_1 \gg 1$$

(A.1.2) 
$$|f(x \pm t) - f(x)| \le \tau_2 \max\{|f(x)|, |f(x \pm t)|\}, \quad \tau_2 \in (0, 1)$$

In practice,  $\tau_1 = 100$  and  $\tau_2 = 0.1$ . The first condition ensures that h is sufficiently large such that the second-order difference is not dominated by noise, while the second condition enforces that the difference is not dominated by a particular evaluation, so that there is "equal" contribution from each function evaluation in the finite-difference approximation. If conditions (A.1.1) and (A.1.2) are satisfied, then we take  $L = \max\{10^{-1}, |\Delta(t)|/t^2\}$ .

For central differencing, we used a theoretical estimate based on knowledge of the true Hessian  $\nabla^2 \phi(x)$ :

(A.1.3) 
$$M = \max\left\{10^{-1}, \frac{|p^T(\nabla^2 \phi(x + \tilde{h}\frac{p}{\|p\|}) - \nabla^2 \phi(x))p|}{\tilde{h}\|p\|^2}\right\}$$

where  $\tilde{h} = \sqrt{\epsilon_M}$ . Since this estimate is ideal, we do not include the cost of evaluating M in the number of function evaluations. In both cases, the maximum is taken to ensure that the bound is strictly bounded away from 0 so that (2.2.9) and (2.2.10) remain well-defined.

One may ask whether or not a refined choice of L or M is necessary for finite-difference L-BFGS. To show the impact of Lipschitz estimation on the performance of finite-difference methods, we focus on forward-difference L-BFGS and compare nine different theoretical Lipschitz estimation schemes. The first five consider techniques where L is fixed for the entire run based on information at the initial point. We test this because it is frequently claimed that using an initial estimate of L is sufficient for the entire run; see [40, 71]. The first approach requires no additional information about the function, while the others incorporate information about the Hessian at the current point. The latter four techniques similarly incorporate information about the Hessian but re-estimate L whenever the finite-difference gradient or directional derivative is evaluated.

All of these techniques rely on the assumption that  $|D_p^2\phi(x)| \approx |D_p^2\phi(\xi)|$  for some  $\xi \in [x, x + h]$ . As we will see, both of these approximations that are based on conventional wisdom are challenged in our experiments.

- (1) Fix L = 1. This requires no additional knowledge about the problem at no added cost.
- (2) Fix  $L = \max\{10^{-1}, \|\nabla^2 \phi(x_0)\|_2\}$ . Similar to fixing L = 1, but incorporates knowledge of the initial Hessian.
- (3) Fix  $L = \max\left\{10^{-1}, \frac{1}{n}\sum_{i=1}^{n} |[\nabla^2 \phi(x_0)]_{ii}|\right\}$ . This can be obtained by choosing h such that the bound on  $||g(x_0) \nabla \phi(x_0)||_1$  at the initial point is minimized.
- (4) Fix  $L = \max\left\{10^{-1}, \frac{1}{\sqrt{n}}\sqrt{\sum_{i=1}^{n} [\nabla^2 \phi(x_0)]_{ii}^2}\right\}$ . This can be obtained by choosing h such that the bound on  $\|g(x_0) \nabla \phi(x_0)\|_2^2$  at the initial point is minimized.
- (5) Fix L to a vector  $(\max\{10^{-1}, |[\nabla^2 \phi(x_0)]_{ii}|\})_{i=1}^n$  and use the *i*th component for estimating the *i*th component of g(x). Uses  $||L||_2/\sqrt{n}$  when estimating h for the directional derivative in the line search.
- (6) Evaluate  $L = \max\{10^{-1}, \|\nabla^2 \phi(x)\|_2\}$  each time finite-differencing is performed.
- (7) Evaluate  $L = \max\left\{10^{-1}, \frac{1}{n}\sum_{i=1}^{n} |[\nabla^2 \phi(x)]_{ii}|\right\}$  each time finite-differencing is performed.
- (8) Evaluate  $L = \max\left\{10^{-1}, \frac{1}{\sqrt{n}}\sqrt{\sum_{i=1}^{n} [\nabla^2 \phi(x)]_{ii}^2}\right\}$  each time finite-differencing is performed.
- (9) Evaluate  $L = \max\{10^{-1}, |[\nabla^2 \phi(x)]_{ii}|\}$  when evaluating  $[g(x)]_i$ . When the directional derivative along  $p \in \mathbb{R}^n$  is computed, evaluates  $L = \max\{10^{-1}, \frac{|p^T \nabla^2 \phi(x)p|}{||p||_2^2}\}$

Each method is terminated when it cannot make more progress over 5 consecutive iterations. The optimal value  $\phi^*$  is obtained by running L-BFGS on the non-noisy function until no more progress can be made.

To compare the solution quality between two algorithms, we will use log-ratio profiles as proposed in [67] over the optimality gaps for each algorithm. The log-ratio profiles report

(A.1.4) 
$$\log_2\left(\frac{\phi_{\text{new}} - \phi^*}{\phi_{\text{old}} - \phi^*}\right)$$

for each problem plotted in increasing order. The area of the shaded region is representative of the general success of the algorithm.

One may ask whether or not using L = 1 is sufficient, without any additional knowledge of the second derivative. As seen in Figure A.1, when compared against the other fixed Lipschitz estimation schemes, the difference is not significant. In particular, the additional information from the Hessian does not yield significant benefits over setting L = 1 because the bound on the second derivative does not remain valid over the entire run of the algorithm. This may be surprising as it has been commonly suggested that it is sufficient to fix the Lipschitz estimate; see [40, 71]. In fact, we will see that it is more crucial for the algorithm to have a right estimate of L during the later stages of the run than at the beginning of the run in order to achieve high accuracy.

To further support why an adaptive L is important to achieve high accuracy, we compare the fixed L approaches against the adaptive L in Figure A.2.

As seen in Figure A.2, we see that using Lipschitz estimation at every iteration yields much higher accuracy than the alternatives. Upon inspection of individual runs, one can observe many cases where fixing the Lipschitz constant is unstable and inadequate, potentially leading to poor gradient approximations that result in early stagnation of the algorithm. This suggests that it is imperative to adaptively re-estimate L in order to



Figure A.1. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing forward difference L-BFGS with L = 1 and other fixed Lipschitz estimation schemes. The noise level is  $\sigma_f = 10^{-3}$ , but is representative for  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ .

reduce the noise in the gradient and squeeze the best possible accuracy out of forward difference L-BFGS.

In addition, since estimating L for each direction is most competitive out of the adaptive variants as seen in Figure A.3, we present the results for a heuristic approach in our main work. However, our investigation reveals that the simple mean or root mean square can provide potentially cheaper alternatives if they can be estimated without knowledge of the componentwise Lipschitz constants.

To our knowledge, there are two weaknesses with the componentwise estimation approach. The first is that for a small subset of problems, it is prone to underestimate the Lipschitz constant, as discussed in Section 2.2. The second weakness is that the approach, if performed at each iteration and approximated through finite-differencing, is far too



Figure A.2. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing forward difference L-BFGS with fixed and adaptive Lipschitz estimation schemes. The noise level is  $\sigma_f = 10^{-3}$ , but is representative for  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ .



Figure A.3. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing forward difference L-BFGS with fixed and adaptive Lipschitz estimation schemes. The noise level is  $\sigma_f = 10^{-3}$ , but is representative for  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ .

expensive. We propose a few possible practical heuristics for handling this in the following section.

## A.1.2. Practical Heuristics for Lipschitz Estimation

Two heuristics were proposed for estimating the Lipschitz constant. Moré and Wild [71] proposed a simple heuristic for estimating the bound on the second derivative of a univariate function with noise, as described in Section 2.2. Gill, et al. [40] proposed a similar procedure that instead enforces that the relative cancellation error lies within an interval

(A.1.5) 
$$\frac{4\epsilon_f}{|\Delta(h)|} \in [0.001, 0.1].$$

Note that the lower bound on the interval in (A.1.5) corresponds to (A.1.1).

Both of these heuristics were proposed with particular initial guesses of h and additional conditions to handle certain cases where the methods can fail. These heuristics were employed only at the beginning of the iteration for each variable component, then remains fixed for the entire run.

The only work to our knowledge that employs re-estimation of the Lipschitz constant for finite-differencing derivative-free optimization is Berahas, et al. [7]. In his work, the Lipschitz constant is re-estimated whenever the line search fails and the recovery procedure is triggered. Similarly, we will use the Moré and Wild heuristic to estimate the Lipschitz constant, but only re-estimate the Lipschitz constant when  $\alpha < 0.5$  after the first iteration, as described in Procedure I in Section 2.2. We compare two variants of the Lipschitz estimation procedure:

(1) Component MW: We use the Moré and Wild heuristic to estimate the Lipschitz constant with respect to each component. This is performed at the first iteration

and subsequent iterations for line search methods when the line search from the prior iteration gives a steplength  $\alpha_k < 0.5$ . When estimating the directional derivative along the search direction  $p_k$ , we use the root mean square of the component-wise estimates.

(2) Random MW: We use the Moré and Wild heuristic to estimate the Lipschitz constant along a random direction sampled uniformly from a sphere, i.e.,  $p \sim S(0, I)$ . This is performed at the first iteration and subsequent iterations for line search methods when the line search from the prior iteration gives a steplength  $\alpha_k < 0.5$ .

We compare both Moré and Wild heuristics against NEWUOA and theoretical componentwise Lipschitz estimation in Figures A.4, A.5, and A.6. In general, Component MW gives a better solution than Random MW. When we compare these methods against NEWUOA in Figure A.6, the performance of the methods are relatively the same as the theoretical Lipschitz estimates.



Figure A.4. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing forward difference L-BFGS with component MW and random MW Lipschitz estimation schemes.


Figure A.5. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing forward difference L-BFGS with theoretical componentwise Lipschitz estimates and the Moré and Wild Lipschitz estimation schemes.



Figure A.6. Accuracy, Noisy Case with  $\sigma_f = 10^{-3}$ . Log-ratio optimality gap profiles comparing NEWUOA against forward difference L-BFGS with the Moré and Wild Lipschitz estimation schemes. We compare L-BFGS with Component MW (left) and Random MW (right).

## A.2. Investigation of Parameters for NEWUOA

In this section, we empirically investigate the influence of the parameters for NEWUOA for the noisy setting. Although these parameters have been optimized for the noiseless setting, it is not generally known how these parameters may impact its performance on noisy unconstrained problems. By default, NEWUOA employs p = 2n + 1 interpolation points when constructing the quadratic model at each iteration with an initial trust region radius of  $\rho_{\text{beg}} = 1$  and a final trust region radius of  $\rho_{\text{end}} = 10^{-6}$ .

#### A.2.1. Number of Interpolation Points

In Section 2.2, we observed that NEWUOA is able to converge to a better quality neighborhood than L-BFGS with forward differencing, but inferior to central differencing. Since NEWUOA by default employs p = 2n + 1 points, it is natural to both ask if: (1) decreasing the number of interpolation points would yield a less accurate quadratic model, with a potentially less accurate gradient; and (2) increasing the number of points used in the interpolation may improve the accuracy of the solution to be competitive with central differencing.

To do this, we run NEWUOA with p = n + 2 and  $p = \min\left\{3n + 1, \frac{(n+1)(n+2)}{2}\right\}$  interpolation points and compare their optimality gaps and number of function evaluations. In Figures A.7 and A.8, we report the log-ratio profiles for the objective function and function evaluations when comparing NEWUOA with p = n + 2 and p = 2n + 1 points. We see that with higher noise levels, NEWUOA with p = n + 2 tends to terminate earlier, while for lower noise levels, it is less efficient than NEWUOA with p = 2n + 1 points. In terms of solution quality, NEWUOA with p = n + 2 points. This is similarly seen when we compare p = 3n + 1 and p = 2n + 1 points in Figures A.9 and A.10.

When we compare NEWUOA with p = n + 2 points against forward difference L-BFGS, we see that L-BFGS is now competitive against NEWUOA, unlike when NEWUOA employed p = 2n + 1 points; see Figure A.11. However, when we increase the number of points to p = 3n + 1, NEWUOA is still not competitive against central differencing, as seen in Figure A.12.



Figure A.7. Noisy Case. Log-ratio optimality gap profiles comparing NEWUOA with p = 2n + 1 and p = n + 2 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).

# A.2.2. Trust Region Radius

One can also ask if decreasing the final trust region radius could allow NEWUOA to converge to a better solution. To test this, we change  $\rho_{end} = 10^{-12}$  and compare against the default  $\rho_{end} = 10^{-6}$  in Figure A.13. From our experiments, changing the final trust region radius makes almost no difference on the solution quality and efficiency.

# A.3. Complete Numerical Results

In this appendix, we present our complete numerical results.



Figure A.8. Noisy Case. Log-ratio function evaluation profiles comparing NEWUOA with p = 2n + 1 and p = n + 2 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).

# A.3.1. Unconstrained Optimization

A.3.1.1. Noiseless Functions. In this section, we list the number of function evaluations (#feval), ratio between number of function evaluations to the number of function evaluations taken by NEWUOA (RATIO), CPU time (CPU), and optimality gap ( $\phi(x) - \phi^*$ ) for each problem instance in Tables A.1-A.5. Function evaluations marked with a \* denote cases where the algorithm reached the maximum number of function evaluations. Instances where NEWUOA samples a point that satisfies (2.2.7) within the initial 2n + 1 evaluations are denoted by \*\*. Problems marked with a <sup>‡</sup> denote cases where NEWUOA and L-BFGS converge to different minimizers. Optimality gaps marked with a <sup>‡</sup> denote failures of the algorithm to converge to a valid solution satisfying (2.2.7).



Figure A.9. Noisy Case. Log-ratio optimality gap profiles comparing NEWUOA with p = 3n + 1 and p = 2n + 1 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).



Figure A.10. Noisy Case. Log-ratio function evaluation profiles comparing NEWUOA with p = 3n + 1 and p = 2n + 1 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).



Figure A.11. Noisy Case. Log-ratio optimality gap profiles comparing forward difference L-BFGS against NEWUOA with p = n + 2 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).



Figure A.12. Noisy Case. Log-ratio optimality gap profiles comparing central difference L-BFGS against NEWUOA with p = 3n + 1 points. The noise levels are  $\sigma_f = 10^{-1}$  (top left),  $\sigma_f = 10^{-3}$  (top right),  $10^{-5}$  (bottom left), and  $10^{-7}$  (bottom right).



Figure A.13. Noisy Case with  $\sigma_f = 10^{-5}$ . Log-ratio optimality gap profiles comparing NEWUOA with  $\rho_{\text{end}} = 10^{-6}$  and  $10^{-12}$ .

		NEWUOA					FD L-	BFGS		CD L-BFGS			
Problem	n	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$
AIRCRFTB	5	661	1.000	0.1	$6.963 \times 10^{-7}$	296	$4.48 \times 10^{-1}$	0.02	$6.518 \times 10^{-7}$	536	$8.11 \times 10^{-1}$	0.03	$8.288 \times 10^{-7}$
ALLINITU	4	41	1.000	0.09	$3.073 \times 10^{-6}$	58	1.415	0.01	$4.240 \times 10^{-6}$	103	2.512	0.01	$4.240 \times 10^{-6}$
ARWHEAD	100	$201^{**}$	1.000	0.09	0.000	1130	5.622	0.06	$2.474 \times 10^{-7}$	2240	$1.1144 \times 10^{1}$	0.07	$2.476 \times 10^{-7}$
BARD	3	69	1.000	0.09	$7.032 \times 10^{-8}$	102	1.478	0.01	$5.244 \times 10^{-7}$	176	2.551	0.01	$5.246 \times 10^{-7}$
BDQRTIC	100	3682	1.000	1.85	$3.731 \times 10^{-4}$	3178	$8.63 \times 10^{-1}$	0.14	$2.923 \times 10^{-4}$	6308	1.713	0.18	$2.914 \times 10^{-4}$
BIGGS3	3	73	1.000	0.09	$4.599 \times 10^{-7}$	85	1.164	0.01	$5.422 \times 10^{-7}$	152	2.082	0.01	$5.423 \times 10^{-7}$
$BIGGS5^{\ddagger}$	5	103	1.000	0.08	$-3.722 \times 10^{-5}$	416	4.039	0.04	$6.651 \times 10^{-7}$	767	7.447	0.05	$4.321 \times 10^{-7}$
$\mathrm{BIGGS6}^{\ddagger}$	6	338	1.000	0.08	$-1.444 \times 10^{-4}$	323	$9.56 \times 10^{-1}$	0.03	$5.357 \times 10^{-7}$	593	1.754	0.03	$5.654 \times 10^{-7}$
BOX2	2	$2^{**}$	1.000	0.08	$5.847 \times 10^{-19}$	57	$2.8500 \times 10^{1}$	0.01	$9.947 \times 10^{-7}$	95	$4.7500 \times 10^{1}$	0.01	$9.947 \times 10^{-7}$
BOX3	3	$2^{**}$	1.000	0.08	$5.847 \times 10^{-19}$	52	$2.6000 \times 10^{1}$	0.01	$7.371 \times 10^{-7}$	91	$4.5500 \times 10^{1}$	0.01	$7.371 \times 10^{-7}$
BRKMCC	2	10	1.000	0.08	$1.440 \times 10^{-7}$	32	3.200	0.0	$4.066 \times 10^{-9}$	51	5.100	0.0	$4.088 \times 10^{-9}$
BROWNAL	10	917	1.000	0.11	$9.998 \times 10^{-7}$	177	$1.93 \times 10^{-1}$	0.01	$1.707 \times 10^{-7}$	331	$3.61 \times 10^{-1}$	0.01	$1.707 \times 10^{-7}$
BROWNAL	100	18069	1.000	9.53	$9.993 \times 10^{-7}$	1039	$5.8 \times 10^{-2}$	0.07	$1.040 \times 10^{-7}$	2051	$1.14 \times 10^{-1}$	0.1	$1.047 \times 10^{-7}$
BROWNAL	200	$100000^{*}$	1.000	207.31	$3.184 \times 10^{-6}$	823	$8 \times 10^{-3}$	0.13	$8.254 \times 10^{-7}$	1626	$1.6 \times 10^{-2}$	0.16	$8.243 \times 10^{-7}$
BROWNDEN	4	140	1.000	0.09	$2.900 \times 10^{-2}$	131	$9.36 \times 10^{-1}$	0.01	$1.448 \times 10^{-2}$	225	1.607	0.01	$1.449 \times 10^{-2}$
CLIFF	2	84	1.000	0.09	$5.220 \times 10^{-7}$	240	2.857	0.02	$3.101 \times 10^{-7}$	365	4.345	0.02	$2.076 \times 10^{-7}$
CRAGGLVY	4	182	1.000	0.09	$8.761 \times 10^{-7}$	135	$7.42 \times 10^{-1}$	0.01	$5.854 \times 10^{-7}$	244	1.341	0.01	$5.854 \times 10^{-7}$
CRAGGLVY	10	511	1.000	0.1	$1.650 \times 10^{-6}$	482	$9.43 \times 10^{-1}$	0.03	$1.176 \times 10^{-6}$	910	1.781	0.04	$1.168 \times 10^{-6}$
CRAGGLVY	50	2670	1.000	0.45	$1.523 \times 10^{-5}$	2301	$8.62 \times 10^{-1}$	0.1	$1.289 \times 10^{-5}$	4544	1.702	0.14	$1.292 \times 10^{-5}$
CRAGGLVY	100	5679	1.000	2.92	$3.219 \times 10^{-5}$	4503	$7.93 \times 10^{-1}$	0.21	$1.932 \times 10^{-5}$	8946	1.575	0.29	$1.934 \times 10^{-5}$
CUBE	2	173	1.000	0.08	$8.157 \times 10^{-7}$	114	$6.59 \times 10^{-1}$	0.01	$3.107 \times 10^{-7}$	190	1.098	0.01	$3.036 \times 10^{-7}$
DENSCHNA	2	22	1.000	0.09	$2.893 \times 10^{-10}$	38	1.727	0.01	$9.517 \times 10^{-10}$	64	2.909	0.01	$9.521 \times 10^{-10}$
DENSCHNB	2	21	1.000	0.09	$9.135 \times 10^{-8}$	25	1.190	0.0	$3.867 \times 10^{-11}$	42	2.000	0.0	$3.877 \times 10^{-11}$
DENSCHNC	2	66	1.000	0.09	$7.333 \times 10^{-8}$	68	1.030	0.01	$8.949 \times 10^{-8}$	111	1.682	0.01	$8.951 \times 10^{-8}$
DENSCHND	3	312	1.000	0.07	$9.794 \times 10^{-7}$	208	$6.67 \times 10^{-1}$	0.02	$9.534 \times 10^{-7}$	356	1.141	0.02	$9.537 \times 10^{-7}$
DENSCHNE	3	127	1.000	0.09	$4.487 \times 10^{-7}$	141	1.110	0.01	$1.599 \times 10^{-7}$	237	1.866	0.01	$1.210 \times 10^{-7}$
DENSCHNF	2	28	1.000	0.09	$5.079 \times 10^{-9}$	48	1.714	0.01	$5.450 \times 10^{-9}$	77	2.750	0.01	$5.468 \times 10^{-9}$

Table A.1. Noiseless Unconstrained CUTEst Problems Tested. n is the number of variables.

			NE	WUOA			FD L-1	BFGS			CD L-	BFGS	
Problem	n	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$
DIXMAANA	15	596	1.000	0.09	$9.909 \times 10^{-7}$	138	$2.32 \times 10^{-1}$	0.01	$1.096 \times 10^{-7}$	265	$4.45 \times 10^{-1}$	0.01	$1.096 \times 10^{-7}$
DIXMAANA	90	1594	1.000	0.61	$9.626 \times 10^{-7}$	738	$4.63 \times 10^{-1}$	0.03	$6.577 \times 10^{-7}$	1465	$9.19 \times 10^{-1}$	0.04	$6.578 \times 10^{-7}$
DIXMAANA	300	8117	1.000	37.53	$9.936 \times 10^{-7}$	2720	$3.35 \times 10^{-1}$	0.15	$2.979 \times 10^{-9}$	5428	$6.69 \times 10^{-1}$	0.21	$2.978 \times 10^{-9}$
DIXMAANB	15	395	1.000	0.1	$8.586 \times 10^{-7}$	122	$3.09 \times 10^{-1}$	0.01	$3.845 \times 10^{-9}$	233	$5.90 \times 10^{-1}$	0.01	$3.845 \times 10^{-9}$
DIXMAANB	90	1491	1.000	0.56	$9.724 \times 10^{-7}$	647	$4.34 \times 10^{-1}$	0.03	$6.536 \times 10^{-10}$	1283	$8.60 \times 10^{-1}$	0.04	$6.537 \times 10^{-10}$
DIXMAANB	300	6446	1.000	27.48	$9.980 \times 10^{-7}$	2117	$3.28 \times 10^{-1}$	0.15	$7.825 \times 10^{-10}$	4223	$6.55 \times 10^{-1}$	0.15	$7.845 \times 10^{-10}$
DIXMAANC	15	197	1.000	0.08	$9.747 \times 10^{-7}$	158	$8.02 \times 10^{-1}$	0.01	$8.216 \times 10^{-7}$	302	1.533	0.01	$8.216 \times 10^{-7}$
DIXMAANC	90	1611	1.000	0.63	$9.914 \times 10^{-7}$	558	$3.46 \times 10^{-1}$	0.03	$1.352 \times 10^{-7}$	1104	$6.85 \times 10^{-1}$	0.03	$1.353 \times 10^{-7}$
DIXMAANC	300	4643	1.000	18.87	$9.964 \times 10^{-7}$	2120	$4.57 \times 10^{-1}$	0.15	$5.674 \times 10^{-7}$	4227	$9.10 \times 10^{-1}$	0.15	$5.672 \times 10^{-7}$
DIXMAAND	15	251	1.000	0.1	$9.036 \times 10^{-7}$	179	$7.13 \times 10^{-1}$	0.01	$4.536 \times 10^{-7}$	340	1.355	0.01	$4.536 \times 10^{-7}$
DIXMAAND	90	2090	1.000	0.81	$9.552 \times 10^{-7}$	745	$3.56 \times 10^{-1}$	0.04	$6.242 \times 10^{-7}$	1474	$7.05 \times 10^{-1}$	0.04	$6.242 \times 10^{-7}$
DIXMAAND	300	7155	1.000	33.18	$9.791 \times 10^{-7}$	2425	$3.39 \times 10^{-1}$	0.18	$3.136 \times 10^{-7}$	4834	$6.76 \times 10^{-1}$	0.18	$3.136 \times 10^{-7}$
DIXMAANE	15	377	1.000	0.1	$7.225 \times 10^{-7}$	293	$7.77 \times 10^{-1}$	0.02	$4.958 \times 10^{-7}$	564	1.496	0.02	$4.959 \times 10^{-7}$
DIXMAANE	90	3081	1.000	1.28	$9.840 \times 10^{-7}$	3499	1.136	0.15	$8.197 \times 10^{-7}$	6956	2.258	0.2	$8.202 \times 10^{-7}$
DIXMAANE	300	18666	1.000	107.25	$9.993 \times 10^{-7}$	19633	1.052	0.96	$5.643 \times 10^{-7}$	39197	2.100	1.42	$6.031 \times 10^{-7}$
DIXMAANF	15	274	1.000	0.09	$9.301 \times 10^{-7}$	226	$8.25 \times 10^{-1}$	0.01	$7.285 \times 10^{-7}$	434	1.584	0.01	$7.286 \times 10^{-7}$
DIXMAANF	90	3104	1.000	1.33	$9.964 \times 10^{-7}$	2674	$8.61 \times 10^{-1}$	0.11	$3.730 \times 10^{-7}$	5313	1.712	0.15	$3.731 \times 10^{-7}$
DIXMAANF	300	16182	1.000	92.5	$9.988 \times 10^{-7}$	14504	$8.96 \times 10^{-1}$	0.69	$9.006 \times 10^{-7}$	28952	1.789	1.06	$9.011 \times 10^{-7}$
DIXMAANG	15	286	1.000	0.08	$8.435\times10^{-7}$	277	$9.69 \times 10^{-1}$	0.01	$2.101\times10^{-7}$	533	1.864	0.02	$2.102\times10^{-7}$
DIXMAANG	90	3429	1.000	1.47	$9.922 \times 10^{-7}$	2857	$8.33 \times 10^{-1}$	0.12	$6.520 \times 10^{-7}$	5678	1.656	0.16	$6.522 \times 10^{-7}$
DIXMAANG	300	21081	1.000	126.91	$9.999 \times 10^{-7}$	14806	$7.02 \times 10^{-1}$	0.78	$9.894 \times 10^{-7}$	30158	1.431	1.1	$8.871 \times 10^{-7}$
DIXMAANH	15	375	1.000	0.1	$8.025 \times 10^{-7}$	280	$7.47 \times 10^{-1}$	0.02	$1.341 \times 10^{-7}$	537	1.432	0.02	$1.342 \times 10^{-7}$
DIXMAANH	90	3106	1.000	1.27	$9.943 \times 10^{-7}$	2588	$8.33 \times 10^{-1}$	0.11	$6.407 \times 10^{-7}$	5138	1.654	0.15	$6.405 \times 10^{-7}$
DIXMAANH	300	35784	1.000	212.15	$9.999 \times 10^{-7}$	14207	$3.97 \times 10^{-1}$	0.69	$7.042 \times 10^{-7}$	28356	$7.92 \times 10^{-1}$	1.05	$6.752 \times 10^{-7}$
DIXMAANI	15	586	1.000	0.1	$9.553 \times 10^{-7}$	549	$9.37 \times 10^{-1}$	0.03	$9.880 \times 10^{-7}$	1060	1.809	0.04	$9.880 \times 10^{-7}$
DIXMAANI	90	13092	1.000	6.05	$9.966 \times 10^{-7}$	15925	1.216	0.62	$9.862 \times 10^{-7}$	33132	2.531	0.91	$9.270 \times 10^{-7}$
DIXMAANI	300	121834	1.000	774.33	$9.999 \times 10^{-7}$	$150107^{*}$	1.232	6.82	$2.119 \times 10^{-6}$	$150155^{*}$	1.232	5.3	$1.720 \times 10^{-5}$
DIXMAANJ	15	563	1.000	0.1	$9.736 \times 10^{-7}$	479	$8.51 \times 10^{-1}$	0.03	$1.809 \times 10^{-7}$	926	1.645	0.03	$1.797 \times 10^{-7}$
DIXMAANJ	90	12417	1.000	6.11	$9.923 \times 10^{-7}$	11508	$9.27 \times 10^{-1}$	0.53	$9.628 \times 10^{-7}$	22883	1.843	0.74	$9.706 \times 10^{-7}$
DIXMAANJ	300	120362	1.000	718.6	$9.987 \times 10^{-7}$	77321	$6.42 \times 10^{-1}$	3.61	$9.974 \times 10^{-7}$	$150156^{*}$	1.248	5.37	$1.028 \times 10^{-6}$
DIXMAANK	15	559	1.000	0.1	$9.181 \times 10^{-7}$	533	$9.53 \times 10^{-1}$	0.03	$3.242 \times 10^{-7}$	1029	1.841	0.04	$3.230 \times 10^{-7}$
DIXMAANK	90	12601	1.000	5.72	$9.969 \times 10^{-7}$	7463	$5.92 \times 10^{-1}$	0.3	$9.902 \times 10^{-7}$	15018	1.192	0.41	$9.893 \times 10^{-7}$
DIXMAANK	300	80808	1.000	455.59	$9.997 \times 10^{-7}$	23567	$2.92 \times 10^{-1}$	1.12	$9.623 \times 10^{-7}$	47649	$5.90 \times 10^{-1}$	1.74	$9.108 \times 10^{-7}$
DIXMAANL	15	752	1.000	0.1	$9.740 \times 10^{-7}$	502	$6.68 \times 10^{-1}$	0.03	$4.740 \times 10^{-7}$	967	1.286	0.03	$4.714 \times 10^{-7}$
DIXMAANL	90	10608	1.000	4.87	$9.970 \times 10^{-7}$	10323	$9.73 \times 10^{-1}$	0.41	$9.519\times10^{-7}$	20517	1.934	0.57	$8.891\times10^{-7}$
DIXMAANL	300	$150000^{*}$	1.000	875.24	$3.192 \times 10^{-6}$	34747	$2.32 \times 10^{-1}$	1.67	$9.946 \times 10^{-7}$	69364	$4.62 \times 10^{-1}$	2.51	$9.926 \times 10^{-7}$

Table A.2. Noiseless Unconstrained CUTEst Problems Tested. n is the number of variables.

		NEWUOA				FD L-	BFGS		CD L-BFGS				
Problem	n	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$
DQRTIC	10	502	1.000	0.1	$9.719 \times 10^{-7}$	258	$5.14 \times 10^{-1}$	0.02	$4.439 \times 10^{-7}$	488	$9.72 \times 10^{-1}$	0.02	$4.439 \times 10^{-7}$
DQRTIC	50	2847	1.000	0.43	$8.852 \times 10^{-7}$	1467	$5.15 \times 10^{-1}$	0.06	$7.890 \times 10^{-7}$	2894	1.017	0.08	$7.888 \times 10^{-7}$
DQRTIC	100	6061	1.000	2.81	$9.769 \times 10^{-7}$	3277	$5.41 \times 10^{-1}$	0.13	$5.471 \times 10^{-7}$	6508	1.074	0.17	$5.462 \times 10^{-7}$
EDENSCH	36	928	1.000	0.14	$2.175 \times 10^{-4}$	577	$6.22 \times 10^{-1}$	0.03	$7.341 \times 10^{-5}$	1131	1.219	0.03	$7.341 \times 10^{-5}$
EIGENALS	6	$5^{**}$	1.000	0.08	0.000	83	$1.6600 \times 10^{1}$	0.01	$7.344 \times 10^{-9}$	152	$3.0400 \times 10^{1}$	0.01	$7.099 \times 10^{-9}$
EIGENALS	110	$55000^{*}$	1.000	40.44	$1.716\times10^{-3}$	42138	$7.66 \times 10^{-1}$	2.12	$9.758 \times 10^{-7}$	$55108^{*}$	1.002	2.13	$2.142 \times 10^{-5}$
EIGENBLS <sup>‡</sup>	6	113	1.000	0.09	$-3.033 \times 10^{-4}$	91	$8.05 \times 10^{-1}$	0.01	$9.727 \times 10^{-9}$	167	1.478	0.01	$9.145 \times 10^{-9}$
EIGENBLS	110	37505	1.000	27.4	$9.993 \times 10^{-7}$	$55010^{*}$	1.467	2.68	$9.280 \times 10^{-4}$	$55088^{*}$	1.469	2.1	$4.002 \times 10^{-2}$
EIGENCLS	30	1630	1.000	0.19	$9.383 \times 10^{-7}$	2634	1.616	0.12	$9.353 \times 10^{-7}$	5174	3.174	0.16	$8.934 \times 10^{-7}$
ENGVAL1	2	32	1.000	0.08	$6.651 \times 10^{-9}$	38	1.188	0.0	$2.191 \times 10^{-7}$	60	1.875	0.0	$2.191 \times 10^{-7}$
ENGVAL1	50	2234	1.000	0.38	$5.211 \times 10^{-5}$	585	$2.62 \times 10^{-1}$	0.02	$3.074 \times 10^{-5}$	1148	$5.14 \times 10^{-1}$	0.03	$3.074 \times 10^{-5}$
ENGVAL1	100	1941	1.000	0.92	$1.076 \times 10^{-4}$	1337	$6.89 \times 10^{-1}$	0.07	$4.197 \times 10^{-6}$	2651	1.366	0.07	$4.198 \times 10^{-6}$
EXPFIT	2	42	1.000	0.09	$9.588 \times 10^{-7}$	46	1.095	0.01	$1.427 \times 10^{-9}$	75	1.786	0.0	$1.420 \times 10^{-9}$
$FLETCBV3^{\ddagger}$	10	1106	1.000	0.11	$1.063 \times 10^{-3}$	261	$2.36 \times 10^{-1}$	0.01	$-1.664 \times 10^{-5}$	521	$4.71 \times 10^{-1}$	0.02	$-4.710 \times 10^{-5}$
$FLETCBV3^{\ddagger}$	100	50000*	1.000	28.66	$1.333 \times 10^5$	19708	$3.94 \times 10^{-1}$	0.87	$-3.486 \times 10^{-1}$	16125	$3.23 \times 10^{-1}$	0.54	$-1.107 \times 10^1$
$FLETCHBV^{\ddagger}$	10	945	1.000	0.09	$1.490 \times 10^{5}$	263	$2.78 \times 10^{-1}$	0.02	$-1.278 \times 10^1$	278	$2.94 \times 10^{-1}$	0.01	$-5.322 \times 10^3$
$FLETCHBV^{\ddagger}$	100	50000*	1.000	27.77	$1.210\times10^{13}$	18068	$3.61 \times 10^{-1}$	0.8	$-6.240 \times 10^6$	20933	$4.19 \times 10^{-1}$	0.67	$-7.016 \times 10^7$
FREUROTH	$^{2}$	55	1.000	0.09	$2.214 \times 10^{-5}$	81	1.473	0.01	$3.932 \times 10^{-5}$	134	2.436	0.01	$3.932 \times 10^{-5}$
FREUROTH	10	367	1.000	0.09	$6.524 \times 10^{-4}$	286	$7.79 \times 10^{-1}$	0.02	$1.194 \times 10^{-6}$	538	1.466	0.02	$1.194 \times 10^{-6}$
FREUROTH	50	5840	1.000	0.86	$5.844 \times 10^{-3}$	1053	$1.80 \times 10^{-1}$	0.04	$4.888 \times 10^{-4}$	2073	$3.55 \times 10^{-1}$	0.06	$4.886 \times 10^{-4}$
FREUROTH	100	2881	1.000	1.44	$1.188\times 10^{-2}$	1950	$6.77 \times 10^{-1}$	0.09	$2.601 \times 10^{-3}$	3868	1.343	0.12	$2.589 \times 10^{-3}$
GENROSE	5	199	1.000	0.08	$9.086 \times 10^{-7}$	226	1.136	0.02	$9.527 \times 10^{-9}$	411	2.065	0.02	$9.367 \times 10^{-9}$
GENROSE	10	645	1.000	0.1	$8.369 \times 10^{-7}$	896	1.389	0.06	$8.050 \times 10^{-7}$	1614	2.502	0.06	$1.032 \times 10^{-7}$
GENROSE	100	17249	1.000	9.45	$9.954 \times 10^{-7}$	25453	1.476	1.01	$8.013 \times 10^{-7}$	$50198^{*}$	2.910	1.4	$1.939 \times 10^{-5}$
GULF	3	876	1.000	0.12	$6.954 \times 10^{-7}$	225	$2.57 \times 10^{-1}$	0.03	$7.913 \times 10^{-8}$	392	$4.47 \times 10^{-1}$	0.04	$6.048 \times 10^{-7}$
HAIRY	2	261	1.000	0.09	$1.214 \times 10^{-6}$	86	$3.30 \times 10^{-1}$	0.01	$2.209 \times 10^{-7}$	149	$5.71 \times 10^{-1}$	0.01	$2.996 \times 10^{-7}$
HELIX	3	65	1.000	0.09	$3.569 \times 10^{-8}$	149	2.292	0.01	$6.812 \times 10^{-7}$	86	1.323	0.0	$8.521 \times 10^{2^{+}}$
JENSMP <sup>‡</sup>	$^{2}$	4**	1.000	0.09	$-5.469 \times 10^2$	16	4.000	0.0	0.000	21	5.250	0.0	0.000
KOWOSB	4	150	1.000	0.09	$7.089 \times 10^{-7}$	183	1.220	0.02	$2.824 \times 10^{-7}$	332	2.213	0.02	$2.825 \times 10^{-7}$
MEXHAT	2	46	1.000	0.09	$1.036\times10^{-2\dagger}$	203	4.413	0.02	$7.395 \times 10^{-7}$	339	7.370	0.02	$4.707 \times 10^{-7}$

Table A.3. Noiseless Unconstrained CUTEst Problems Tested. n is the number of variables.

		NEWUOA					FD I	-BFGS		CD L-BFGS			
Problem	n	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x)-\phi^*$
MOREBV	10	379	1.000	0.09	$9.458 \times 10^{-7}$	366	$9.66 \times 10^{-1}$	0.02	$6.782 \times 10^{-7}$	695	1.834	0.03	$6.843 \times 10^{-7}$
MOREBV	50	21488	1.000	3.71	$9.988 \times 10^{-7}$	$25029^{*}$	1.165	0.98	$1.368 \times 10^{-6}$	$25035^{*}$	1.165	0.68	$6.088 \times 10^{-6}$
MOREBV	100	11143	1.000	6.29	$9.999 \times 10^{-7}$	15512	1.392	0.59	$9.985 \times 10^{-7}$	31067	2.788	0.82	$9.992 \times 10^{-7}$
NCB20B	21	478	1.000	0.11	$4.093 \times 10^{-5}$	169	$3.54 \times 10^{-1}$	0.01	$-4.368 \times 10^{-11}$	322	$6.74 \times 10^{-1}$	0.01	$-4.334 \times 10^{-11}$
NCB20B	22	627	1.000	0.11	$4.361 \times 10^{-5}$	154	$2.46 \times 10^{-1}$	0.01	$2.493 \times 10^{-5}$	292	$4.66 \times 10^{-1}$	0.01	$2.493 \times 10^{-5}$
NCB20B	50	3345	1.000	0.59	$9.936 \times 10^{-5}$	1465	$4.38 \times 10^{-1}$	0.07	$9.906 \times 10^{-5}$	2892	$8.65 \times 10^{-1}$	0.1	$9.912 \times 10^{-5}$
NCB20B	100	7309	1.000	4.02	$1.963 \times 10^{-4}$	3376	$4.62 \times 10^{-1}$	0.21	$1.794 \times 10^{-4}$	6708	$9.18 \times 10^{-1}$	0.3	$1.797 \times 10^{-4}$
NCB20B	180	11313	1.000	22.07	$3.470 \times 10^{-4}$	4375	$3.87 \times 10^{-1}$	0.35	$3.384 \times 10^{-4}$	8718	$7.71 \times 10^{-1}$	0.56	$3.379 \times 10^{-4}$
NONDIA	10	198	1.000	0.1	$4.267 \times 10^{-7}$	167	$8.43 \times 10^{-1}$	0.01	$1.459 \times 10^{-11}$	298	1.505	0.01	$1.053 \times 10^{-11}$
NONDIA	20	426	1.000	0.1	$8.903 \times 10^{-7}$	341	$8.00 \times 10^{-1}$	0.02	$2.571 \times 10^{-10}$	655	1.538	0.02	$2.929 \times 10^{-10}$
NONDIA	30	569	1.000	0.12	$9.936 \times 10^{-7}$	492	$8.65 \times 10^{-1}$	0.02	$3.712 \times 10^{-7}$	956	1.680	0.03	$3.738 \times 10^{-7}$
NONDIA	50	799	1.000	0.2	$9.829 \times 10^{-7}$	792	$9.91 \times 10^{-1}$	0.03	$3.825 \times 10^{-8}$	1556	1.947	0.04	$3.691 \times 10^{-8}$
NONDIA	90	1302	1.000	0.58	$9.192 \times 10^{-7}$	1301	$9.99 \times 10^{-1}$	0.06	$1.233 \times 10^{-7}$	2574	1.977	0.07	$1.188 \times 10^{-7}$
NONDIA	100	1683	1.000	0.89	$9.873 \times 10^{-7}$	1543	$9.17 \times 10^{-1}$	0.07	$1.028 \times 10^{-10}$	3057	1.816	0.09	$1.125 \times 10^{-10}$
NONDQUAR	100	$50000^{*}$	1.000	29.28	$8.675 \times 10^{-6}$	50056*	1.001	1.88	$1.137 \times 10^{-5}$	$50189^{*}$	1.004	1.32	$3.611 \times 10^{-5}$
$OSBORNEA^{\ddagger}$	5	1094	1.000	0.09	$1.204 \times 10^{-5}$	477	$4.36 \times 10^{-1}$	0.05	$8.442 \times 10^{-7}$	888	$8.12 \times 10^{-1}$	0.06	$8.390 \times 10^{-7}$
OSBORNEB	11	1529	1.000	0.14	$9.678 \times 10^{-7}$	1044	$6.83 \times 10^{-1}$	0.08	$9.551 \times 10^{-7}$	2216	1.449	0.11	$6.312 \times 10^{-7}$
PENALTY1	4	577	1.000	0.09	$9.994 \times 10^{-7}$	868	1.504	0.09	$9.134 \times 10^{-7}$	1425	2.470	0.09	$9.856 \times 10^{-7}$
PENALTY1	10	1332	1.000	0.11	$9.944 \times 10^{-7}$	1523	1.143	0.09	$9.903 \times 10^{-7}$	2942	2.209	0.11	$8.873 \times 10^{-7}$
PENALTY1	50	6082	1.000	0.89	$9.985 \times 10^{-7}$	5773	$9.49 \times 10^{-1}$	0.22	$9.994 \times 10^{-7}$	11797	1.940	0.32	$9.652 \times 10^{-7}$
PENALTY1	100	13215	1.000	6.44	$9.969 \times 10^{-7}$	11175	$8.46 \times 10^{-1}$	0.43	$7.672 \times 10^{-7}$	21381	1.618	0.57	$9.919 \times 10^{-7}$
$\rm PFIT1LS^{\ddagger}$	3	417	1.000	0.11	$2.892 \times 10^{-4}$	351	$8.42 \times 10^{-1}$	0.03	$9.871 \times 10^{-7}$	403	$9.66 \times 10^{-1}$	0.03	$6.795 \times 10^{-7}$
$PFIT2LS^{\ddagger}$	3	767	1.000	0.1	$1.243 \times 10^{-2}$	$1502^{*}$	1.958	0.17	$2.818 \times 10^{-4}$	$1509^{*}$	1.967	0.11	$1.528 \times 10^{-3}$
$PFIT3LS^{\ddagger}$	3	907	1.000	0.1	$8.228\times 10^{-2}$	1502*	1.656	0.17	$1.504 \times 10^{-2}$	$1501^{*}$	1.655	0.11	$3.622 \times 10^{-2}$
PFIT4LS <sup>‡</sup>	3	1088	1.000	0.1	$2.673 \times 10^{-1}$	1505*	1.383	0.15	$5.424 \times 10^{-2}$	$1508^{*}$	1.386	0.1	$9.180 \times 10^{-2}$
QUARTC	25	1004	1.000	0.13	$8.886 \times 10^{-7}$	765	$7.62 \times 10^{-1}$	0.04	$4.947 \times 10^{-7}$	1492	1.486	0.05	$4.946 \times 10^{-7}$
QUARTC	100	6061	1.000	2.82	$9.769\times10^{-7}$	3277	$5.41 \times 10^{-1}$	0.13	$5.471 \times 10^{-7}$	6508	1.074	0.16	$5.462 \times 10^{-7}$
SINEVAL	2	254	1.000	0.08	$6.706\times10^{-9}$	309	1.217	0.04	$3.585 \times 10^{-9}$	522	2.055	0.04	$1.305 \times 10^{-7}$

Table A.4. Noiseless Unconstrained CUTEst Problems Tested. n is the number of variables.

		NEWUOA			FD L-BFGS				CD L-BFGS				
Problem	n	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$	#feval	RATIO	CPU	$\phi(x) - \phi^*$
SINQUAD	5	125	1.000	0.09	$6.682\times10^{-6}$	61	$4.88 \times 10^{-1}$	0.01	$3.061\times10^{-6}$	108	$8.64 \times 10^{-1}$	0.01	$3.061 \times 10^{-6}$
SINQUAD	50	3325	1.000	0.55	$1.132 \times 10^{-3}$	795	$2.39 \times 10^{-1}$	0.04	$4.489 \times 10^{-6}$	1562	$4.70 \times 10^{-1}$	0.05	$4.553 \times 10^{-6}$
SINQUAD	100	8318	1.000	4.3	$3.985 \times 10^{-3}$	1545	$1.86 \times 10^{-1}$	0.07	$2.755 \times 10^{-3}$	3062	$3.68 \times 10^{-1}$	0.09	$2.758 \times 10^{-3}$
SISSER	2	24	1.000	0.09	$8.489 \times 10^{-7}$	54	2.250	0.01	$5.876 \times 10^{-7}$	92	3.833	0.01	$5.876 \times 10^{-7}$
SPARSQUR	10	184	1.000	0.09	$8.876 \times 10^{-7}$	184	1.000	0.01	$6.336 \times 10^{-7}$	348	1.891	0.01	$6.336 \times 10^{-7}$
SPARSQUR	50	1645	1.000	0.29	$6.775 \times 10^{-7}$	1046	$6.36 \times 10^{-1}$	0.04	$3.744 \times 10^{-7}$	2065	1.255	0.06	$3.744 \times 10^{-7}$
SPARSQUR	100	2932	1.000	1.37	$8.999 \times 10^{-7}$	2150	$7.33 \times 10^{-1}$	0.09	$3.880 \times 10^{-7}$	4270	1.456	0.12	$3.880 \times 10^{-7}$
TOINTGSS	10	234	1.000	0.09	$1.038 \times 10^{-5}$	24	$1.03 \times 10^{-1}$	0.0	0.000	45	$1.92 \times 10^{-1}$	0.0	0.000
TOINTGSS	50	775	1.000	0.18	$9.723 \times 10^{-6}$	104	$1.34 \times 10^{-1}$	0.0	$-3.000 \times 10^{-9}$	205	$2.65 \times 10^{-1}$	0.01	$-3.000 \times 10^{-9}$
TOINTGSS	100	1199	1.000	0.55	$9.778 \times 10^{-6}$	204	$1.70 \times 10^{-1}$	0.02	$4.000 \times 10^{-9}$	405	$3.38 \times 10^{-1}$	0.01	$4.000 \times 10^{-9}$
TQUARTIC	5	92	1.000	0.09	$5.444 \times 10^{-7}$	85	$9.24 \times 10^{-1}$	0.01	$1.078 \times 10^{-8}$	150	1.630	0.01	$1.074 \times 10^{-8}$
TQUARTIC	10	277	1.000	0.09	$9.847 \times 10^{-7}$	172	$6.21 \times 10^{-1}$	0.01	$5.228 \times 10^{-9}$	325	1.173	0.01	$5.313 \times 10^{-9}$
TQUARTIC	50	5022	1.000	0.75	$9.949 \times 10^{-7}$	646	$1.29 \times 10^{-1}$	0.03	$1.681 \times 10^{-8}$	1265	$2.52 \times 10^{-1}$	0.03	$1.725 \times 10^{-8}$
TQUARTIC	100	20515	1.000	10.51	$9.991 \times 10^{-7}$	1538	$7.5 \times 10^{-2}$	0.07	$6.818 \times 10^{-9}$	3053	$1.49 \times 10^{-1}$	0.08	$6.223 \times 10^{-9}$
TRIDIA	10	186	1.000	0.09	$9.019 \times 10^{-7}$	210	1.129	0.01	$5.451 \times 10^{-8}$	396	2.129	0.01	$5.446 \times 10^{-8}$
TRIDIA	20	446	1.000	0.1	$8.062 \times 10^{-7}$	690	1.547	0.04	$8.068 \times 10^{-7}$	1340	3.004	0.04	$8.064 \times 10^{-7}$
TRIDIA	30	768	1.000	0.13	$9.962 \times 10^{-7}$	1481	1.928	0.07	$6.309 \times 10^{-7}$	2906	3.784	0.08	$6.315 \times 10^{-7}$
TRIDIA	50	1446	1.000	0.3	$8.806 \times 10^{-7}$	3128	2.163	0.15	$9.300 \times 10^{-7}$	6187	4.279	0.19	$9.292 \times 10^{-7}$
TRIDIA	100	3450	1.000	2.11	$9.968 \times 10^{-7}$	8783	2.546	0.32	$9.677 \times 10^{-7}$	17468	5.063	0.43	$9.208 \times 10^{-7}$
WATSON	12	4966	1.000	0.21	$9.982 \times 10^{-7}$	1324	$2.67 \times 10^{-1}$	0.08	$9.773 \times 10^{-7}$	2708	$5.45 \times 10^{-1}$	0.11	$9.950 \times 10^{-7}$
WATSON	31	$15500^{*}$	1.000	1.11	$6.438 \times 10^{-5}$	13875	$8.95 \times 10^{-1}$	0.65	$9.957 \times 10^{-7}$	$15530^{*}$	1.002	0.51	$5.664 \times 10^{-6}$
WOODS	4	497	1.000	0.08	$1.434 \times 10^{-7}$	178	$3.58 \times 10^{-1}$	0.02	$4.447 \times 10^{-9}$	312	$6.28 \times 10^{-1}$	0.02	$4.353 \times 10^{-9}$
WOODS	100	$50000^{*}$	1.000	28.69	$1.371 \times 10^{-4}$	2972	$5.9 \times 10^{-2}$	0.13	$7.208 \times 10^{-7}$	5902	$1.18 \times 10^{-1}$	0.16	$6.252 \times 10^{-7}$
ZANGWIL2	2	12	1.000	0.09	$5.085\times10^{-7}$	11	$9.17 \times 10^{-1}$	0.0	$-1.000 \times 10^{-10}$	19	1.583	0.0	$-9.999 \times 10^{-11}$

Table A.5. Noiseless Unconstrained CUTEst Problems Tested. n is the number of variables.

A.3.1.2. Noisy Functions. In this section, we list the best optimality gap  $(\phi(x) - \phi^*)$ and the number of function evaluations (#feval) needed to achieve this for each problem instance, varying the noise level  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ , in Tables A.6-A.18. Function evaluations marked with a \* denote cases where the algorithm reached the maximum number of function evaluations. Instances where NEWUOA samples a point that satisfies (2.2.7) within the initial 2d + 1 evaluations are denoted by \*\*. Problems marked with a <sup>‡</sup> denote cases where NEWUOA and L-BFGS converge to different minimizers.

#### A.3.2. Nonlinear Least Squares Problems

A.3.2.1. Noiseless Functions. In this section, we list the number of function evaluations (#feval), CPU time (CPU), and optimality gap  $(\phi(x) - \phi^*)$  for each problem instance in Tables A.19. Function evaluations marked with a \* denote cases where the algorithm reached the maximum number of function evaluations.

A.3.2.2. Noisy Functions. In this section, we list the best optimality gap  $(\phi(x) - \phi^*)$  and the number of function evaluations (#feval) needed to achieve this for each problem instance, varying the noise level  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ , in Tables A.21-A.24. Function evaluations marked with a \* denote cases where the algorithm reached the maximum number of function evaluations.

## A.3.3. Constrained Optimization

# A.3.3.1. Test Problem Summary.

			$  NEWUOA \\ f_f   #feval \phi(x) - \phi^*$		FI	O L-BFGS	Cl	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
AIRCRFTB	5	$1 \times 10^{-1}$	43	$7.257\times10^{-1}$	53	$7.277\times 10^{-1}$	231	$2.623\times 10^{-1}$
AIRCRFTB	5	$1 \times 10^{-3}$	70	$1.772 \times 10^{-1}$	715	$9.938\times10^{-2}$	904	$6.648  imes 10^{-5}$
AIRCRFTB	5	$1 \times 10^{-5}$	159	$1.006\times 10^{-2}$	1221	$3.355\times 10^{-4}$	927	$1.042\times 10^{-5}$
AIRCRFTB	5	$1 \times 10^{-7}$	317	$2.893 imes10^{-4}$	1176	$4.119\times10^{-6}$	1106	$7.139  imes 10^{-12}$
ALLINITU	4	$1 \times 10^{-1}$	26	$1.020  imes 10^{-1}$	24	2.881	108	$1.980 imes10^{-3}$
ALLINITU	4	$1 \times 10^{-3}$	38	$9.901  imes 10^{-4}$	97	$2.614\times 10^{-4}$	189	$7.467\times10^{-7}$
ALLINITU	4	$1 \times 10^{-5}$	44	$1.337 \times 10^{-6}$	62	$1.721 \times 10^{-5}$	320	$5.606 \times 10^{-9}$
ALLINITU	4	$1 \times 10^{-7}$	58	$1.186  imes 10^{-8}$	74	$1.612 \times 10^{-7}$	170	$3.518 \times 10^{-10}$
ARWHEAD	100	$1 \times 10^{-1}$	$201^{**}$	0.000	508	$8.433 \times 10^{-2}$	1022	$5.894 \times 10^{-2}$
ARWHEAD	100	$1 \times 10^{-3}$	$201^{**}$	0.000	1035	$3.525 \times 10^{-2}$	2037	$1.305 \times 10^{-3}$
ARWHEAD	100	$1 \times 10^{-5}$	$201^{**}$	0.000	1765	$1.925 \times 10^{-4}$	2444	$5.061 \times 10^{-6}$
ARWHEAD	100	$1 \times 10^{-7}$	$201^{**}$	0.000	1434	$1.346 \times 10^{-5}$	2660	$6.841 \times 10^{-9}$
BARD	3	$1 \times 10^{-1}$	20	$1.665 \times 10^{-1}$	53	$9.347 \times 10^{-1}$	102	$2.119 \times 10^{-3}$
BARD	3	$1 \times 10^{-3}$	36	$5.403 \times 10^{-4}$	77	$1.030 \times 10^{-3}$	116	$1.876 \times 10^{-3}$
BARD	3	$1 \times 10^{-5}$	47	$1.256 \times 10^{-3}$	294	$1.898 \times 10^{-3}$	310	$2.731 \times 10^{-6}$
BARD	3	$1 \times 10^{-7}$	87	$1.685 \times 10^{-7}$	198	$6.076 \times 10^{-8}$	230	$3.753 \times 10^{-9}$
BDQRTIC	100	$1 \times 10^{-1}$	2283	1.132	2482	5.246	3476	$2.697 \times 10^{-1}$
BDQRTIC	100	$1 \times 10^{-3}$	4736	$9.554 \times 10^{-2}$	2664	$1.099 \times 10^{-1}$	6583	$4.061 \times 10^{-3}$
BDQRTIC	100	$1 \times 10^{-5}$	7755	$4.134 \times 10^{-3}$	6619	$3.957 \times 10^{-4}$	11431	$9.216 \times 10^{-7}$
BDQRTIC	100	$1 \times 10^{-7}$	5083	$6.419 \times 10^{-6}$	6543	$6.829 \times 10^{-6}$	9934	$4.423 \times 10^{-9}$
BIGGS3 <sup>‡</sup>	3	$1 \times 10^{-1}$	8	1.366	160	1.311	168	$9.756 \times 10^{-1}$
BIGGS3 <sup>‡</sup>	3	$1 \times 10^{-3}$	67	$2.251 \times 10^{-3}$	155	$6.556 \times 10^{-4}$	389	$1.958 \times 10^{-5}$
BIGGS3 <sup>‡</sup>	3	$1 \times 10^{-5}$	73	$1.776 \times 10^{-6}$	78	$2.858 \times 10^{-6}$	293	$1.380 \times 10^{-6}$
BIGGS3 <sup>‡</sup>	3	$1 \times 10^{-7}$	85	$1.154 \times 10^{-8}$	95	$1.538 \times 10^{-6}$	389	$8.357 \times 10^{-9}$
BIGGS5 <sup>‡</sup>	5	$1 \times 10^{-1}$	30	1.300	15	1.373	435	$1.113 \times 10^{-1}$
BIGGS5 <sup>‡</sup>	5	$1 \times 10^{-3}$	113	$3.932 \times 10^{-2}$	266	$1.181 \times 10^{-1}$	253	$3.142 \times 10^{-2}$
BIGGS5+	5	$1 \times 10^{-5}$	176	$-4.496 \times 10^{-3}$	285	$1.808 \times 10^{-2}$	818	$8.487 \times 10^{-6}$
BIGGS <sup>‡</sup> 5	5	$1 \times 10^{-7}$	265	$-5.529 \times 10^{-3}$	580	$6.905 \times 10^{-5}$	1077	$2.321 \times 10^{-8}$
BIGGS6	6	$1 \times 10^{-1}$	30	$3.144 \times 10^{-1}$	121	$3.515 \times 10^{-1}$	283	$2.855 \times 10^{-1}$
BIGGS6	6	$1 \times 10^{-3}$	43	$2.937 \times 10^{-1}$	126	$2.897 \times 10^{-1}$	1549	$3.092 \times 10^{-2}$
BIGGS6	6	$1 \times 10^{-5}$	363	$-4.550 \times 10^{-3}$	586	$-3.402 \times 10^{-3}$	1324	$-5.568 \times 10^{-3}$
BIGGS6	6	$1 \times 10^{-7}$	657	$-5.612 \times 10^{-3}$	1612	$-5.608 \times 10^{-3}$	1778	$-5.648 \times 10^{-3}$
BOX2	2	$1 \times 10^{-1}$	2**	$5.847 \times 10^{-19}$	187	$7.590 \times 10^{-2}$	26	$3.168 \times 10^{-1}$
BOX2	2	$1 \times 10^{-3}$	2**	$5.847 \times 10^{-19}$	66	$4.861 \times 10^{-1}$	387	$1.276 \times 10^{-6}$
BOX2	2	$1 \times 10^{-3}$	2**	$5.847 \times 10^{-19}$	42	$4.089 \times 10^{-6}$	132	$1.362 \times 10^{-6}$
BOX2	2	$1 \times 10^{-7}$	2**	$5.847 \times 10^{-19}$	56	$1.849 \times 10^{-6}$	108	$2.693 \times 10^{-8}$
BOX3	3	$1 \times 10^{-1}$	2**	$5.847 \times 10^{-19}$	28	$2.056 \times 10^{-2}$	55	$2.102 \times 10^{-2}$
BOX3	3	$1 \times 10^{-3}$	2**	$5.847 \times 10^{-19}$	181	$1.408 \times 10^{-3}$	94	$4.054 \times 10^{-6}$
BOX3	3	$1 \times 10^{-3}$	2**	$5.847 \times 10^{-19}$	86	$9.697 \times 10^{-6}$		$6.273 \times 10^{-7}$
BOX3	3	$1 \times 10^{-7}$	2**	$5.847 \times 10^{-19}$	67	$0.791 \times 10^{-7}$	370	$4.093 \times 10^{-8}$
BRKMCC	2	$1 \times 10^{-1}$	14	$1.160 \times 10^{-2}$	107	$1.085 \times 10^{-1}$		$4.542 \times 10^{-3}$
BRKMCC	2	$1 \times 10^{-3}$ 1 10 <sup>-5</sup>	22	$8.434 \times 10^{-8}$	55	$3.649 \times 10^{-4}$		$1.407 \times 10^{-0}$
BRKMUU	2	$1 \times 10^{-7}$	10	$1.3(1 \times 10^{-1})$	45	$5.221 \times 10^{-5}$	62	$2.505 \times 10^{-10}$
BRKMCC	2	$1 \times 10^{-7}$	19	$3.325 \times 10^{-10}$	34	$4.961 \times 10^{-8}$	156	$1.995 \times 10^{-10}$

Table A.6. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			N	EWUOA	FD	L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
BROWNAL	10	$1 \times 10^{-1}$	46	$9.212\times 10^{-2}$	53	$2.719\times 10^{-2}$	384	$4.202\times 10^{-5}$
BROWNAL	10	$1 \times 10^{-3}$	123	$1.232\times 10^{-2}$	65	$5.759\times10^{-5}$	223	$2.969\times 10^{-5}$
BROWNAL	10	$1 \times 10^{-5}$	248	$3.218\times 10^{-4}$	77	$3.062\times 10^{-5}$	740	$1.092 \times 10^{-5}$
BROWNAL	10	$1 \times 10^{-7}$	458	$6.915\times10^{-6}$	77	$2.974\times10^{-5}$	821	$2.971\times 10^{-5}$
BROWNAL	100	$1 \times 10^{-1}$	610	1.164	725	$5.055\times10^{-2}$	2664	$1.297  imes 10^{-5}$
BROWNAL	100	$1 \times 10^{-3}$	4084	$2.877\times10^{-2}$	725	$5.541\times10^{-4}$	3677	$1.287\times10^{-5}$
BROWNAL	100	$1 \times 10^{-5}$	7907	$3.839\times10^{-4}$	725	$1.908\times 10^{-5}$	3763	$3.096 imes10^{-8}$
BROWNAL	100	$1 \times 10^{-7}$	11449	$7.999 imes10^{-5}$	827	$1.336\times 10^{-5}$	3540	$2.445\times 10^{-8}$
BROWNAL	200	$1 \times 10^{-1}$	1514	$7.945 \times 10^{-1}$	821	$5.423 \times 10^{-3}$	7494	$8.719 \times 10^{-7}$
BROWNAL	200	$1 \times 10^{-3}$	11839	$3.645\times10^{-2}$	821	$4.205\times10^{-4}$	7494	$8.234\times10^{-7}$
BROWNAL	200	$1 \times 10^{-5}$	22817	$3.574 imes10^{-4}$	1833	$1.073  imes 10^{-5}$	5291	$9.198  imes 10^{-9}$
BROWNAL	200	$1 \times 10^{-7}$	26754	$5.228 \times 10^{-6}$	1833	$9.284 \times 10^{-7}$	1626	$8.243 \times 10^{-7}$
BROWNDEN	4	$1 \times 10^{-1}$	219	$3.762\times10^{-2}$	151	$5.209  imes 10^{-1}$	327	$1.623  imes 10^{-4}$
BROWNDEN	4	$1 \times 10^{-3}$	163	$3.183 \times 10^{-4}$	358	$3.298 \times 10^{-3}$	276	$5.083 \times 10^{-7}$
BROWNDEN	4	$1 \times 10^{-5}$	239	$1.428 \times 10^{-6}$	351	$3.116 \times 10^{-5}$	751	$2.212 \times 10^{-9}$
BROWNDEN	4	$1 \times 10^{-7}$	224	$1.020 \times 10^{-8}$	159	$1.909\times10^{-7}$	292	$-4.075 \times 10^{-10}$
CLIFF	2	$1 \times 10^{-1}$	19	$9.252 \times 10^{-2}$	51	$2.902 \times 10^2$	809	$5.153 \times 10^{-1}$
CLIFF	2	$1 \times 10^{-3}$	31	$1.389 \times 10^{-3}$	201	$2.902 \times 10^2$	412	$3.190 \times 10^{-4}$
CLIFF	2	$1 \times 10^{-5}$	25	$8.089 \times 10^{-4}$	644	1.027	272	$2.209 \times 10^{-4}$
CLIFF	2	$1 \times 10^{-7}$	68	$1.705 \times 10^{-8}$	407	$4.054 \times 10^{-2}$	335	$2.192 \times 10^{-4}$
CRAGGLVY	4	$1 \times 10^{-1}$	14	$8.946 \times 10^{-1}$	253	$7.327 \times 10^{-1}$	285	$4.782 \times 10^{-2}$
CRAGGLVY	4	$1 \times 10^{-3}$	72	$4.048 \times 10^{-3}$	271	$7.551 \times 10^{-3}$	147	$5.397 \times 10^{-4}$
CRAGGLVY	4	$1 \times 10^{-5}$	173	$1.960 \times 10^{-5}$	262	$8.396 \times 10^{-5}$	338	$6.253 \times 10^{-6}$
CRAGGLVY	4	$1 \times 10^{-7}$	199	$1.186 \times 10^{-6}$	160	$3.749 \times 10^{-4}$	620	$3.188 \times 10^{-8}$
CRAGGLVY	10	$1 \times 10^{-1}$	195	$6.511 \times 10^{-1}$	663	$3.023 \times 10^{1}$	912	$2.336 \times 10^{-1}$
CRAGGLVY	10	$1 \times 10^{-3}$	345	$4.337 \times 10^{-2}$	731	$2.444 \times 10^{-2}$	828	$5.647 \times 10^{-4}$
CRAGGLVY	10	$1 \times 10^{-5}$	573	$7.120 \times 10^{-4}$	716	$6.673 \times 10^{-4}$	1466	$4.751 \times 10^{-7}$
CRAGGLVY	10	$1 \times 10^{-7}$	1066	$7.874 \times 10^{-7}$	1602	$5.150 \times 10^{-6}$	1466	$4.121 \times 10^{-9}$
CRAGGLVY	50	$1 \times 10^{-1}$	1364	1.968	2041	5.834	3529	$8.585 \times 10^{-1}$
CRAGGLVY	50	$1 \times 10^{-3}$	1853	$4.989 \times 10^{-2}$	4908	$5.233 \times 10^{-2}$	4192	$9.901 \times 10^{-3}$
CRAGGLVY	50	$1 \times 10^{-3}$	2312	$4.363 \times 10^{-4}$	3101	$9.898 \times 10^{-4}$	5860	$1.281 \times 10^{-3}$
CRAGGLVY	50	$1 \times 10^{-7}$	3654	$2.357 \times 10^{-6}$		$9.978 \times 10^{-6}$	7120	$4.856 \times 10^{-8}$
CRAGGLVY	100	$1 \times 10^{-1}$	3091	2.974	5275	$1.228 \times 10^{1}$	7272	1.469
CRAGGLVY	100	$1 \times 10^{-5}$	5041	$1.878 \times 10^{-4}$	3918	$6.386 \times 10^{-1}$	7959	$1.298 \times 10^{-2}$
CRAGGLVY	100	$1 \times 10^{-5}$	4955	$4.727 \times 10^{-4}$	6241	$5.729 \times 10^{-5}$	15314	$1.573 \times 10^{-5}$
CRAGGLV Y	100	$1 \times 10^{-1}$	7008	$5.223 \times 10^{-6}$	11106	$3.053 \times 10^{-9}$	12688	$2.635 \times 10^{-1}$
CUBE	2	$1 \times 10^{-1}$	26	2.098	34	$6.758 \times 10^{-2}$	65	$4.380 \times 10^{-2}$
CUBE	2	$1 \times 10^{-5}$	195	$9.524 \times 10^{-2}$	43	$4.104 \times 10^{-2}$	340	$7.350 \times 10^{-7}$
CUBE	2	$1 \times 10^{-7}$	135	$1.075 \times 10^{-6}$	40	$4.207 \times 10^{-2}$	213	$2.049 \times 10^{-8}$
DENICUMA	2	$1 \times 10^{-1}$	140	$3.014 \times 10^{-0}$	141	$9.280 \times 10^{-1}$	203	$8.023 \times 10^{-8}$ 7.120 $\times 10^{-2}$
DENSCHNA	2	$1 \times 10^{-1}$	12	$1.802 \times 10^{-1}$		$1.032 \times 10^{-1}$	12	$(.139 \times 10^{-2})$
DENSCHNA	2	$1 \times 10^{-5}$ $1 \times 10^{-5}$	18	$2.990 \times 10^{-7}$	00	$1.190 \times 10^{-5}$ $1.087 \times 10^{-5}$	387	$1.373 \times 10^{-9}$
DENSORINA	2	$1 \times 10^{-7}$ $1 \times 10^{-7}$	22	$3.337 \times 10^{-10}$		$1.007 \times 10^{-8}$	02	$0.002 \times 10^{-11}$
DENSORINA	2	1 × 10 '	22	$1.233 \times 10^{-10}$	44	$0.991 \times 10^{-5}$	101	$1.405 \times 10^{-11}$

Table A.7. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

A.3.3.2. Noiseless Functions. In this section, we list the final objective value( $\phi(x)$ ), number of function evaluations (#feval), CPU time (CPU), and feasibility error(feaserr) for each problem instance in Tables A.28-A.29.

			N	IEWUOA	FI	) L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
DENSCHNB	2	$1 \times 10^{-1}$	14	$2.282\times 10^{-2}$	23	$2.369\times 10^{-1}$	293	$4.982\times 10^{-4}$
DENSCHNB	2	$1 \times 10^{-3}$	17	$1.807 \times 10^{-5}$	40	$8.803\times10^{-4}$	387	$1.432\times 10^{-6}$
DENSCHNB	2	$1 \times 10^{-5}$	27	$2.693\times 10^{-7}$	61	$1.017\times 10^{-6}$	42	$5.983 \times 10^{-11}$
DENSCHNB	2	$1 \times 10^{-7}$	26	$7.582 \times 10^{-9}$	32	$2.194\times10^{-8}$	136	$9.498 \times 10^{-12}$
DENSCHNC	2	$1 \times 10^{-1}$	30	$-7.755 \times 10^{-2}$	18	$5.033  imes 10^{-2}$	281	$3.323\times 10^{-2}$
DENSCHNC	2	$1 \times 10^{-3}$	62	$2.780 \times 10^{-5}$	80	$1.775\times10^{-2}$	219	$2.208\times10^{-5}$
DENSCHNC	2	$1 \times 10^{-5}$	55	$-1.834 \times 10^{-1}$	160	$6.362 \times 10^{-6}$	387	$1.055\times10^{-8}$
DENSCHNC	2	$1 \times 10^{-7}$	54	$-1.834\times10^{-1}$	218	$2.200\times10^{-6}$	136	$-1.285 \times 10^{-10}$
DENSCHND	3	$1 \times 10^{-1}$	142	$3.571  imes 10^1$	269	$3.030  imes 10^1$	228	$1.318\times 10^{-3}$
DENSCHND	3	$1 \times 10^{-3}$	279	$2.258\times 10^{-3}$	216	$1.288 \times 10^1$	414	$4.836 \times 10^{-5}$
DENSCHND	3	$1 \times 10^{-5}$	216	$2.909\times 10^{-4}$	240	$2.218\times 10^{-5}$	326	$7.702 \times 10^{-6}$
DENSCHND	3	$1 \times 10^{-7}$	319	$1.084\times 10^{-6}$	432	$1.318\times 10^{-7}$	679	$8.619\times 10^{-8}$
DENSCHNE	3	$1 \times 10^{-1}$	29	1.570	44	1.003	66	1.038
DENSCHNE	3	$1 \times 10^{-3}$	42	1.003	132	$9.994 \times 10^{-1}$	138	$9.993  imes 10^{-1}$
DENSCHNE	3	$1 \times 10^{-5}$	73	$9.999  imes 10^{-1}$	75	$9.993 imes10^{-1}$	389	$9.993\times10^{-1}$
DENSCHNE	3	$1 \times 10^{-7}$	116	$1.080 \times 10^{-8}$	346	$9.993  imes 10^{-1}$	295	$1.651 \times 10^{-10}$
DENSCHNF	2	$1 \times 10^{-1}$	19	$3.023 \times 10^{-3}$	199	$1.909\times10^{-2}$	100	$2.349\times10^{-3}$
DENSCHNF	2	$1 \times 10^{-3}$	26	$1.388 \times 10^{-5}$	80	$2.299\times 10^{-4}$	100	$8.045\times10^{-6}$
DENSCHNF	2	$1 \times 10^{-5}$	32	$2.913\times10^{-8}$	59	$6.314 \times 10^{-7}$	77	$2.843\times10^{-8}$
DENSCHNF	2	$1 \times 10^{-7}$	38	$2.403\times10^{-9}$	58	$8.740 \times 10^{-10}$	136	$4.048 \times 10^{-11}$
DIXMAANA	15	$1 \times 10^{-1}$	202	2.011	278	3.989	727	$3.807 \times 10^{-3}$
DIXMAANA	15	$1 \times 10^{-3}$	401	$1.671\times 10^{-2}$	426	$4.111 \times 10^{-3}$	516	$7.512\times10^{-6}$
DIXMAANA	15	$1 \times 10^{-5}$	564	$5.826 \times 10^{-5}$	291	$2.041\times10^{-5}$	413	$2.145\times10^{-8}$
DIXMAANA	15	$1 \times 10^{-7}$	650	$7.482\times10^{-7}$	170	$1.288\times 10^{-6}$	660	$5.438 \times 10^{-11}$
DIXMAANA	90	$1 \times 10^{-1}$	1357	$9.089 \times 10^{-1}$	1596	$1.935  imes 10^1$	4121	$1.338\times 10^{-2}$
DIXMAANA	90	$1 \times 10^{-3}$	1820	$1.564\times 10^{-2}$	6783	$4.443\times10^{-2}$	3657	$3.093  imes 10^{-5}$
DIXMAANA	90	$1 \times 10^{-5}$	2384	$2.123\times 10^{-4}$	3552	$2.365\times10^{-4}$	3657	$1.304\times 10^{-7}$
DIXMAANA	90	$1 \times 10^{-7}$	2302	$1.072\times 10^{-6}$	2031	$1.824\times10^{-6}$	4663	$1.690 \times 10^{-10}$
DIXMAANA	300	$1 \times 10^{-1}$	11454	2.772	2091	$5.752  imes 10^1$	13555	$6.279\times10^{-2}$
DIXMAANA	300	$1 \times 10^{-3}$	12409	$3.370\times10^{-2}$	5739	$4.750\times10^{-2}$	7272	$1.397  imes 10^{-4}$
DIXMAANA	300	$1 \times 10^{-5}$	21687	$2.406\times 10^{-4}$	4530	$6.880\times10^{-4}$	17895	$3.066\times 10^{-7}$
DIXMAANA	300	$1 \times 10^{-7}$	18612	$3.404\times10^{-6}$	5134	$6.485\times10^{-6}$	7860	$8.727 \times 10^{-10}$
DIXMAANB	15	$1 \times 10^{-1}$	199	1.796	136	2.402	278	$6.595  imes 10^{-3}$
DIXMAANB	15	$1 \times 10^{-3}$	384	$1.563\times10^{-2}$	700	$7.149  imes 10^{-3}$	270	$1.424\times10^{-5}$
DIXMAANB	15	$1 \times 10^{-5}$	389	$1.644\times 10^{-5}$	351	$5.285  imes 10^{-5}$	505	$1.552\times 10^{-8}$
DIXMAANB	15	$1 \times 10^{-7}$	373	$2.098\times 10^{-7}$	274	$6.291  imes 10^{-7}$	491	$5.018\times10^{-10}$
DIXMAANB	90	$1 \times 10^{-1}$	1588	1.101	277	$1.692  imes 10^1$	1880	$2.910\times 10^{-2}$
DIXMAANB	90	$1 \times 10^{-3}$	1406	$1.029\times 10^{-2}$	4163	$1.213\times10^{-2}$	1101	$5.186 imes10^{-5}$
DIXMAANB	90	$1 \times 10^{-5}$	1918	$2.728\times 10^{-4}$	1568	$9.443\times10^{-5}$	2937	$2.263\times 10^{-7}$
DIXMAANB	90	$1 \times 10^{-7}$	1932	$1.523\times 10^{-6}$	2282	$5.621\times10^{-7}$	2937	$1.318\times 10^{-9}$
DIXMAANB	300	$1 \times 10^{-1}$	7971	2.504	905	$5.660  imes 10^1$	11680	$5.947\times10^{-2}$
DIXMAANB	300	$1 \times 10^{-3}$	12927	$3.555\times 10^{-2}$	11239	$9.529\times10^{-2}$	7371	$2.935\times10^{-4}$
DIXMAANB	300	$1 \times 10^{-5}$	14656	$2.960\times 10^{-4}$	8463	$2.769\times10^{-3}$	14914	$1.849\times10^{-6}$
DIXMAANB	300	$1 \times 10^{-7}$	19729	$3.256\times 10^{-6}$	5138	$2.711\times 10^{-5}$	7282	$1.165\times 10^{-9}$

Table A.8. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

We now comment on outcomes (ii)-(iv) mentioned in Section 2.4.1.

			N]	EWUOA	FD	L-BFGS	CI	) L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
DIXMAANC	15	$1 \times 10^{-1}$	123	$3.678\times10^{-1}$	696	$9.417 \times 10^{-1}$	413	$4.277 \times 10^{-3}$
DIXMAANC	15	$1 \times 10^{-3}$	214	$1.694  imes 10^{-3}$	784	$9.657  imes 10^{-3}$	627	$7.834 \times 10^{-5}$
DIXMAANC	15	$1 \times 10^{-5}$	252	$1.523 \times 10^{-4}$	621	$3.658 \times 10^{-5}$	413	$2.695 \times 10^{-8}$
DIXMAANC	15	$1 \times 10^{-7}$	249	$2.681\times 10^{-7}$	528	$4.150 \times 10^{-8}$	627	$1.283\times10^{-10}$
DIXMAANC	90	$1 \times 10^{-1}$	853	1.678	2077	3.626	2649	$1.112\times 10^{-2}$
DIXMAANC	90	$1 \times 10^{-3}$	2546	$1.886\times 10^{-2}$	1016	$6.657\times 10^{-2}$	1880	$4.973  imes 10^{-5}$
DIXMAANC	90	$1 \times 10^{-5}$	2245	$1.953\times 10^{-4}$	2971	$4.662 \times 10^{-4}$	1880	$1.398 \times 10^{-7}$
DIXMAANC	90	$1 \times 10^{-7}$	2873	$2.169 \times 10^{-6}$	2869	$9.707 \times 10^{-6}$	5905	$1.894 \times 10^{-10}$
DIXMAANC	300	$1 \times 10^{-1}$	5706	2.660	908	$2.909  imes 10^1$	12855	$5.278  imes 10^{-2}$
DIXMAANC	300	$1 \times 10^{-3}$	15396	$2.200\times10^{-2}$	9990	$1.601 \times 10^{-1}$	13555	$1.924 \times 10^{-4}$
DIXMAANC	300	$1 \times 10^{-5}$	17952	$3.248 \times 10^{-4}$	8167	$2.040\times 10^{-4}$	9111	$2.439\times10^{-7}$
DIXMAANC	300	$1 \times 10^{-7}$	20832	$3.234 \times 10^{-6}$	7261	$1.070 \times 10^{-5}$	6639	$9.021 \times 10^{-10}$
DIXMAAND	15	$1 \times 10^{-1}$	197	2.209	71	2.464	394	$1.834\times 10^{-2}$
DIXMAAND	15	$1 \times 10^{-3}$	189	$1.568  imes 10^{-3}$	210	$3.272 \times 10^{-3}$	1730	$6.885 \times 10^{-6}$
DIXMAAND	15	$1 \times 10^{-5}$	265	$2.640\times 10^{-5}$	517	$1.972\times 10^{-4}$	660	$1.423  imes 10^{-7}$
DIXMAAND	15	$1 \times 10^{-7}$	245	$3.120 \times 10^{-7}$	583	$1.876\times10^{-6}$	1730	$6.047\times10^{-11}$
DIXMAAND	90	$1 \times 10^{-1}$	1468	1.597	647	4.471	1880	$2.476\times 10^{-2}$
DIXMAAND	90	$1 \times 10^{-3}$	2428	$2.428\times 10^{-2}$	2637	$6.034 \times 10^{-2}$	3075	$1.075  imes 10^{-4}$
DIXMAAND	90	$1 \times 10^{-5}$	1511	$2.151\times 10^{-4}$	5883	$2.672\times10^{-4}$	2449	$2.126 \times 10^{-7}$
DIXMAAND	90	$1 \times 10^{-7}$	2339	$1.634\times 10^{-6}$	2957	$1.455\times10^{-6}$	3970	$4.576 \times 10^{-10}$
DIXMAAND	300	$1 \times 10^{-1}$	6087	3.295	1211	$2.487 \times 10^1$	11680	$6.232 \times 10^{-2}$
DIXMAAND	300	$1 \times 10^{-3}$	28875	$4.695\times 10^{-2}$	15937	$2.861\times 10^{-2}$	13555	$2.324 \times 10^{-4}$
DIXMAAND	300	$1 \times 10^{-5}$	36583	$1.968\times 10^{-4}$	7259	$1.237 \times 10^{-3}$	6695	$8.301 \times 10^{-7}$
DIXMAAND	300	$1 \times 10^{-7}$	31821	$2.379\times 10^{-6}$	6958	$1.154\times 10^{-5}$	7282	$1.251 \times 10^{-9}$
DIXMAANE	15	$1 \times 10^{-1}$	117	2.680	228	2.593	516	$4.750 \times 10^{-2}$
DIXMAANE	15	$1 \times 10^{-3}$	483	$6.719\times10^{-2}$	907	$2.567\times 10^{-3}$	759	$9.723  imes 10^{-5}$
DIXMAANE	15	$1 \times 10^{-5}$	359	$6.662  imes 10^{-5}$	496	$4.389 \times 10^{-5}$	875	$3.665  imes 10^{-8}$
DIXMAANE	15	$1 \times 10^{-7}$	473	$5.981 \times 10^{-8}$	515	$4.413\times10^{-7}$	1730	$2.043 \times 10^{-10}$
DIXMAANE	90	$1 \times 10^{-1}$	957	2.077	2228	$1.048 \times 10^1$	5905	$1.522 \times 10^{-1}$
DIXMAANE	90	$1 \times 10^{-3}$	2231	$5.295\times10^{-2}$	2558	$5.724\times10^{-2}$	7854	$2.745 \times 10^{-3}$
DIXMAANE	90	$1 \times 10^{-5}$	2422	$1.047  imes 10^{-3}$	3592	$1.658\times 10^{-3}$	14178	$2.990 \times 10^{-6}$
DIXMAANE	90	$1 \times 10^{-7}$	3565	$3.468 \times 10^{-6}$	6643	$5.020 \times 10^{-6}$	9472	$1.434 \times 10^{-8}$
DIXMAANE	300	$1 \times 10^{-1}$	5797	6.472	5878	$1.576  imes 10^1$	17895	$5.082 \times 10^{-1}$
DIXMAANE	300	$1 \times 10^{-3}$	39321	$1.160\times10^{-1}$	11096	$1.216\times 10^{-1}$	42978	$5.040 \times 10^{-3}$
DIXMAANE	300	$1 \times 10^{-5}$	35575	$3.907  imes 10^{-3}$	16015	$3.572  imes 10^{-3}$	41288	$6.960  imes 10^{-5}$
DIXMAANE	300	$1 \times 10^{-7}$	29189	$1.864\times 10^{-5}$	27787	$3.545\times 10^{-5}$	51466	$5.905  imes 10^{-8}$
DIXMAANF	15	$1 \times 10^{-1}$	186	1.292	133	$9.307  imes 10^{-1}$	394	$8.259\times10^{-2}$
DIXMAANF	15	$1 \times 10^{-3}$	189	$1.051\times 10^{-2}$	409	$5.640  imes 10^{-3}$	413	$3.139  imes 10^{-5}$
DIXMAANF	15	$1 \times 10^{-5}$	296	$2.329\times 10^{-5}$	574	$1.101\times 10^{-4}$	825	$8.220 \times 10^{-8}$
DIXMAANF	15	$1 \times 10^{-7}$	351	$2.811\times 10^{-7}$	546	$1.188 \times 10^{-6}$	1730	$1.116\times10^{-10}$
DIXMAANF	90	$1 \times 10^{-1}$	1055	1.623	277	6.676	3143	$7.567 \times 10^{-2}$
DIXMAANF	90	$1 \times 10^{-3}$	1855	$1.150  imes 10^{-1}$	2279	$3.255\times 10^{-2}$	6862	$1.656\times 10^{-3}$
DIXMAANF	90	$1 \times 10^{-5}$	3117	$2.379\times 10^{-4}$	4615	$5.649 \times 10^{-4}$	6955	$1.941 \times 10^{-6}$
DIXMAANF	90	$1 \times 10^{-7}$	6296	$1.248\times 10^{-5}$	4327	$6.372\times 10^{-6}$	9524	$6.810\times10^{-9}$
DIXMAANF	300	$1 \times 10^{-1}$	4652	2.581	904	$2.266 \times 10^1$	11680	$1.549\times10^{-1}$
DIXMAANF	300	$1 \times 10^{-3}$	10797	$1.173\times 10^{-1}$	4840	$2.153\times 10^{-1}$	31414	$5.266\times 10^{-3}$
DIXMAANF	300	$1 \times 10^{-5}$	41760	$2.501\times 10^{-3}$	15493	$6.550\times 10^{-3}$	41516	$2.248\times 10^{-5}$
DIXMAANF	300	$1 \times 10^{-7}$	17405	$3.175\times 10^{-5}$	22658	$3.595\times 10^{-5}$	59871	$2.087\times 10^{-8}$
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Table A.9. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			NI	EWUOA	FD	L-BFGS	CD	L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
DIXMAANG	15	$1 \times 10^{-1}$	118	1.301	53	$8.760\times10^{-1}$	1730	$2.446\times 10^{-2}$
DIXMAANG	15	$1 \times 10^{-3}$	219	$3.590  imes 10^{-3}$	532	$8.830 \times 10^{-3}$	1104	$4.679\times 10^{-5}$
DIXMAANG	15	$1 \times 10^{-5}$	259	$1.628 \times 10^{-5}$	310	$9.965  imes 10^{-4}$	1730	$9.830  imes 10^{-8}$
DIXMAANG	15	$1 \times 10^{-7}$	354	$1.419 \times 10^{-7}$	1235	$2.106 \times 10^{-6}$	1730	$3.410\times10^{-10}$
DIXMAANG	90	$1 \times 10^{-1}$	1253	1.630	2284	1.752	4496	$9.061 \times 10^{-2}$
DIXMAANG	90	$1 \times 10^{-3}$	3690	$9.785 \times 10^{-2}$	2011	$3.834 \times 10^{-2}$	6194	$1.981 \times 10^{-3}$
DIXMAANG	90	$1 \times 10^{-5}$	2545	$3.136  imes 10^{-4}$	2579	$8.481\times 10^{-4}$	6124	$7.697\times10^{-6}$
DIXMAANG	90	$1 \times 10^{-7}$	3047	$8.060 \times 10^{-6}$	7112	$6.479 \times 10^{-6}$	8861	$7.244 \times 10^{-9}$
DIXMAANG	300	$1 \times 10^{-1}$	7837	3.509	908	$1.259 \times 10^{1}$	12260	$3.429  imes 10^{-1}$
DIXMAANG	300	$1 \times 10^{-3}$	40449	$1.096 \times 10^{-1}$	9353	$2.259\times 10^{-1}$	17101	$4.560 \times 10^{-3}$
DIXMAANG	300	$1 \times 10^{-5}$	33893	$3.070 \times 10^{-3}$	8463	$2.846 \times 10^{-3}$	35863	$3.047 \times 10^{-5}$
DIXMAANG	300	$1 \times 10^{-7}$	149998	$4.680 \times 10^{-4}$	21147	$6.501 \times 10^{-5}$	44939	$9.867 \times 10^{-8}$
DIXMAANH	15	$1 \times 10^{-1}$	167	$5.656 \times 10^{-1}$	54	$9.557 \times 10^{-1}$	394	$2.593 \times 10^{-2}$
DIXMAANH	15	$1 \times 10^{-3}$	280	$7.738 \times 10^{-3}$	395	$4.275 \times 10^{-3}$	705	$3.527 \times 10^{-5}$
DIXMAANH	15	$1 \times 10^{-5}$	365	$3.603 \times 10^{-5}$	419	$6.578 \times 10^{-5}$	825	$7.652 \times 10^{-8}$
DIXMAANH	15	$1 \times 10^{-7}$	400	$1.836 \times 10^{-7}$	767	$2.555\times 10^{-7}$	830	$9.052\times10^{-10}$
DIXMAANH	90	$1 \times 10^{-1}$	1598	2.375	963	1.997	3163	$8.808 \times 10^{-2}$
DIXMAANH	90	$1 \times 10^{-3}$	2287	$1.070 \times 10^{-1}$	3112	$9.064 \times 10^{-2}$	8780	$3.029 \times 10^{-3}$
DIXMAANH	90	$1 \times 10^{-5}$	3252	$2.565 \times 10^{-4}$	5492	$8.586 \times 10^{-4}$	6194	$4.423 \times 10^{-6}$
DIXMAANH	90	$1 \times 10^{-7}$	5217	$3.169 \times 10^{-6}$	4430	$6.061 \times 10^{-6}$	8780	$4.295 \times 10^{-9}$
DIXMAANH	300	$1 \times 10^{-1}$	19419	5.492	1211	$1.205 \times 10^1$	6759	$3.222 \times 10^{-1}$
DIXMAANH	300	$1 \times 10^{-3}$	33582	$2.678\times10^{-1}$	15009	$1.093 \times 10^{-1}$	17702	$4.861 \times 10^{-3}$
DIXMAANH	300	$1 \times 10^{-5}$	25154	$4.098 \times 10^{-3}$	15460	$4.812 \times 10^{-3}$	40314	$1.688 \times 10^{-5}$
DIXMAANH	300	$1 \times 10^{-7}$	25753	$1.920 \times 10^{-5}$	22065	$2.289 \times 10^{-5}$	41288	$8.904 \times 10^{-8}$
DIXMAANI	15	$1 \times 10^{-1}$	113	$5.926 \times 10^{-1}$	1218	1.032	516	$3.909  imes 10^{-2}$
DIXMAANI	15	$1 \times 10^{-3}$	358	$6.615  imes 10^{-3}$	504	$1.797 \times 10^{-2}$	1053	$2.345 \times 10^{-3}$
DIXMAANI	15	$1 \times 10^{-5}$	467	$6.109 \times 10^{-4}$	1343	$3.261 \times 10^{-4}$	1476	$1.374 \times 10^{-6}$
DIXMAANI	15	$1 \times 10^{-7}$	636	$6.946 \times 10^{-6}$	792	$1.308 \times 10^{-6}$	1611	$4.603 \times 10^{-9}$
DIXMAANI	90	$1 \times 10^{-1}$	951	2.135	3000	4.382	7696	$2.633 \times 10^{-1}$
DIXMAANI	90	$1 \times 10^{-3}$	1824	$4.215\times 10^{-2}$	2824	$5.241\times10^{-2}$	7466	$5.757 imes10^{-3}$
DIXMAANI	90	$1 \times 10^{-5}$	3278	$7.114 \times 10^{-3}$	5618	$4.747 \times 10^{-3}$	18792	$1.816  imes 10^{-4}$
DIXMAANI	90	$1 \times 10^{-7}$	10692	$2.134\times 10^{-4}$	12783	$1.281\times 10^{-4}$	36022	$5.632  imes 10^{-7}$
DIXMAANI	300	$1 \times 10^{-1}$	7480	5.731	12317	$1.749  imes 10^1$	21554	$9.401  imes 10^{-1}$
DIXMAANI	300	$1 \times 10^{-3}$	12194	$1.431\times 10^{-1}$	10792	$1.467\times 10^{-1}$	43977	$1.445\times 10^{-2}$
DIXMAANI	300	$1 \times 10^{-5}$	25477	$9.917 imes10^{-3}$	34341	$6.574 imes10^{-3}$	70880	$3.511 imes10^{-4}$
DIXMAANI	300	$1 \times 10^{-7}$	36578	$1.928\times 10^{-4}$	74070	$2.406\times 10^{-4}$	150151*	$1.779\times10^{-5}$
DIXMAANJ	15	$1 \times 10^{-1}$	156	$1.786\times10^{-1}$	52	$6.219 imes10^{-1}$	1124	$3.543\times10^{-2}$
DIXMAANJ	15	$1 \times 10^{-3}$	309	$2.567\times 10^{-2}$	584	$1.557\times 10^{-2}$	1536	$7.979\times10^{-4}$
DIXMAANJ	15	$1 \times 10^{-5}$	672	$7.838  imes 10^{-5}$	835	$2.748\times10^{-4}$	1124	$7.681\times10^{-7}$
DIXMAANJ	15	$1 \times 10^{-7}$	685	$4.621\times 10^{-7}$	789	$3.246  imes 10^{-5}$	2031	$1.804\times10^{-8}$
DIXMAANJ	90	$1 \times 10^{-1}$	1141	1.187	277	3.668	1880	$7.149\times10^{-2}$
DIXMAANJ	90	$1 \times 10^{-3}$	2186	$4.912\times10^{-2}$	2853	$5.441 \times 10^{-2}$	7696	$2.858\times10^{-3}$
DIXMAANJ	90	$1 \times 10^{-5}$	4146	$3.591\times 10^{-3}$	3791	$3.090\times 10^{-3}$	11399	$4.242\times10^{-5}$
DIXMAANJ	90	$1  imes 10^{-7}$	13470	$6.305\times10^{-4}$	7251	$3.692\times 10^{-5}$	29130	$2.883\times 10^{-7}$
DIXMAANJ	300	$1 \times 10^{-1}$	4325	2.391	904	$1.236  imes 10^1$	11680	$1.834\times10^{-1}$
DIXMAANJ	300	$1  imes 10^{-3}$	23360	$1.285\times 10^{-1}$	8042	$1.123\times 10^{-1}$	36345	$8.119\times 10^{-3}$
DIXMAANJ	300	$1 \times 10^{-5}$	19573	$3.531\times 10^{-3}$	18352	$5.437\times10^{-3}$	41369	$5.491\times10^{-5}$
DIXMAANJ	300	$1 \times 10^{-7}$	31059	$1.801\times 10^{-4}$	33637	$5.060\times10^{-5}$	84610	$2.777\times10^{-6}$

Table A.10. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			NI	EWUOA	FD	L-BFGS	CI	) L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
DIXMAANK	15	$1 \times 10^{-1}$	123	$2.036\times 10^{-1}$	70	$3.310\times 10^{-1}$	270	$5.345  imes 10^{-2}$
DIXMAANK	15	$1 \times 10^{-3}$	226	$4.148\times 10^{-2}$	1229	$1.660\times 10^{-2}$	1024	$8.734  imes 10^{-4}$
DIXMAANK	15	$1 \times 10^{-5}$	533	$4.054\times10^{-4}$	759	$7.894 \times 10^{-5}$	1730	$3.452\times10^{-7}$
DIXMAANK	15	$1 \times 10^{-7}$	710	$1.111 \times 10^{-6}$	724	$3.219  imes 10^{-7}$	1595	$4.578 \times 10^{-9}$
DIXMAANK	90	$1 \times 10^{-1}$	962	2.115	961	$9.536  imes 10^{-1}$	1880	$3.049 \times 10^{-1}$
DIXMAANK	90	$1 \times 10^{-3}$	1921	$4.371 \times 10^{-2}$	2718	$9.449 \times 10^{-2}$	10211	$3.178 \times 10^{-3}$
DIXMAANK	90	$1 \times 10^{-5}$	3646	$8.827\times 10^{-4}$	7155	$1.234 \times 10^{-3}$	8157	$3.217 \times 10^{-5}$
DIXMAANK	90	$1 \times 10^{-7}$	6961	$5.780 \times 10^{-4}$	5158	$2.208 \times 10^{-5}$	14626	$1.152 \times 10^{-6}$
DIXMAANK	300	$1 \times 10^{-1}$	4131	3.657	908	6.765	12777	$4.528 \times 10^{-1}$
DIXMAANK	300	$1 \times 10^{-3}$	30156	$1.396  imes 10^{-1}$	10483	$3.126  imes 10^{-1}$	21925	$3.813  imes 10^{-3}$
DIXMAANK	300	$1 \times 10^{-5}$	17158	$7.029 \times 10^{-3}$	16876	$3.556 \times 10^{-3}$	26675	$6.708 \times 10^{-5}$
DIXMAANK	300	$1 \times 10^{-7}$	52657	$1.527\times 10^{-4}$	15408	$7.813\times10^{-5}$	82360	$6.627\times 10^{-7}$
DIXMAANL	15	$1 \times 10^{-1}$	152	1.451	88	$5.246  imes 10^{-1}$	1246	$5.972 \times 10^{-2}$
DIXMAANL	15	$1 \times 10^{-3}$	370	$3.861\times 10^{-2}$	866	$1.485\times10^{-2}$	1694	$5.900 \times 10^{-4}$
DIXMAANL	15	$1 \times 10^{-5}$	782	$6.330\times 10^{-5}$	1050	$1.029\times 10^{-4}$	1476	$1.103\times10^{-6}$
DIXMAANL	15	$1 \times 10^{-7}$	873	$1.152\times 10^{-6}$	638	$1.039\times10^{-6}$	1730	$2.064\times 10^{-8}$
DIXMAANL	90	$1 \times 10^{-1}$	1230	1.224	1363	1.218	2817	$1.846\times10^{-1}$
DIXMAANL	90	$1 \times 10^{-3}$	2191	$4.680\times 10^{-2}$	2539	$1.429\times 10^{-1}$	5193	$2.433\times 10^{-3}$
DIXMAANL	90	$1 \times 10^{-5}$	3048	$1.131 \times 10^{-3}$	3152	$2.455 \times 10^{-3}$	8780	$3.050 \times 10^{-5}$
DIXMAANL	90	$1 \times 10^{-7}$	7705	$2.888 \times 10^{-4}$	5990	$1.848 \times 10^{-5}$	32927	$1.567 \times 10^{-7}$
DIXMAANL	300	$1 \times 10^{-1}$	7660	3.606	1211	6.900	11674	$1.739 \times 10^{-1}$
DIXMAANL	300	$1 \times 10^{-3}$	33194	$1.813 \times 10^{-1}$	11079	$2.767 \times 10^{-1}$	19562	$4.733 \times 10^{-3}$
DIXMAANL	300	$1 \times 10^{-5}$	149998	$3.705\times10^{-2}$	10571	$5.175  imes 10^{-3}$	26767	$4.075\times10^{-5}$
DIXMAANL	300	$1 \times 10^{-7}$	61830	$7.296\times10^{-4}$	24171	$5.364\times10^{-5}$	57101	$1.891\times10^{-6}$
DQRTIC	10	$1 \times 10^{-1}$	139	$1.786 \times 10^{-1}$	927	$1.449 \times 10^{-1}$	762	$2.816 \times 10^{-4}$
DORTIC	10	$1 \times 10^{-3}$	152	$4.676\times 10^{-3}$	304	$6.109  imes 10^{-3}$	749	$3.770 \times 10^{-5}$
DQRTIC	10	$1 \times 10^{-5}$	429	$8.020 \times 10^{-5}$	943	$3.905  imes 10^{-5}$	675	$8.733 \times 10^{-9}$
DQRTIC	10	$1 \times 10^{-7}$	405	$3.782\times10^{-8}$	676	$1.675\times 10^{-6}$	675	$1.032 \times 10^{-8}$
DORTIC	50	$1 \times 10^{-1}$	1672	$7.132\times10^{-1}$	1197	$8.681\times 10^{-1}$	2857	$7.958\times10^{-2}$
DQRTIC	50	$1 \times 10^{-3}$	1719	$7.324 \times 10^{-3}$	2008	$7.825 \times 10^{-3}$	2651	$3.038 \times 10^{-4}$
DORTIC	50	$1 \times 10^{-5}$	2369	$2.603\times10^{-5}$	2629	$1.445\times 10^{-4}$	3341	$1.758 \times 10^{-6}$
DORTIC	50	$1 \times 10^{-7}$	2435	$1.473\times 10^{-6}$	2493	$1.915\times 10^{-6}$	3341	$2.992\times 10^{-8}$
DQRTIC	100	$1 \times 10^{-1}$	4920	$3.885 \times 10^{-1}$	4766	6.422	7146	$4.248 \times 10^{-2}$
DORTIC	100	$1 \times 10^{-3}$	5612	$9.569  imes 10^{-3}$	4517	$2.474\times 10^{-3}$	5960	$1.807 \times 10^{-4}$
DQRTIC	100	$1 \times 10^{-5}$	6494	$5.767  imes 10^{-5}$	6534	$7.941\times10^{-4}$	7146	$2.060\times 10^{-6}$
DQRTIC	100	$1 \times 10^{-7}$	6913	$8.633\times 10^{-7}$	1119	$1.906 \times 10^7$	7344	$2.041 \times 10^{-8}$
EDENSCH	36	$1 \times 10^{-1}$	873	$9.721  imes 10^{-1}$	1071	1.223	2104	$1.290 \times 10^{-1}$
EDENSCH	36	$1 \times 10^{-3}$	1175	$1.187\times 10^{-2}$	2031	$1.116\times 10^{-2}$	1772	$3.795  imes 10^{-4}$
EDENSCH	36	$1 \times 10^{-5}$	1166	$1.001  imes 10^{-4}$	1324	$5.176 imes10^{-4}$	1772	$8.173\times10^{-7}$
EDENSCH	36	$1 \times 10^{-7}$	1331	$1.083\times 10^{-6}$	1300	$3.572\times 10^{-6}$	1824	$2.417\times 10^{-9}$
EIGENALS	6	$1 \times 10^{-1}$	5**	0.000	96	$5.364 \times 10^{-1}$	509	$2.108 \times 10^{-1}$
EIGENALS	6	$1 \times 10^{-3}$	5**	0.000	82	$1.041\times 10^{-2}$	252	$5.299 \times 10^{-4}$
EIGENALS	6	$1 \times 10^{-5}$	5**	0.000	250	$4.298\times 10^{-4}$	152	$1.284\times 10^{-7}$
EIGENALS	6	$1 \times 10^{-7}$	5**	0.000	294	$5.331 \times 10^{-6}$	395	$1.289 \times 10^{-10}$
EIGENALS	110	$1 \times 10^{-1}$	2393	9.021	3176	$1.692  imes 10^1$	8820	$1.404\times 10^{-1}$
EIGENALS	110	$1 \times 10^{-3}$	6471	$8.081\times 10^{-1}$	3365	$2.275\times 10^{-1}$	5137	$1.253\times 10^{-2}$
EIGENALS	110	$1 \times 10^{-5}$	36272	$2.114\times 10^{-2}$	6187	$1.407\times 10^{-2}$	41789	$2.313\times10^{-4}$
EIGENALS	110	$1  imes 10^{-7}$	54998	$2.423\times 10^{-4}$	32572	$1.418\times 10^{-4}$	$55103^{*}$	$2.471\times 10^{-5}$

Table A.11. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			N	EWUOA	FD	L-BFGS	CI	) L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
EIGENBLS <sup>‡</sup>	6	$1 \times 10^{-1}$	51	$1.421\times 10^{-1}$	32	1.096	252	$3.986\times 10^{-3}$
EIGENBLS <sup>‡</sup>	6	$1 \times 10^{-3}$	109	$3.843\times 10^{-2}$	240	$1.223\times 10^{-3}$	376	$4.336\times 10^{-5}$
EIGENBLS <sup>‡</sup>	6	$1 \times 10^{-5}$	217	$-1.840 \times 10^{-1}$	247	$2.792\times10^{-5}$	299	$1.340 \times 10^{-7}$
EIGENBLS <sup>‡</sup>	6	$1 \times 10^{-7}$	307	$-1.849\times10^{-1}$	190	$9.576\times10^{-8}$	256	$5.404\times10^{-10}$
EIGENBLS	110	$1 \times 10^{-1}$	705	3.389	1761	$1.122 \times 10^1$	4023	1.707
EIGENBLS	110	$1 \times 10^{-3}$	5295	1.452	2754	1.596	7292	1.552
EIGENBLS	110	$1 \times 10^{-5}$	18185	$2.700\times10^{-2}$	21793	$2.398\times10^{-2}$	55096*	$5.814\times10^{-3}$
EIGENBLS	110	$1 \times 10^{-7}$	30763	$5.299 \times 10^{-4}$	46722	$9.959  imes 10^{-4}$	55088*	$2.591\times10^{-2}$
EIGENCLS	30	$1 \times 10^{-1}$	400	1.610	1107	$1.185  imes 10^1$	1812	$4.235\times10^{-1}$
EIGENCLS	30	$1 \times 10^{-3}$	686	$3.905  imes 10^{-1}$	1546	$2.774\times10^{-2}$	3424	$1.360 \times 10^{-3}$
EIGENCLS	30	$1 \times 10^{-5}$	1590	$1.176 \times 10^{-3}$	1935	$1.016 \times 10^{-3}$	4427	$1.116 \times 10^{-5}$
EIGENCLS	30	$1 \times 10^{-7}$	2212	$6.465 \times 10^{-6}$	3159	$1.266 \times 10^{-5}$	6325	$3.770 imes10^{-8}$
ENGVAL1	2	$1 \times 10^{-1}$	22	$3.602 \times 10^{-2}$	24	$5.487 \times 10^{-1}$	146	$1.987 \times 10^{-2}$
ENGVAL1	2	$1 \times 10^{-3}$	30	$9.036 \times 10^{-5}$	203	$4.722\times10^{-3}$	146	$2.480\times10^{-5}$
ENGVAL1	2	$1 \times 10^{-5}$	33	$5.987 \times 10^{-8}$	54	$2.397 \times 10^{-6}$	91	$2.265\times10^{-8}$
ENGVAL1	2	$1 \times 10^{-7}$	36	$3.361 \times 10^{-9}$	56	$3.129  imes 10^{-8}$	91	$1.713 \times 10^{-11}$
ENGVAL1	50	$1 \times 10^{-1}$	1298	1.868	2612	6.147	1160	$2.136 \times 10^{-1}$
ENGVAL1	50	$1 \times 10^{-3}$	1125	$2.268 \times 10^{-2}$	634	$1.857 \times 10^{-2}$	3720	$2.678 \times 10^{-4}$
ENGVAL1	50	$1 \times 10^{-5}$	2643	$1.160 \times 10^{-4}$	894	$5.713 \times 10^{-5}$	1460	$2.359 \times 10^{-6}$
ENGVAL1	50	$1 \times 10^{-7}$	2802	$9.129 \times 10^{-7}$	1354	$1.628 \times 10^{-6}$	2651	$2.889 \times 10^{-9}$
ENGVAL1	100	$1 \times 10^{-1}$	2362	1.585	3555	$1.187 \times 10^{1}$	1638	$3.397 \times 10^{-1}$
ENGVAL1	100	$1 \times 10^{-3}$	3400	$1.462 \times 10^{-2}$	2805	$4.521 \times 10^{-2}$	4885	$8.850 \times 10^{-4}$
ENGVAL1	100	$1 \times 10^{-5}$	3501	$1.729 \times 10^{-4}$	1334	$5.313 \times 10^{-4}$	5960	$2.445 \times 10^{-6}$
ENGVAL1	100	$1 \times 10^{-7}$	3166	$1.293 \times 10^{-6}$	4913	$3.878 \times 10^{-6}$	4095	$6.088 \times 10^{-9}$
EXPFIT	2	$1 \times 10^{-1}$	17	3.207	54	4.954	126	$4.794 \times 10^{-2}$
EXPFIT	2	$1 \times 10^{-3}$	33	$3.946 \times 10^{-3}$	58	$4.299 \times 10^{-3}$	67	$4.256 \times 10^{-5}$
EXPFIT	2	$1 \times 10^{-5}$	38	$3.463 \times 10^{-7}$	107	$2.064 \times 10^{-6}$	260	$7.468 \times 10^{-10}$
EXPET	2	$1 \times 10^{-7}$	49	$8.778 \times 10^{-10}$	68	$2.677 \times 10^{-8}$	91	$5.019 \times 10^{-11}$
FLETCBV3*	10	$1 \times 10^{-1}$	21**	$3.228 \times 10^{-2}$	32	$3.228 \times 10^{-2}$	21	$3.228 \times 10^{-2}$
FLETCBV3*	10	$1 \times 10^{-5}$	21**	$3.228 \times 10^{-2}$	29	$3.228 \times 10^{-2}$	21	$3.228 \times 10^{-2}$
FLETCBV3*	10	$1 \times 10^{-3}$	22	$3.228 \times 10^{-2}$	27	$3.228 \times 10^{-2}$	21	$3.228 \times 10^{-2}$
FLETCBV3*	100	$1 \times 10^{-1}$	965	$8.774 \times 10^{-4}$	27	$3.228 \times 10^{-2}$	403	$3.228 \times 10^{-2}$
FLETCBV3*	100	$1 \times 10^{-1}$	202	$1.785 \times 10^{5}$	2486	$1.785 \times 10^{5}$	1952	$1.785 \times 10^{5}$
FLETCBV3*	100	$1 \times 10^{-5}$	202	$1.785 \times 10^{5}$	447	$1.785 \times 10^{5}$	4029	$1.785 \times 10^{3}$
FLETCBV3*	100	$1 \times 10^{-3}$	50000*	$1.349 \times 10^{3}$	11892	$1.785 \times 10^{5}$	49799	$2.028 \times 10^{2}$
FLETCBV3*	100	$1 \times 10^{-1}$	49999	$5.761 \times 10^{4}$	222	$1.785 \times 10^{\circ}$	50010*	$1.015 \times 10^{4}$
FLETCHBV*	10	$1 \times 10^{-1}$	863	$2.240 \times 10^{4}$	639	$2.653 \times 10^{3}$	1087	$1.218 \times 10^{4}$
FLETCHBV*	10	$1 \times 10^{-5}$	1289	$1.540 \times 10^{5}$	1118	$8.389 \times 10^{3}$	1356	$-4.764 \times 10^{3}$
FLETCHBV*	10	$1 \times 10^{-7}$	790	$2.228 \times 10^{6}$	708	$-4.605 \times 10^{3}$	1445	$-7.519 \times 10^{3}$
FLETCHBV*	100	$1 \times 10^{-1}$	977	$8.461 \times 10^{4}$	759	$(.233 \times 10^2)$	952	$4.230 \times 10^{2}$
FLETCHBV*	100	$1 \times 10^{-1}$	50000*	$1.581 \times 10^{13}$ 1.721 $10^{13}$	40962	$-1.187 \times 10^{9}$	50076*	$-0.387 \times 10^{9}$
FLETCHBV*	100	$1 \times 10^{-5}$ $1 \cdots 10^{-5}$	50000*	$1.731 \times 10^{13}$ 1.704 ··· 10^{13}	50015*	$-9.179 \times 10^{9}$	50085*	$-3.093 \times 10^{9}$
FLETCHBV*	100	$1 \times 10^{-7}$	50000*	$1.704 \times 10^{10}$ 1.201 $\times 10^{13}$	33119	$-8.084 \times 10^{-6}$	50068*	$-5.710 \times 10^{\circ}$
FLEICHBV <sup>+</sup>	100	$1 \times 10^{-7}$	50000*	$1.391 \times 10^{13}$	8369	$(.425 \times 10^{10})$	50080*	$-5.308 \times 10^{9}$

Table A.12. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

(ii) The solvers converged to feasible points with different objective function values. There are three such problems, for all of which COBYLA terminated due to the limit on

			N	EWUOA	FD	L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x)-\phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
FREUROTH	2	$1 \times 10^{-1}$	38	$6.304\times10^{-2}$	35	$5.015  imes 10^1$	368	$3.877 \times 10^{-4}$
FREUROTH	2	$1 \times 10^{-3}$	35	$5.507 \times 10^{-1}$	347	$1.660 \times 10^{-3}$	141	$1.496 \times 10^{-5}$
FREUROTH	2	$1 \times 10^{-5}$	70	$8.659\times 10^{-7}$	83	$9.381\times10^{-5}$	387	$3.302 \times 10^{-8}$
FREUROTH	2	$1 \times 10^{-7}$	73	$4.594\times10^{-10}$	138	$1.112\times 10^{-6}$	261	$4.996 \times 10^{-10}$
FREUROTH	10	$1 \times 10^{-1}$	125	$4.555 \times 10^1$	137	$5.548  imes 10^1$	1466	$8.708 \times 10^{-3}$
FREUROTH	10	$1 \times 10^{-3}$	371	$3.964\times10^{-3}$	467	$1.341\times 10^{-1}$	562	$4.686 \times 10^{-5}$
FREUROTH	10	$1 \times 10^{-5}$	433	$3.004 \times 10^{-4}$	295	$1.857 \times 10^{-3}$	842	$5.755 \times 10^{-8}$
FREUROTH	10	$1 \times 10^{-7}$	435	$1.512\times 10^{-6}$	602	$1.348\times 10^{-5}$	630	$-3.206 \times 10^{-10}$
FREUROTH	50	$1 \times 10^{-1}$	1629	$5.744 \times 10^{1}$	2466	$5.503  imes 10^1$	2651	$6.734\times10^{-2}$
FREUROTH	50	$1 \times 10^{-3}$	6822	6.137	3983	$3.723\times10^{-2}$	2857	$5.299 \times 10^{-5}$
FREUROTH	50	$1 \times 10^{-5}$	7662	$1.340\times10^{-3}$	5024	$3.387\times 10^{-4}$	2487	$2.101\times 10^{-7}$
FREUROTH	50	$1 \times 10^{-7}$	6619	$9.762 \times 10^{-6}$	1257	$1.453\times10^{-6}$	3063	$5.975 \times 10^{-10}$
FREUROTH	100	$1 \times 10^{-1}$	1868	$5.751  imes 10^1$	920	$6.619  imes 10^1$	8609	$6.268\times 10^{-2}$
FREUROTH	100	$1 \times 10^{-3}$	8308	9.577	3192	$5.430\times10^{-2}$	6050	$1.752\times 10^{-4}$
FREUROTH	100	$1 \times 10^{-5}$	5029	$1.824\times10^{-3}$	4095	$2.290\times 10^{-4}$	5769	$4.168 \times 10^{-7}$
FREUROTH	100	$1 \times 10^{-7}$	5293	$7.505\times10^{-6}$	6187	$1.098\times 10^{-5}$	5907	$-5.191\times10^{-9}$
GENROSE	5	$1 \times 10^{-1}$	35	3.849	68	3.246	180	2.711
GENROSE	5	$1 \times 10^{-3}$	162	$4.182\times10^{-2}$	234	2.566	446	$9.023  imes 10^{-4}$
GENROSE	5	$1 \times 10^{-5}$	221	$8.543\times10^{-5}$	260	$2.548\times 10^{-3}$	413	$2.278\times 10^{-7}$
GENROSE	5	$1 \times 10^{-7}$	261	$9.023  imes 10^{-7}$	317	$4.476\times 10^{-5}$	454	$1.113\times10^{-9}$
GENROSE	10	$1 \times 10^{-1}$	84	9.259	221	8.902	805	8.892
GENROSE	10	$1 \times 10^{-3}$	585	$4.321\times10^{-3}$	446	8.816	1564	$1.860 \times 10^{-4}$
GENROSE	10	$1 \times 10^{-5}$	654	$3.114 imes10^{-5}$	1016	$7.485\times10^{-3}$	1935	$1.029\times 10^{-6}$
GENROSE	10	$1 \times 10^{-7}$	726	$7.303\times10^{-7}$	964	$8.881\times 10^{-5}$	1904	$5.181\times10^{-10}$
GENROSE	100	$1 \times 10^{-1}$	1868	$1.189  imes 10^2$	3231	$1.335  imes 10^2$	7344	$1.119  imes 10^2$
GENROSE	100	$1 \times 10^{-3}$	15359	$1.081  imes 10^2$	2247	$1.108  imes 10^2$	50024*	$6.559  imes 10^1$
GENROSE	100	$1 \times 10^{-5}$	18075	$6.885 \times 10^{-3}$	29786	$5.081 \times 10^{-3}$	50009*	$2.770 \times 10^{-1}$
GENROSE	100	$1 \times 10^{-7}$	18030	$1.719  imes 10^{-6}$	29847	$4.748 \times 10^{-5}$	50005*	$2.194\times10^{-2}$
GULF	3	$1 \times 10^{-1}$	26	7.047	41	6.637	389	6.622
GULF	3	$1 \times 10^{-3}$	38	7.032	38	6.622	229	$4.534 \times 10^{-3}$
GULF	3	$1 \times 10^{-5}$	115	$6.882 \times 10^{-2}$	267	$3.649 \times 10^{-3}$	370	$4.462 \times 10^{-3}$
GULF	3	$1 \times 10^{-7}$	333	$6.930 \times 10^{-3}$	217	$4.381 \times 10^{-3}$	688	$2.460 \times 10^{-7}$
HAIRY	2	$1 \times 10^{-1}$	49	$3.471 \times 10^{2}$	71	$2.721 \times 10^{2}$	387	$1.950 \times 10^{-4}$
HAIRY	2	$1 \times 10^{-3}$	351	$2.426 \times 10^{-4}$	186	$8.522 \times 10^{-4}$	413	$2.326 \times 10^{-8}$
HAIRY	2	$1 \times 10^{-5}$	226	$1.256 \times 10^{-8}$	175	$9.401 \times 10^{-7}$	516	$5.074 \times 10^{-11}$
HAIRY	2	$1 \times 10^{-7}$	252	$2.339 \times 10^{-8}$	87	$1.508 \times 10^{-6}$	210	$4.974 \times 10^{-14}$
HELIX	3	$1 \times 10^{-1}$	26	$9.616 \times 10^{-2}$	188	7.615	374	$2.785 \times 10^{-1}$
HELIX	3	$1 \times 10^{-3}$	38	$1.865 \times 10^{-2}$	314	4.147	23	$8.522 \times 10^{2}$
HELIX	3	$1 \times 10^{-5}$	72	$1.822 \times 10^{-6}$	453	$3.021 \times 10^{-6}$	23	$8.521 \times 10^{2}$
HELIX	3	$1 \times 10^{-7}$	84	$4.567 \times 10^{-7}$	167	$1.171 \times 10^{-4}$	23	$8.521 \times 10^2$
JENSMP <sup>‡</sup>	2	$1 \times 10^{-1}$	24	$-1.786 \times 10^{3}$	18	0.000	21	0.000
JENSMP <sup>‡</sup>	2	$1 \times 10^{-3}$	78	$-1.896 \times 10^{3}$	18	0.000	21	0.000
JENSMP <sup>‡</sup>	2	$1 \times 10^{-5}$	87	$-1.896 \times 10^{3}$	18	0.000	21	0.000
JENSMP <sup>‡</sup>	2	$1 \times 10^{-7}$	67	$-1.896 \times 10^{3}$	18	0.000	21	0.000
KOWOSB	4	$1 \times 10^{-1}$	1**	$5.006 \times 10^{-3}$	72	$1.383 \times 10^{-3}$	9	$5.006 \times 10^{-3}$
KOWOSB	4	$1 \times 10^{-3}$	20	$2.240 \times 10^{-3}$	261	$3.305 \times 10^{-3}$	104	$2.435 \times 10^{-4}$
KOWOSB	4	$1 \times 10^{-5}$	44	$1.249 \times 10^{-4}$	112	$1.767 \times 10^{-4}$	220	$6.118 \times 10^{-5}$
KOWOSB	4	$1 \times 10^{-7}$	111	$2.694 \times 10^{-5}$	1309	$6.915 \times 10^{-7}$	314	$5.638 \times 10^{-8}$

Table A.13. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			N	EWUOA	FD	L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
MEXHAT	2	$1 \times 10^{-1}$	30	$1.166\times 10^{-2}$	98	$4.477 \times 10^{-1}$	146	$4.231 \times 10^{-1}$
MEXHAT	2	$1 \times 10^{-3}$	43	$1.102 \times 10^{-2}$	80	$8.450 \times 10^{-1}$	709	$4.227 \times 10^{-1}$
MEXHAT	2	$1 \times 10^{-5}$	45	$1.110\times 10^{-2}$	203	$4.173\times10^{-1}$	232	$1.568 \times 10^{-3}$
MEXHAT	2	$1 \times 10^{-7}$	44	$1.036\times 10^{-2}$	210	$4.226 \times 10^{-1}$	660	$1.383  imes 10^{-4}$
MOREBV	10	$1 \times 10^{-1}$	$1^{**}$	$1.599\times 10^{-2}$	19	$1.599\times 10^{-2}$	21	$1.599\times 10^{-2}$
MOREBV	10	$1 \times 10^{-3}$	81	$1.262\times 10^{-2}$	131	$1.451\times 10^{-2}$	937	$1.506\times 10^{-3}$
MOREBV	10	$1 \times 10^{-5}$	184	$2.345\times10^{-3}$	433	$1.847 \times 10^{-3}$	796	$1.441 \times 10^{-6}$
MOREBV	10	$1 \times 10^{-7}$	386	$1.283\times 10^{-7}$	988	$7.402\times10^{-5}$	1003	$5.675 \times 10^{-9}$
MOREBV	50	$1 \times 10^{-1}$	$1^{**}$	$1.321 \times 10^{-3}$	93	$1.321\times 10^{-3}$	101	$1.321\times 10^{-3}$
MOREBV	50	$1 \times 10^{-3}$	105	$1.122\times 10^{-3}$	101	$1.321\times 10^{-3}$	828	$4.204\times10^{-4}$
MOREBV	50	$1 \times 10^{-5}$	746	$3.981\times 10^{-4}$	997	$7.841\times10^{-4}$	2064	$2.290\times 10^{-5}$
MOREBV	50	$1 \times 10^{-7}$	1322	$3.021  imes 10^{-5}$	1283	$4.621\times 10^{-5}$	12825	$7.628\times10^{-6}$
MOREBV	100	$1 \times 10^{-1}$	$1^{**}$	$3.633 \times 10^{-4}$	190	$3.633\times 10^{-4}$	201	$3.633\times10^{-4}$
MOREBV	100	$1 \times 10^{-3}$	$1^{**}$	$3.633 \times 10^{-4}$	201	$3.633\times 10^{-4}$	2669	$1.495\times10^{-4}$
MOREBV	100	$1 \times 10^{-5}$	320	$2.538\times10^{-4}$	309	$3.571  imes 10^{-4}$	3457	$6.772\times10^{-6}$
MOREBV	100	$1 \times 10^{-7}$	3087	$1.417\times 10^{-5}$	2809	$1.540 imes10^{-5}$	14045	$1.546\times 10^{-6}$
NCB20B	21	$1 \times 10^{-1}$	1**	$3.000 \times 10^{-3}$	22	$3.000  imes 10^{-3}$	232	$3.177  imes 10^{-4}$
NCB20B	21	$1 \times 10^{-3}$	113	$1.441 \times 10^{-5}$	74	$2.964\times10^{-4}$	421	$6.287\times10^{-6}$
NCB20B	21	$1 \times 10^{-5}$	260	$1.402\times10^{-4}$	1332	$5.741 \times 10^{-6}$	686	$6.098 \times 10^{-9}$
NCB20B	21	$1 \times 10^{-7}$	708	$1.224 \times 10^{-5}$	260	$4.816 \times 10^{-8}$	650	$-5.574 \times 10^{-11}$
NCB20B	22	$1 \times 10^{-1}$	121	$5.384  imes 10^{-3}$	23	$6.000 \times 10^{-3}$	388	$1.036 \times 10^{-3}$
NCB20B	22	$1 \times 10^{-3}$	121	$1.817 \times 10^{-3}$	174	$8.594 imes10^{-4}$	935	$8.921 \times 10^{-4}$
NCB20B	22	$1 \times 10^{-5}$	265	$2.971 \times 10^{-4}$	298	$1.075 \times 10^{-5}$	578	$3.829 \times 10^{-6}$
NCB20B	22	$1 \times 10^{-7}$	715	$2.307 \times 10^{-5}$	345	$1.163 \times 10^{-7}$	908	$2.372 \times 10^{-9}$
NCB20B	50	$1 \times 10^{-1}$	212	$9.199 \times 10^{-2}$	121	$1.099 \times 10^{-1}$	938	$1.883 \times 10^{-2}$
NCB20B	50	$1 \times 10^{-3}$	734	$7.836 \times 10^{-3}$	1513	$1.296 \times 10^{-3}$	1866	$3.565 \times 10^{-4}$
NCB20B	50	$1 \times 10^{-5}$	1757	$4.549 \times 10^{-4}$	838	$3.965 \times 10^{-4}$	3623	$9.630 \times 10^{-5}$
NCB20B	50	$1 \times 10^{-7}$	4316	$6.696 \times 10^{-5}$	2253	$8.619 \times 10^{-5}$	23370	$2.289 \times 10^{-6}$
NCB20B	100	$1 \times 10^{-1}$	481	$3.871 \times 10^{-1}$	1565	$3.458 \times 10^{-1}$	2938	$3.874 \times 10^{-2}$
NCB20B	100	$1 \times 10^{-3}$	1748	$1.447 \times 10^{-2}$	2557	$2.895 \times 10^{-2}$	5703	$8.588 \times 10^{-4}$
NCB20B	100	$1 \times 10^{-3}$	5976	$7.143 \times 10^{-4}$	5154	$3.922 \times 10^{-4}$	10509	$8.218 \times 10^{-5}$
NCB20B	100	$1 \times 10^{-7}$	15651	$4.301 \times 10^{-3}$	6944	$6.476 \times 10^{-3}$	32343	$1.546 \times 10^{-3}$
NCB20B	180	$1 \times 10^{-1}$	1664	$4.700 \times 10^{-1}$	1648	1.184	5833	$1.772 \times 10^{-1}$
NCB20B	180	$1 \times 10^{-5}$	6604	$4.261 \times 10^{-2}$	2007	$5.420 \times 10^{-2}$	12753	$1.807 \times 10^{-5}$
NCB20B	180	$1 \times 10^{-7}$	12066	$1.067 \times 10^{-5}$	10462	$3.299 \times 10^{-4}$	22900	$2.856 \times 10^{-5}$
NCB20B	180	$1 \times 10^{-1}$	26507	$2.789 \times 10^{-6}$	11118	$2.804 \times 10^{-6}$	28345	$2.440 \times 10^{-6}$
NONDIA	10	$1 \times 10^{-3}$	84 144	4.198	407	0.347	124 570	0.487
NONDIA	10	$1 \times 10^{-5}$	144	$1.406 \times 10^{-5}$	19	0.481	579 240	$1.101 \times 10^{-1}$
NONDIA	10	$1 \times 10^{-7}$	108	$3.507 \times 10^{-7}$	93	0.480	349	$5.342 \times 10^{-10}$
NONDIA	10	$1 \times 10^{-1}$	190	$2.093 \times 10^{-1}$	180	$1.373 \times 10^{-1}$	00E	4.429 × 10 -*
NONDIA	20	$1 \times 10^{-1}$ $1 \times 10^{-3}$	104 956	2.000 $2.000 \times 10^{-3}$	140	4.149 4 495	000 744	4.400 5 560 $\times$ 10 <sup>-4</sup>
NONDIA	20	$1 \times 10^{-5}$	300 409	$2.292 \times 10$ 2.162 $\times 10^{-5}$	200	4.435 1 700 × 10 <sup>-3</sup>	006	$3.300 \times 10$ $4.728 \times 10^{-7}$
NONDIA	20 20	$1 \times 10^{-7}$	403 468	$3.103 \times 10^{-6}$ 1 300 $\times 10^{-6}$	282	$1.709 \times 10^{-7}$ $1.588 \times 10^{-7}$	1033	$4.730 \times 10^{-9}$ $3.743 \times 10^{-9}$
NONDIA	20	$1 \times 10^{-1}$	100	1 /07	160	2 986	1900 517	2.740 × 10 2.971
NONDIA	30	$1 \times 10^{-3}$	199 566	1.497 6.818 $\vee 10^{-3}$	780	$1.606 \times 10^{-2}$	1210	3.271 1 828 $\vee$ 10 <sup>-3</sup>
NONDIA	30	$1 \times 10^{-5}$	580	$1.475 \times 10^{-4}$	680	$8.898 \times 10^{-4}$	1082	$5.285 \times 10^{-6}$
NONDIA	30	$1 \times 10^{-7}$	634	$3.771 \times 10^{-7}$	1310	$6.491 \times 10^{-7}$	1551	$4.326 \times 10^{-9}$
NONDIA	30	1 \ 10	054	0.111 × 10	1019	0.431 × 10	1001	$4.020 \times 10$

Table A.14. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			NEWUOA		FD	) L-BFGS	CD L-BFGS	
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
NONDIA	50	$1 \times 10^{-1}$	317	$8.665\times 10^{-1}$	269	1.505	423	1.520
NONDIA	50	$1 \times 10^{-3}$	635	$8.814\times10^{-3}$	375	1.498	1659	$2.467\times10^{-3}$
NONDIA	50	$1 \times 10^{-5}$	851	$7.108 \times 10^{-5}$	3599	$4.640\times10^{-4}$	1659	$6.908\times10^{-6}$
NONDIA	50	$1 \times 10^{-7}$	1132	$1.517\times 10^{-6}$	1731	$1.971\times 10^{-5}$	1556	$1.127\times 10^{-8}$
NONDIA	90	$1 \times 10^{-1}$	573	$3.893 \times 10^{-1}$	1087	$6.154\times10^{-1}$	744	$5.843\times10^{-1}$
NONDIA	90	$1 \times 10^{-3}$	1374	$3.301  imes 10^{-1}$	563	$5.766\times10^{-1}$	2052	$9.232\times10^{-3}$
NONDIA	90	$1 \times 10^{-5}$	1720	$9.430 \times 10^{-5}$	564	$5.781\times10^{-1}$	3319	$2.610\times10^{-5}$
NONDIA	90	$1 \times 10^{-7}$	1788	$2.403\times10^{-6}$	1306	$4.184\times10^{-6}$	4057	$5.888\times 10^{-8}$
NONDIA	100	$1 \times 10^{-1}$	419	$3.342 \times 10^{-1}$	418	$4.850 \times 10^{-1}$	824	$4.930 \times 10^{-1}$
NONDIA	100	$1 \times 10^{-3}$	1610	$2.086\times10^{-1}$	927	$4.803\times10^{-1}$	4544	$1.815\times10^{-2}$
NONDIA	100	$1 \times 10^{-5}$	1907	$1.474\times10^{-4}$	934	$4.658\times10^{-1}$	4975	$4.317\times10^{-5}$
NONDIA	100	$1 \times 10^{-7}$	2235	$2.091\times 10^{-6}$	2668	$5.756 imes10^{-5}$	4072	$1.010  imes 10^{-7}$
NONDQUAR	100	$1 \times 10^{-1}$	828	$5.744 \times 10^{-1}$	1433	$7.564 \times 10^{-1}$	3532	$1.006 \times 10^{-1}$
NONDQUAR	100	$1 \times 10^{-3}$	1842	$1.889\times 10^{-2}$	5568	$2.000\times 10^{-2}$	7272	$5.244  imes 10^{-3}$
NONDQUAR	100	$1 \times 10^{-5}$	4661	$1.475 \times 10^{-3}$	8339	$1.946\times10^{-3}$	20194	$3.394 \times 10^{-4}$
NONDQUAR	100	$1 \times 10^{-7}$	26783	$1.700\times 10^{-4}$	21822	$1.549\times10^{-4}$	50037*	$3.497\times 10^{-5}$
OSBORNEA <sup>‡</sup>	5	$1 \times 10^{-1}$	52	$1.127\times 10^{-2}$	8	$8.790 \times 10^{-1}$	170	$1.485 \times 10^{-1}$
OSBORNEA <sup>‡</sup>	5	$1 \times 10^{-3}$	64	$5.240\times10^{-2}$	9	$8.790\times10^{-1}$	496	$7.589 \times 10^{-4}$
OSBORNEA <sup>‡</sup>	5	$1 \times 10^{-5}$	135	$3.040 \times 10^{-3}$	260	$1.902 \times 10^{-3}$	1638	$2.432\times 10^{-5}$
OSBORNEA <sup>‡</sup>	5	$1 \times 10^{-7}$	504	$4.404\times10^{-5}$	677	$2.252\times 10^{-5}$	1033	$2.247\times 10^{-5}$
OSBORNEB	11	$1 \times 10^{-1}$	89	$5.190 imes10^{-1}$	308	$6.677\times10^{-1}$	1031	$2.972\times 10^{-1}$
OSBORNEB	11	$1 \times 10^{-3}$	172	$3.075 \times 10^{-1}$	829	$3.219  imes 10^{-1}$	1238	$7.480\times10^{-2}$
OSBORNEB	11	$1 \times 10^{-5}$	1052	$7.968  imes 10^{-3}$	1249	$2.882\times10^{-3}$	2334	$7.854\times10^{-5}$
OSBORNEB	11	$1 \times 10^{-7}$	1708	$8.191\times 10^{-6}$	1555	$6.681\times10^{-6}$	2395	$9.143 \times 10^{-9}$
PENALTY1	4	$1 \times 10^{-1}$	49	$1.917\times 10^{-2}$	40	$8.686\times 10^{-3}$	100	$4.980\times 10^{-3}$
PENALTY1	4	$1 \times 10^{-3}$	87	$2.835\times 10^{-5}$	165	$3.661\times 10^{-5}$	95	$4.498\times 10^{-5}$
PENALTY1	4	$1 \times 10^{-5}$	117	$3.573\times10^{-5}$	55	$1.741\times 10^{-4}$	106	$3.829\times 10^{-5}$
PENALTY1	4	$1 \times 10^{-7}$	113	$2.254\times 10^{-5}$	97	$3.820\times 10^{-5}$	140	$3.826\times 10^{-5}$
PENALTY1	10	$1 \times 10^{-1}$	358	$4.733  imes 10^{-1}$	1167	$5.627\times 10^{-2}$	284	$1.327  imes 10^{-3}$
PENALTY1	10	$1 \times 10^{-3}$	178	$3.567  imes 10^{-5}$	347	$2.207\times 10^{-4}$	353	$9.096  imes 10^{-5}$
PENALTY1	10	$1 \times 10^{-5}$	194	$2.463\times10^{-5}$	335	$6.078  imes 10^{-5}$	595	$6.005  imes 10^{-5}$
PENALTY1	10	$1 \times 10^{-7}$	245	$2.985\times10^{-5}$	352	$5.882  imes 10^{-5}$	477	$5.965  imes 10^{-5}$
PENALTY1	50	$1 \times 10^{-1}$	2527	$5.606 imes10^{-1}$	2274	1.423	2076	$8.180 imes10^{-3}$
PENALTY1	50	$1 \times 10^{-3}$	2455	$6.140\times10^{-5}$	2697	$1.201\times 10^{-4}$	2281	$2.999\times 10^{-4}$
PENALTY1	50	$1 \times 10^{-5}$	2298	$5.576 imes10^{-5}$	1786	$9.926\times 10^{-5}$	2384	$1.322\times 10^{-4}$
PENALTY1	50	$1 \times 10^{-7}$	3919	$8.208\times10^{-6}$	1523	$8.616\times 10^{-5}$	16116	$7.338\times10^{-6}$
PENALTY1	100	$1 \times 10^{-1}$	5649	$8.758\times10^{-1}$	6442	6.236	4687	$8.428\times 10^{-3}$
PENALTY1	100	$1 \times 10^{-3}$	6701	$1.019\times 10^{-4}$	3923	$1.489\times10^{-4}$	5093	$1.859\times 10^{-4}$
PENALTY1	100	$1 \times 10^{-5}$	5940	$9.287  imes 10^{-5}$	4516	$1.904\times10^{-4}$	5907	$1.878\times 10^{-4}$
PENALTY1	100	$1 \times 10^{-7}$	10954	$1.256\times 10^{-5}$	6607	$1.677\times 10^{-4}$	28399	$5.191  imes 10^{-6}$
$\rm PFIT1LS^{\ddagger}$	3	$1 \times 10^{-1}$	35	4.034	67	9.281	389	5.246
PFIT1LS <sup>‡</sup>	3	$1 \times 10^{-3}$	111	$9.255\times10^{-1}$	1415	7.002	447	$2.288\times 10^{-4}$
PFIT1LS <sup>‡</sup>	3	$1 \times 10^{-5}$	325	$1.154\times 10^{-2}$	342	$2.235\times 10^{-6}$	502	$2.581\times10^{-6}$
PFIT1LS <sup>‡</sup>	3	$1 \times 10^{-7}$	482	$2.955\times 10^{-4}$	258	$4.493  imes 10^{-5}$	453	$2.489\times 10^{-6}$
$\rm PFIT2LS^{\ddagger}$	3	$1 \times 10^{-1}$	113	$5.174 \times 10^1$	81	$9.874 \times 10^1$	620	$2.712\times10^{-1}$
$PFIT2LS^{\ddagger}$	3	$1 \times 10^{-3}$	410	2.136	123	$1.245\times 10^2$	537	$6.269\times10^{-3}$
$\mathbf{PFIT2LS}^{\ddagger}$	3	$1 \times 10^{-5}$	618	$5.570\times10^{-2}$	430	$4.757\times10^{-3}$	537	$1.772\times 10^{-3}$
$\mathbf{PFIT2LS}^{\ddagger}$	3	$1 \times 10^{-7}$	706	$1.251\times 10^{-2}$	340	$2.704\times10^{-3}$	870	$1.648\times 10^{-3}$

Table A.15. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			N	EWUOA	FD	L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
PFIT3LS <sup>‡</sup>	3	$1 \times 10^{-1}$	162	$1.868 \times 10^2$	390	$7.944 \times 10^{1}$	817	$3.682 \times 10^{-1}$
$PFIT3LS^{\ddagger}$	3	$1 \times 10^{-3}$	546	7.561	640	$9.159  imes 10^{-1}$	612	$3.276 \times 10^{-2}$
PFIT3LS <sup>‡</sup>	3	$1 \times 10^{-5}$	1012	$9.384 \times 10^{-2}$	615	$4.587  imes 10^{-2}$	702	$3.154 \times 10^{-2}$
PFIT3LS <sup>‡</sup>	3	$1 \times 10^{-7}$	1012	$8.229 \times 10^{-2}$	444	$4.374\times10^{-2}$	$1504^{*}$	$2.782\times10^{-2}$
PFIT4LS <sup>‡</sup>	3	$1 \times 10^{-1}$	234	$2.464 \times 10^{2}$	78	$2.144 \times 10^{3}$	723	$1.842 \times 10^{-1}$
$PFIT4LS^{\ddagger}$	3	$1 \times 10^{-3}$	612	$2.736  imes 10^1$	758	6.570	715	$1.234 \times 10^{-1}$
PFIT4LS <sup>‡</sup>	3	$1 \times 10^{-5}$	1181	$8.871 \times 10^{-1}$	471	$1.677 \times 10^{-1}$	$1506^{*}$	$1.433 \times 10^{-1}$
$PFIT4LS^{\ddagger}$	3	$1 \times 10^{-7}$	1374	$2.619\times 10^{-1}$	597	$1.249\times 10^{-1}$	$1502^{*}$	$1.086 \times 10^{-1}$
QUARTC	25	$1 \times 10^{-1}$	599	$4.172\times10^{-1}$	1716	2.004	1203	$1.376 \times 10^{-2}$
QUARTC	25	$1 \times 10^{-3}$	759	$7.392  imes 10^{-3}$	767	$3.811\times 10^{-2}$	1496	$8.968  imes 10^{-5}$
QUARTC	25	$1 \times 10^{-5}$	828	$6.201\times10^{-6}$	1245	$8.620\times 10^{-5}$	1750	$3.508  imes 10^{-7}$
QUARTC	25	$1 \times 10^{-7}$	1040	$4.136\times 10^{-7}$	1584	$1.148 \times 10^{-6}$	2807	$1.313 \times 10^{-8}$
QUARTC	100	$1 \times 10^{-1}$	4920	$3.885  imes 10^{-1}$	4766	6.422	7146	$4.248\times 10^{-2}$
QUARTC	100	$1 \times 10^{-3}$	5612	$9.569  imes 10^{-3}$	4517	$2.474\times10^{-3}$	5960	$1.807 \times 10^{-4}$
QUARTC	100	$1 \times 10^{-5}$	6494	$5.767  imes 10^{-5}$	6534	$7.941\times10^{-4}$	7146	$2.060 \times 10^{-6}$
QUARTC	100	$1 \times 10^{-7}$	6913	$8.633 \times 10^{-7}$	1119	$1.906 \times 10^7$	7344	$2.041 \times 10^{-8}$
SINEVAL	2	$1 \times 10^{-1}$	12	5.231	4	5.552	108	4.632
SINEVAL	2	$1 \times 10^{-3}$	152	$9.339\times10^{-1}$	194	3.587	700	$2.082\times 10^{-4}$
SINEVAL	2	$1 \times 10^{-5}$	226	$1.682 \times 10^{-5}$	450	$1.856 \times 10^{-1}$	488	$1.360 \times 10^{-7}$
SINEVAL	2	$1 \times 10^{-7}$	295	$1.841\times 10^{-8}$	425	$1.955\times10^{-5}$	563	$9.899 \times 10^{-13}$
SINQUAD	5	$1 \times 10^{-1}$	36	1.503	333	5.219	107	$1.121\times 10^{-2}$
SINQUAD	5	$1 \times 10^{-3}$	98	$8.682\times 10^{-3}$	79	$6.186\times10^{-3}$	121	$3.326\times 10^{-5}$
SINQUAD	5	$1 \times 10^{-5}$	140	$1.602\times 10^{-5}$	77	$3.260  imes 10^{-5}$	121	$9.620 \times 10^{-8}$
SINQUAD	5	$1 \times 10^{-7}$	170	$1.177\times 10^{-6}$	135	$5.826 \times 10^{-9}$	205	$-3.938 \times 10^{-10}$
SINQUAD	50	$1 \times 10^{-1}$	1724	1.740	742	$2.222 \times 10^1$	3929	$1.376  imes 10^{-1}$
SINQUAD	50	$1 \times 10^{-3}$	3724	$3.119\times 10^{-2}$	1101	$9.615\times10^{-3}$	3616	$2.544\times10^{-4}$
SINQUAD	50	$1 \times 10^{-5}$	4234	$3.236  imes 10^{-4}$	849	$2.589\times 10^{-4}$	2857	$6.467  imes 10^{-7}$
SINQUAD	50	$1  imes 10^{-7}$	6464	$1.670\times 10^{-6}$	1261	$1.644\times 10^{-6}$	2083	$5.443 \times 10^{-9}$
SINQUAD	100	$1 \times 10^{-1}$	4978	4.096	1247	$1.048 \times 10^1$	3062	$1.026\times10^{-1}$
SINQUAD	100	$1 \times 10^{-3}$	7193	$2.876\times10^{-2}$	2470	$6.295\times10^{-2}$	3265	$2.132\times 10^{-4}$
SINQUAD	100	$1 \times 10^{-5}$	10426	$3.206\times 10^{-4}$	1948	$9.257\times10^{-4}$	3468	$5.694 imes10^{-7}$
SINQUAD	100	$1 \times 10^{-7}$	15008	$2.263\times10^{-6}$	2152	$3.046  imes 10^{-6}$	4079	$3.929  imes 10^{-9}$
SISSER	2	$1 \times 10^{-1}$	4**	$3.000\times10^{-4}$	20	$2.326\times 10^{-2}$	65	$2.424\times10^{-4}$
SISSER	2	$1 \times 10^{-3}$	15	$1.764\times10^{-4}$	139	$2.826 imes10^{-5}$	204	$1.415\times 10^{-4}$
SISSER	2	$1 \times 10^{-5}$	20	$2.270 \times 10^{-5}$	44	$3.731 \times 10^{-5}$	85	$2.188\times10^{-9}$
SISSER	2	$1 \times 10^{-7}$	27	$1.467\times10^{-8}$	71	$4.270 \times 10^{-7}$	116	$6.695 \times 10^{-12}$
SPARSQUR	10	$1 \times 10^{-1}$	42	$2.779  imes 10^{-1}$	104	$4.304 imes10^{-2}$	403	$5.561 \times 10^{-3}$
SPARSQUR	10	$1 \times 10^{-3}$	85	$2.675\times10^{-3}$	212	$6.281  imes 10^{-3}$	384	$6.274 \times 10^{-5}$
SPARSQUR	10	$1 \times 10^{-5}$	184	$4.915 \times 10^{-5}$	740	$2.770 \times 10^{-5}$	384	$3.317 \times 10^{-7}$
SPARSQUR	10	$1 \times 10^{-7}$	241	$1.377 \times 10^{-6}$	398	$1.386 \times 10^{-7}$	579	$1.307 \times 10^{-8}$
SPARSQUR	50	$1 \times 10^{-1}$	1070	$7.980\times10^{-1}$	1303	$6.397 imes10^{-1}$	1316	$1.918\times10^{-2}$
SPARSQUR	50	$1 \times 10^{-3}$	1034	$4.918\times10^{-3}$	1078	$1.952\times 10^{-2}$	3577	$9.425\times10^{-5}$
SPARSQUR	50	$1 \times 10^{-5}$	1472	$2.315\times10^{-5}$	2755	$4.135 \times 10^{-5}$	2651	$5.338 \times 10^{-7}$
SPARSQUR	50	$1 \times 10^{-7}$	1847	$3.804\times10^{-7}$	1706	$5.903  imes 10^{-6}$	2651	$6.858\times10^{-9}$
SPARSQUR	100	$1 \times 10^{-1}$	2912	$5.988  imes 10^{-1}$	3387	1.953	5960	$1.100\times 10^{-2}$
SPARSQUR	100	$1 \times 10^{-3}$	2484	$4.672 \times 10^{-3}$	2435	$5.029\times10^{-2}$	3677	$5.412 \times 10^{-5}$
SPARSQUR	100	$1 \times 10^{-5}$	2664	$7.151\times10^{-5}$	5794	$1.737\times10^{-4}$	4292	$6.552\times10^{-6}$
SPARSQUR	100	$1 \times 10^{-7}$	3757	$3.002\times10^{-7}$	4591	$2.610\times10^{-6}$	5960	$1.082\times10^{-8}$

Table A.16. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			NI	EWUOA	FD	L-BFGS	CI	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
TOINTGSS	10	$1 \times 10^{-1}$	65	$4.850 \times 10^{-1}$	759	2.072	101	$1.831 \times 10^{-2}$
TOINTGSS	10	$1 \times 10^{-3}$	164	$1.448\times 10^{-2}$	301	$3.304\times10^{-2}$	165	$8.167 \times 10^{-6}$
TOINTGSS	10	$1 \times 10^{-5}$	211	$2.845 \times 10^{-5}$	124	$3.624 \times 10^{-4}$	200	$2.744 \times 10^{-8}$
TOINTGSS	10	$1 \times 10^{-7}$	277	$2.113\times 10^{-7}$	124	$3.539\times 10^{-6}$	200	$5.802\times10^{-11}$
TOINTGSS	50	$1 \times 10^{-1}$	388	1.033	563	$1.183 \times 10^1$	724	$6.155 \times 10^{-2}$
TOINTGSS	50	$1 \times 10^{-3}$	878	$1.276\times 10^{-2}$	415	$4.651\times 10^{-2}$	2132	$1.554\times10^{-5}$
TOINTGSS	50	$1 \times 10^{-5}$	923	$1.748\times10^{-4}$	693	$1.140\times 10^{-3}$	724	$5.405  imes 10^{-8}$
TOINTGSS	50	$1 \times 10^{-7}$	1574	$1.030\times 10^{-6}$	693	$1.136\times 10^{-5}$	724	$-2.877\times10^{-9}$
TOINTGSS	100	$1 \times 10^{-1}$	1638	1.693	590	7.090	2660	$6.483\times10^{-2}$
TOINTGSS	100	$1 \times 10^{-3}$	1521	$2.165\times 10^{-2}$	509	$7.545\times10^{-2}$	1646	$3.228 \times 10^{-5}$
TOINTGSS	100	$1 \times 10^{-5}$	2692	$2.385\times 10^{-4}$	508	$1.898\times 10^{-3}$	2526	$2.620\times 10^{-8}$
TOINTGSS	100	$1 \times 10^{-7}$	2312	$3.442\times 10^{-6}$	508	$1.893\times 10^{-5}$	2526	$4.046\times10^{-9}$
TQUARTIC	5	$1 \times 10^{-1}$	22	$4.614\times 10^{-1}$	24	$6.142\times 10^{-1}$	40	$4.540\times10^{-1}$
TQUARTIC	5	$1 \times 10^{-3}$	63	$5.640 \times 10^{-3}$	242	$9.902\times10^{-2}$	234	$5.290 \times 10^{-4}$
TQUARTIC	5	$1 \times 10^{-5}$	107	$1.155\times 10^{-5}$	202	$8.160\times 10^{-4}$	818	$1.004 \times 10^{-6}$
TQUARTIC	5	$1 \times 10^{-7}$	137	$7.045\times10^{-8}$	95	$1.098 \times 10^{-6}$	163	$9.192 \times 10^{-9}$
TQUARTIC	10	$1 \times 10^{-1}$	44	$6.151\times 10^{-1}$	274	$7.097\times10^{-1}$	92	$5.464 \times 10^{-1}$
TQUARTIC	10	$1 \times 10^{-3}$	174	$4.406\times10^{-2}$	543	$2.485\times10^{-1}$	344	$5.223 \times 10^{-3}$
TQUARTIC	10	$1 \times 10^{-5}$	277	$2.548\times 10^{-5}$	466	$4.217\times 10^{-2}$	722	$7.389\times10^{-6}$
TQUARTIC	10	$1 \times 10^{-7}$	388	$8.693\times10^{-7}$	805	$6.160\times10^{-5}$	535	$4.256 \times 10^{-8}$
TQUARTIC	50	$1 \times 10^{-1}$	191	$6.732  imes 10^{-1}$	170	$7.480\times10^{-1}$	310	$7.002 \times 10^{-1}$
TQUARTIC	50	$1 \times 10^{-3}$	1487	$3.648 \times 10^{-1}$	1155	$5.611\times10^{-1}$	1534	$6.796 \times 10^{-1}$
TQUARTIC	50	$1 \times 10^{-5}$	3112	$6.435\times10^{-3}$	1441	$2.121\times 10^{-2}$	2117	$1.390  imes 10^{-4}$
TQUARTIC	50	$1 \times 10^{-7}$	5704	$8.027\times 10^{-5}$	3736	$3.333\times 10^{-4}$	1785	$3.649\times 10^{-7}$
TQUARTIC	100	$1 \times 10^{-1}$	226	$7.236\times10^{-1}$	236	$7.828\times 10^{-1}$	610	$7.423  imes 10^{-1}$
TQUARTIC	100	$1 \times 10^{-3}$	1085	$7.089\times10^{-1}$	790	$7.232\times 10^{-1}$	3629	$7.219\times10^{-1}$
TQUARTIC	100	$1 \times 10^{-5}$	26574	$2.458\times10^{-2}$	2675	$1.765\times10^{-1}$	4298	$5.460 imes10^{-4}$
TQUARTIC	100	$1 \times 10^{-7}$	31937	$2.808\times 10^{-4}$	4860	$4.761\times10^{-4}$	5838	$2.894  imes 10^{-7}$
TRIDIA	10	$1 \times 10^{-1}$	128	1.555	124	1.392	594	$3.847  imes 10^{-4}$
TRIDIA	10	$1 \times 10^{-3}$	174	$1.649\times 10^{-3}$	305	$2.065\times 10^{-2}$	495	$1.429\times 10^{-6}$
TRIDIA	10	$1 \times 10^{-5}$	176	$2.353\times10^{-5}$	242	$3.449\times 10^{-4}$	698	$4.786 \times 10^{-9}$
TRIDIA	10	$1 \times 10^{-7}$	226	$1.205\times 10^{-7}$	279	$3.036\times 10^{-6}$	848	$2.043 \times 10^{-12}$
TRIDIA	20	$1 \times 10^{-1}$	301	$7.391\times10^{-1}$	812	1.072	2309	$4.788 \times 10^{-4}$
TRIDIA	20	$1 \times 10^{-3}$	312	$9.505  imes 10^{-3}$	443	$1.352\times10^{-1}$	1432	$1.038 \times 10^{-6}$
TRIDIA	20	$1 \times 10^{-5}$	396	$4.890 \times 10^{-5}$	821	$4.616\times10^{-4}$	2029	$3.474 \times 10^{-9}$
TRIDIA	20	$1 \times 10^{-7}$	500	$5.289  imes 10^{-7}$	864	$3.588  imes 10^{-6}$	1992	$7.669 \times 10^{-12}$
TRIDIA	30	$1 \times 10^{-1}$	454	$8.863 \times 10^{-1}$	3249	7.566	2694	$2.364 \times 10^{-3}$
TRIDIA	30	$1 \times 10^{-3}$	599	$7.428\times10^{-3}$	2486	$6.494\times10^{-2}$	2403	$4.586 \times 10^{-5}$
TRIDIA	30	$1 \times 10^{-5}$	753	$1.083\times10^{-4}$	1254	$1.150 \times 10^{-3}$	3111	$7.820 \times 10^{-7}$
TRIDIA	30	$1 \times 10^{-7}$	849	$7.455 \times 10^{-7}$	2292	$5.997\times10^{-6}$	4598	$6.846 \times 10^{-11}$
TRIDIA	50	$1 \times 10^{-1}$	670	1.625	1604	8.841	3723	$1.190 \times 10^{-1}$
TRIDIA	50	$1 \times 10^{-3}$	1124	$1.616\times10^{-2}$	1824	$2.452\times10^{-1}$	6195	$7.991 \times 10^{-6}$
TRIDIA	50	$1 \times 10^{-5}$	1056	$4.101\times10^{-4}$	3602	$1.237\times 10^{-3}$	5985	$2.190\times10^{-6}$
TRIDIA	50	$1 \times 10^{-7}$	1784	$1.022\times 10^{-6}$	3450	$2.572\times10^{-5}$	7640	$1.634\times 10^{-8}$
TRIDIA	100	$1 \times 10^{-1}$	2183	5.214	4778	$2.202 \times 10^1$	9774	$1.524\times10^{-1}$
TRIDIA	100	$1 \times 10^{-3}$	4301	$3.587\times10^{-2}$	3474	$6.746\times10^{-1}$	15035	$9.421\times10^{-5}$
TRIDIA	100	$1 \times 10^{-5}$	4682	$4.705\times10^{-4}$	10097	$2.340\times10^{-3}$	20327	$4.890\times10^{-7}$
TRIDIA	100	$1 \times 10^{-7}$	3650	$4.093 \times 10^{-6}$	11716	$3.468 \times 10^{-5}$	22968	$8.918 \times 10^{-9}$

Table A.17. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

			N	EWUOA	FD	L-BFGS	C	D L-BFGS
Problem	n	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
WATSON	12	$1 \times 10^{-1}$	91	$1.602\times 10^{-1}$	260	$2.747\times10^{-1}$	407	$9.362\times 10^{-2}$
WATSON	12	$1 \times 10^{-3}$	170	$7.112\times10^{-2}$	752	$7.265\times10^{-3}$	737	$8.831 \times 10^{-3}$
WATSON	12	$1 \times 10^{-5}$	708	$6.531\times10^{-4}$	1487	$2.403\times10^{-3}$	1470	$1.808\times10^{-4}$
WATSON	12	$1 \times 10^{-7}$	1664	$2.355\times10^{-4}$	1077	$1.868\times 10^{-4}$	1559	$1.367 \times 10^{-5}$
WATSON	31	$1 \times 10^{-1}$	283	1.444	567	2.596	1379	$2.727 \times 10^{-1}$
WATSON	31	$1 \times 10^{-3}$	982	$1.556 \times 10^{-1}$	1465	$1.648 \times 10^{-1}$	4316	$6.367\times10^{-3}$
WATSON	31	$1 \times 10^{-5}$	4133	$4.936 \times 10^{-3}$	2191	$1.682\times 10^{-2}$	3826	$1.420 \times 10^{-3}$
WATSON	31	$1 \times 10^{-7}$	13951	$1.468\times10^{-4}$	4582	$8.715\times10^{-4}$	8624	$1.013\times 10^{-4}$
WOODS	4	$1 \times 10^{-1}$	86	8.655	378	$5.200  imes 10^{-1}$	391	$2.498\times 10^{-1}$
WOODS	4	$1 \times 10^{-3}$	87	7.873	300	$1.818\times 10^{-1}$	567	$1.248\times10^{-5}$
WOODS	4	$1 \times 10^{-5}$	432	$1.680  imes 10^{-5}$	786	$1.295  imes 10^{-3}$	513	$6.023  imes 10^{-7}$
WOODS	4	$1 \times 10^{-7}$	515	$4.005\times 10^{-7}$	278	$3.224\times 10^{-6}$	391	$6.204\times10^{-10}$
WOODS	100	$1 \times 10^{-1}$	8056	$1.617  imes 10^2$	5272	$4.980  imes 10^1$	5925	6.411
WOODS	100	$1 \times 10^{-3}$	27605	$4.700\times10^{-1}$	9371	$4.582\times 10^{-1}$	7344	$1.666\times 10^{-3}$
WOODS	100	$1 \times 10^{-5}$	35366	$1.415\times 10^{-2}$	4207	$4.455\times10^{-3}$	9097	$9.510  imes 10^{-6}$
WOODS	100	$1 \times 10^{-7}$	50000*	$6.216\times10^{-2}$	4756	$7.326  imes 10^{-5}$	7528	$3.725\times 10^{-7}$
ZANGWIL2	2	$1 \times 10^{-1}$	7	$1.455  imes 10^{-3}$	14	$2.786\times10^{-2}$	31	$1.009 \times 10^{-3}$
ZANGWIL2	2	$1 \times 10^{-3}$	17	$1.572 \times 10^{-5}$	13	$1.557  imes 10^{-3}$	31	$1.987\times10^{-6}$
ZANGWIL2	2	$1 \times 10^{-5}$	21	$6.543  imes 10^{-8}$	13	$1.164\times 10^{-6}$	516	$3.655\times 10^{-9}$
ZANGWIL2	2	$1  imes 10^{-7}$	16	$3.964\times10^{-10}$	13	$1.167\times 10^{-8}$	516	$-9.150\times10^{-11}$

Table A.18. Noisy Unconstrained CUTEst Problems Tested. n is the number of variables.  $\sigma_f$  is the standard deviation of the noise.

the number of function evaluations. We removed the limit on the number of evaluations for COBYLA to see if it converges to a different local solution. For HS67, COBYLA terminates with a final objective function value of -1116.415 after 32606 evaluations, and for CRESC4 it converges to a feasible solution with f = 1.03523 after 4706634 evaluations.

(*iii*) One of the solvers terminated at an infeasible point. There are seven problems for which KNITRO terminated at a feasible point but COBYLA at an infeasible one. For all of these problems, COBYLA hits the limit on the number of function evaluations; it can make further progress in feasibility if the budget of evaluations is increased. For problems HS101, HS102, and HS103, KNITRO also hits the evaluation limit; however, it gets a better solution –a feasible one with a lower objective value– for all three problems.

			LMDER		DFOLS
Problem	(n,m)	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
LINEAR(FR)	(9, 45)	21	$1.421 \times 10^{-14}$	45	$-1.421 \times 10^{-14}$
LINEAR(FR)	(9, 45)	21	$1.279\times10^{-13}$	49	0.000
LINEAR(R1)	(7, 35)	16	$-3.099 \times 10^{-7}$	61	$-3.099 \times 10^{-7}$
LINEAR(R1)	(7, 35)	16	$-3.099 \times 10^{-7}$	56	$-3.099 \times 10^{-7}$
LINEAR (R10RC)	(7, 35)	16	$1.493\times 10^{-8}$	56	$1.493\times10^{-8}$
LINEAR (R10RC)	(7, 35)	16	$1.493\times 10^{-8}$	62	$1.493\times 10^{-8}$
ROSENBR	(2,2)	39	0.000	50	$3.788\times10^{-24}$
ROSENBR	(2,2)	17	0.000	14	$1.498\times10^{-20}$
HELIX	(3,3)	49	$3.941 \times 10^{-59}$	44	$2.532\times10^{-28}$
HELIX	(3,3)	57	$5.921\times10^{-47}$	80	$5.476 \times 10^{-20}$
POWELLSG	(4, 4)	331	$6.610 \times 10^{-35}$	54	$1.074 \times 10^{-18}$
POWELLSG	(4, 4)	346	$7.106 \times 10^{-35}$	47	$1.467 \times 10^{-14}$
FREUROTH	(2,2)	60	$3.688\times 10^{-6}$	122	$3.679\times10^{-6}$
FREUROTH	(2,2)	87	$3.793\times10^{-6}$	107	$3.679\times10^{-6}$
BARD	(3, 15)	21	$3.066 \times 10^{-10}$	32	$3.066 \times 10^{-10}$
BARD	(3, 15)	169	$1.742 \times 10^1$	107	$1.066 \times 10^{-1}$
KOWOSB	(4, 11)	99	$4.429\times10^{-12}$	77	$3.849\times10^{-12}$
MEYER	(3, 16)	456	$-4.829\times10^{-6}$	1500*	$3.715  imes 10^4$
WATSON	(6, 31)	57	$5.359\times10^{-11}$	81	$5.355\times10^{-11}$
WATSON	(6, 31)	85	$5.357 \times 10^{-11}$	105	$5.355 \times 10^{-11}$
WATSON	(9, 31)	144	$1.393\times10^{-13}$	409	$1.468 \times 10^{-13}$
WATSON	(9, 31)	134	$1.382 \times 10^{-13}$	499	$1.454 \times 10^{-13}$
WATSON	(12, 31)	276	$1.198 \times 10^{-15}$	340	$2.214\times10^{-9}$
WATSON	(12, 31)	181	$3.636 \times 10^{-16}$	1013	$2.246\times10^{-9}$
BOX3D	(3, 10)	25	$3.390\times10^{-32}$	52	$1.921\times10^{-27}$
JENSMP	(2, 10)	46	$-1.764 \times 10^{-5}$	70	$-1.764 \times 10^{-5}$
BROWNDEN	(4, 20)	94	$1.684\times10^{-3}$	167	$1.626\times 10^{-3}$
BROWNDEN	(4, 20)	96	$1.733 \times 10^{-3}$	189	$1.626 \times 10^{-3}$
CHEBYQUAD	(6,6)	59	$4.434 \times 10^{-32}$	54	$1.239 \times 10^{-23}$
CHEBYQUAD	(7,7)	58	$6.915 \times 10^{-32}$	49	$5.120 \times 10^{-29}$
CHEBYQUAD	(8,8)	168	$-2.725 \times 10^{-10}$	258	$-2.743 \times 10^{-10}$
CHEBYQUAD	(9,9)	94	$4.143 \times 10^{-32}$	87	$4.535 \times 10^{-28}$
CHEBYQUAD	(10, 10)	108	$1.731 \times 10^{-3}$	191	$-3.036 \times 10^{-10}$
CHEBYQUAD	(11, 11)	320	$-4.428 \times 10^{-10}$	352	$-4.481 \times 10^{-10}$
BROWNAL	(10, 10)	79	$2.840 \times 10^{-29}$	52	$7.676 \times 10^{-19}$
OSBORNE1	(5, 33)	157	$-3.025 \times 10^{-12}$	783	$3.737 \times 10^{-12}$
OSBORNE2	(11, 65)	101	$-3.705 \times 10^{-9}$	150	$-3.706 \times 10^{-9}$
OSBORNE2	(11, 65)	148	1.750	138	1.750
BDQRTIC	(8,8)	508	$3.759 \times 10^{-6}$	336	$8.111 \times 10^{-6}$
BDQRTIC	(10, 12)	662	$3.152 \times 10^{-6}$	457	$3.874 \times 10^{-6}$
BDQRTIC	(11, 14)	1046	$3.296 \times 10^{-6}$	554	$2.046 \times 10^{-5}$
BDQRTIC	(12, 16)	1119	$-1.268 \times 10^{-6}$	561	$5.775 \times 10^{-5}$

Table A.19. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

		]	LMDER	DFOLS		
Problem	(n,m)	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$	
CUBE	(5,5)	388	0.000	295	$3.119\times 10^{-18}$	
CUBE	(6, 6)	1024	0.000	714	$4.127\times 10^{-24}$	
CUBE	(8,8)	$4008^{*}$	$2.090\times 10^{-8}$	$4000^{*}$	$4.753 \times 10^{-11}$	
MANCINO	(5,5)	19	$4.224\times10^{-22}$	25	$6.176\times10^{-20}$	
MANCINO	(5,5)	25	$4.686 \times 10^{-22}$	27	$8.536\times10^{-18}$	
MANCINO	(8,8)	28	$5.795 \times 10^{-22}$	23	$3.152\times10^{-12}$	
MANCINO	(10, 10)	34	$8.218\times10^{-22}$	25	$3.069\times10^{-11}$	
MANCINO	(12, 12)	53	$1.322\times 10^{-22}$	32	$2.454\times10^{-17}$	
MANCINO	(12, 12)	53	$5.201\times10^{-22}$	35	$2.798\times10^{-9}$	
HEART8LS	(8,8)	37	$3.402\times10^{-30}$	46	$1.082\times10^{-23}$	
HEART8LS	(8,8)	109	$3.402\times10^{-30}$	2293	$1.001\times10^{-10}$	

Table A.20. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

For two problems, COBYLA terminated at a feasible point but KNITRO at an infeasible one. For problem POLAK6, KNITRO stalls at an infeasible point. For SPIRAL, it gets close to feasibility but runs out its evaluation budget before it can reduce the feasibility error further.

(*iv*) Both solvers terminated at infeasible points. For all three of those problems, KNITRO declares convergence to infeasible stationary points. For problems S365 and BURKEHAN, COBYLA terminates due to the condition on the final trust region radius whereas for POLAK2 it stops due to the function evaluation limit.

			LMDER			DFOLS
Problem	(n,m)	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
LINEAR(FR)	(9, 45)	$1 \times 10^{-1}$	65	$3.565 \times 10^{1}$	70	$3.641 \times 10^{1}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-3}$	126	$5.008 \times 10^{-2}$	97	$6.606 \times 10^{-2}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-5}$	115	$5.170 \times 10^{-4}$	81	$2.889 \times 10^{-5}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-7}$	114	$5.523 \times 10^{-6}$	76	$2.296 \times 10^{-6}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-1}$	229	$2.822 \times 10^2$	77	$1.112 \times 10^3$
LINEAR(FR)	(9, 45)	$1 \times 10^{-3}$	115	$6.086 \times 10^{-2}$	89	$2.179 \times 10^{-2}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-5}$	115	$5.482 \times 10^{-4}$	78	$2.305  imes 10^{-4}$
LINEAR(FR)	(9, 45)	$1 \times 10^{-7}$	114	$5.546  imes 10^{-6}$	76	$2.240\times10^{-6}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-1}$	69	$6.470\times10^{-2}$	74	3.081
LINEAR(R1)	(7, 35)	$1 \times 10^{-3}$	65	$3.221 \times 10^{-3}$	84	$1.482 \times 10^{-4}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-5}$	70	$-2.904\times10^{-7}$	74	$-2.969 \times 10^{-7}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-7}$	108	$-3.099 \times 10^{-7}$	73	$-3.098 \times 10^{-7}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-1}$	68	$6.471 \times 10^{-2}$	78	2.155
LINEAR(R1)	(7, 35)	$1 \times 10^{-3}$	118	$7.437 \times 10^{-5}$	86	$1.577 \times 10^{-4}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-5}$	106	$-3.049 \times 10^{-7}$	84	$-2.926 \times 10^{-7}$
LINEAR(R1)	(7, 35)	$1 \times 10^{-7}$	103	$-3.099 \times 10^{-7}$	79	$-3.099 \times 10^{-7}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-1}$	63	$3.666\times 10^{-1}$	71	$8.280\times10^{-1}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-3}$	82	$7.472\times10^{-6}$	82	$2.771\times 10^{-4}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-5}$	106	$2.101\times 10^{-8}$	75	$2.978\times 10^{-8}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-7}$	72	$1.493\times 10^{-8}$	73	$1.493\times 10^{-8}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-1}$	62	$3.666\times 10^{-1}$	79	$9.209\times10^{-1}$
LINEAR (R10RC)	(7, 35)	$1 \times 10^{-3}$	67	$1.272 \times 10^{-3}$	89	$2.258\times 10^{-4}$
LINEAR(R10RC)	(7, 35)	$1  imes 10^{-5}$	106	$2.103\times10^{-8}$	82	$1.660\times 10^{-8}$
LINEAR(R10RC)	(7, 35)	$1 \times 10^{-7}$	109	$1.495\times 10^{-8}$	75	$1.985\times 10^{-8}$
ROSENBR	(2, 2)	$1 \times 10^{-1}$	44	$7.600 \times 10^{-2}$	51	$4.584 \times 10^{-1}$
ROSENBR	(2, 2)	$1 \times 10^{-3}$	77	$3.078 \times 10^{-7}$	50	$6.040 \times 10^{-8}$
ROSENBR	(2, 2)	$1 \times 10^{-5}$	67	$4.501 \times 10^{-10}$	55	$1.122 \times 10^{-10}$
ROSENBR	(2, 2)	$1 \times 10^{-7}$	65	$3.663 \times 10^{-14}$	49	$2.148 \times 10^{-14}$
ROSENBR	(2, 2)	$1 \times 10^{-1}$	105	$4.118 \times 10^{-3}$	71	$3.578 \times 10^{-2}$
ROSENBR	(2, 2)	$1 \times 10^{-3}$	49	$7.177 \times 10^{-7}$	35	$6.097 \times 10^{-11}$
ROSENBR	(2, 2)	$1 \times 10^{-5}$	44	$4.539 \times 10^{-10}$	54	$3.365 \times 10^{-10}$
ROSENBR	(2, 2)	$1 \times 10^{-7}$	29	$3.064 \times 10^{-14}$	28	$2.191 \times 10^{-15}$
HELIX	(3,3)	$1 \times 10^{-1}$	61	$4.158 \times 10^{-2}$	59	$5.023 \times 10^{-2}$
HELIX	(3,3)	$1 \times 10^{-3}$	61	$1.328 \times 10^{-6}$	62	$1.265 \times 10^{-5}$
HELIX	(3,3)	$1 \times 10^{-5}$	71	$7.012 \times 10^{-10}$	66	$7.646 \times 10^{-10}$
HELIX	(3,3)	$1 \times 10^{-7}$	70	$4.893 \times 10^{-14}$	77	$2.031 \times 10^{-14}$
HELIX	(3,3)	$1 \times 10^{-1}$	61	1.767	66	$8.701 \times 10^{-3}$
HELIX	(3,3)	$1 \times 10^{-3}$	75	$5.707 \times 10^{-6}$	76	$1.429 \times 10^{-6}$
HELIX	(3,3)	$1 \times 10^{-5}$	73	$4.865 \times 10^{-10}$	74	$8.032 \times 10^{-10}$
HELIX	(3,3)	$1 \times 10^{-7}$	70	$5.016 \times 10^{-14}$	47	$3.085 \times 10^{-16}$
POWELLSG	(4, 4)	$1 \times 10^{-1}$	78	$6.286 \times 10^{-2}$	62	$9.436 \times 10^{-2}$
POWELLSG	(4, 4)	$1 \times 10^{-3}$	82	$3.495 \times 10^{-6}$	54	$7.641 \times 10^{-5}$
POWELLSG	(4, 4)	$1 \times 10^{-5}$	86	$3.364 \times 10^{-9}$	54	$7.255 \times 10^{-10}$
POWELLSG	(4, 4)	$1 \times 10^{-7}$	131	$2.803 \times 10^{-14}$	55	$4.174 \times 10^{-15}$
POWELLSG	(4, 4)	$1 \times 10^{-1}$	79	$6.338 \times 10^{-2}$	59	$2.036 \times 10^{-3}$
POWELLSG	(4, 4)	$1 \times 10^{-3}$	86	$3.404 \times 10^{-5}$	64	$2.126 \times 10^{-7}$
POWELLSG	(4, 4)	$1 \times 10^{-5}$	131	$2.607 \times 10^{-10}$	61	$1.502 \times 10^{-8}$
POWELLSG	(4, 4)	$1 \times 10^{-7}$	121	$1.694 \times 10^{-13}$	56	$2.861 \times 10^{-15}$

Table A.21. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

				LMDER		DFOLS
Problem	(n,m)	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
FREUROTH	(2,2)	$1 \times 10^{-1}$	42	1.724	68	$4.634\times10^{-2}$
FREUROTH	(2,2)	$1 \times 10^{-3}$	48	$7.582 \times 10^{-3}$	75	$4.555 \times 10^{-6}$
FREUROTH	(2,2)	$1 \times 10^{-5}$	59	$7.704 \times 10^{-5}$	113	$3.717 \times 10^{-6}$
FREUROTH	(2,2)	$1 \times 10^{-7}$	64	$4.523 \times 10^{-6}$	91	$3.679 \times 10^{-6}$
FREUROTH	(2,2)	$1 \times 10^{-1}$	55	2.546	83	$1.310 \times 10^{-2}$
FREUROTH	(2,2)	$1 \times 10^{-3}$	66	$1.265 \times 10^{-2}$	78	$1.331 \times 10^{-5}$
FREUROTH	(2,2)	$1 \times 10^{-5}$	80	$4.987 \times 10^{-5}$	102	$3.681 \times 10^{-6}$
FREUROTH	(2,2)	$1 \times 10^{-7}$	91	$7.027\times10^{-6}$	122	$3.679 \times 10^{-6}$
BARD	(3, 15)	$1 \times 10^{-1}$	45	$2.370\times10^{-1}$	46	$5.987\times10^{-2}$
BARD	(3, 15)	$1 \times 10^{-3}$	41	$1.168\times10^{-5}$	38	$4.673 \times 10^{-5}$
BARD	(3, 15)	$1 \times 10^{-5}$	53	$4.713\times10^{-8}$	34	$1.404 \times 10^{-9}$
BARD	(3, 15)	$1 \times 10^{-7}$	44	$8.436 \times 10^{-10}$	35	$3.068 \times 10^{-10}$
BARD	(3, 15)	$1 \times 10^{-1}$	41	$1.355 \times 10^1$	56	$9.256 \times 10^{-2}$
BARD	(3, 15)	$1 \times 10^{-3}$	73	$2.035  imes 10^1$	65	$4.982\times 10^{-5}$
BARD	(3, 15)	$1 \times 10^{-5}$	98	2.282	54	$1.114\times10^{-9}$
BARD	(3, 15)	$1 \times 10^{-7}$	115	$1.611 \times 10^1$	102	$1.066  imes 10^{-1}$
KOWOSB	(4, 11)	$1 \times 10^{-1}$	35	$5.005 \times 10^{-3}$	42	$3.088\times10^{-2}$
KOWOSB	(4, 11)	$1 \times 10^{-3}$	57	$1.418\times 10^{-4}$	44	$1.561\times 10^{-4}$
KOWOSB	(4, 11)	$1 \times 10^{-5}$	92	$4.081\times 10^{-8}$	60	$4.916\times 10^{-8}$
KOWOSB	(4, 11)	$1 \times 10^{-7}$	84	$1.302\times 10^{-8}$	64	$2.981 \times 10^{-10}$
MEYER	(3, 16)	$1 \times 10^{-1}$	182	$8.301 \times 10^4$	67	$1.113\times 10^5$
MEYER	(3, 16)	$1 \times 10^{-3}$	666	4.770	252	$7.380 \times 10^4$
MEYER	(3, 16)	$1 \times 10^{-5}$	504	$8.423\times10^{-3}$	1498	$4.022 \times 10^4$
MEYER	(3, 16)	$1 \times 10^{-7}$	497	$-1.300 \times 10^{-6}$	$1500^{*}$	$4.031 \times 10^4$
WATSON	(6, 31)	$1 \times 10^{-1}$	96	$1.592 \times 10^{-1}$	71	$1.232 \times 10^{-1}$
WATSON	(6, 31)	$1 \times 10^{-3}$	83	$1.818 \times 10^{-4}$	79	$1.328 \times 10^{-4}$
WATSON	(6, 31)	$1 \times 10^{-5}$	82	$6.905 \times 10^{-5}$	82	$7.375 \times 10^{-7}$
WATSON	(6, 31)	$1 \times 10^{-7}$	103	$5.333 \times 10^{-9}$	78	$8.703 \times 10^{-10}$
WATSON	(6, 31)	$1 \times 10^{-1}$	95	$2.562 \times 10^{-1}$	80	$1.150 \times 10^{-1}$
WATSON	(6, 31)	$1 \times 10^{-3}$	236	$4.267 \times 10^{-4}$	82	$2.052 \times 10^{-4}$
WATSON	(6, 31)	$1 \times 10^{-5}$	116	$5.311 \times 10^{-5}$	87	$1.676 \times 10^{-7}$
WATSON	(6, 31)	$1 \times 10^{-7}$	116	$2.197 \times 10^{-8}$	85	$6.039 \times 10^{-10}$
WATSON	(9, 31)	$1 \times 10^{-1}$	95	$1.791 \times 10^{-1}$	95	$6.326 \times 10^{-2}$
WATSON	(9, 31)	$1 \times 10^{-3}$	183	$2.761 \times 10^{-4}$	113	$1.112 \times 10^{-4}$
WATSON	(9, 31)	$1 \times 10^{-5}$	386	$3.603 \times 10^{-5}$	114	$5.717 \times 10^{-6}$
WATSON	(9, 31)	$1 \times 10^{-7}$	531	$1.175 \times 10^{-7}$	220	$9.384 \times 10^{-10}$
WATSON	(9, 31)	$1 \times 10^{-1}$	134	1.338	128	$3.406 \times 10^{-2}$
WATSON	(9, 31)	$1 \times 10^{-3}$	390	$7.293 \times 10^{-3}$	115	$2.543 \times 10^{-4}$
WATSON	(9, 31)	$1 \times 10^{-5}$	580	$1.436 \times 10^{-4}$	113	$4.061 \times 10^{-6}$
WATSON	(9, 31)	$1 \times 10^{-7}$	533	$4.147 \times 10^{-7}$	337	$5.464 \times 10^{-9}$
WATSON	(12, 31)	$1 \times 10^{-1}$	147	$1.829 \times 10^{-1}$	133	$4.518 \times 10^{-1}$
WATSON	(12, 31)	$1 \times 10^{-5}$	199	$1.942 \times 10^{-4}$	129	$4.431 \times 10^{-4}$
WATSON	(12, 31)	$1 \times 10^{-5}$	736	$1.863 \times 10^{-3}$	130	$1.202 \times 10^{-7}$
WATSON	(12, 31)	$1 \times 10^{-7}$	713	$4.164 \times 10^{-6}$	234	$1.337 \times 10^{-8}$
WATSON	(12, 31)	$1 \times 10^{-1}$	188	$7.230 \times 10^{-1}$	127	$1.000 \times 10^{-1}$
WATSON	(12, 31)	$1 \times 10^{-5}$ 1 × 10^{-5}	255	$1.080 \times 10^{-2}$	139	$2.802 \times 10^{-4}$
WAISON	(12, 31)	$1 \times 10^{-5}$ $1 \times 10^{-7}$	530	$3.237 \times 10^{-2}$	196	$1.911 \times 10^{-8}$
WAISON	(12, 31)	$1 \times 10^{-7}$	851	$1.854 \times 10^{-4}$	183	$1.500 \times 10^{-8}$

Table A.22. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

			]	LMDER		DFOLS
Problem	(n,m)	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
BOX3D	(3, 10)	$1 \times 10^{-1}$	44	$1.674\times 10^{-1}$	53	$1.181 \times 10^{-1}$
BOX3D	(3, 10)	$1 \times 10^{-3}$	63	$1.532\times 10^{-3}$	43	$5.394\times10^{-6}$
BOX3D	(3, 10)	$1 \times 10^{-5}$	56	$1.738 \times 10^{-9}$	46	$5.818 \times 10^{-12}$
BOX3D	(3, 10)	$1 \times 10^{-7}$	55	$1.758 \times 10^{-13}$	57	$2.248 \times 10^{-13}$
JENSMP	(2, 10)	$1 \times 10^{-1}$	46	$3.728  imes 10^3$	47	$2.684\times 10^{-1}$
JENSMP	(2, 10)	$1 \times 10^{-3}$	44	$4.102\times10^{-2}$	51	$1.660 \times 10^{-2}$
JENSMP	(2, 10)	$1 \times 10^{-5}$	53	$1.445  imes 10^{-3}$	59	$6.415 imes10^{-5}$
JENSMP	(2, 10)	$1 \times 10^{-7}$	51	$-7.215\times10^{-6}$	47	$-8.720 \times 10^{-6}$
BROWNDEN	(4, 20)	$1 \times 10^{-1}$	75	$1.521  imes 10^3$	84	$3.922  imes 10^2$
BROWNDEN	(4, 20)	$1 \times 10^{-3}$	86	$3.277  imes 10^1$	97	2.897
BROWNDEN	(4, 20)	$1 \times 10^{-5}$	98	$7.498\times10^{-2}$	140	$9.534\times10^{-3}$
BROWNDEN	(4, 20)	$1 \times 10^{-7}$	105	$2.529\times 10^{-3}$	144	$4.056 imes10^{-3}$
BROWNDEN	(4, 20)	$1 \times 10^{-1}$	90	$2.384 \times 10^3$	102	$8.007  imes 10^2$
BROWNDEN	(4, 20)	$1 \times 10^{-3}$	102	$4.493  imes 10^1$	97	3.732
BROWNDEN	(4, 20)	$1 \times 10^{-5}$	117	$4.912\times 10^{-2}$	115	$1.440 \times 10^{-1}$
BROWNDEN	(4, 20)	$1 \times 10^{-7}$	114	$2.743 imes10^{-3}$	130	$2.798\times10^{-3}$
CHEBYQUAD	(6, 6)	$1 \times 10^{-1}$	54	$4.243\times10^{-2}$	64	$2.304\times10^{-2}$
CHEBYQUAD	(6, 6)	$1 \times 10^{-3}$	57	$3.020\times 10^{-2}$	73	$9.161\times 10^{-7}$
CHEBYQUAD	(6, 6)	$1 \times 10^{-5}$	93	$4.322\times10^{-10}$	61	$1.447\times10^{-11}$
CHEBYQUAD	(6, 6)	$1 \times 10^{-7}$	91	$4.785 \times 10^{-14}$	63	$1.792 \times 10^{-13}$
CHEBYQUAD	(7,7)	$1 \times 10^{-1}$	46	$3.377 imes10^{-2}$	67	$1.045\times 10^{-2}$
CHEBYQUAD	(7, 7)	$1 \times 10^{-3}$	118	$6.182\times10^{-6}$	64	$1.901\times10^{-5}$
CHEBYQUAD	(7,7)	$1 \times 10^{-5}$	97	$9.566 \times 10^{-10}$	73	$2.030 \times 10^{-10}$
CHEBYQUAD	(7, 7)	$1 \times 10^{-7}$	95	$1.043 \times 10^{-13}$	73	$6.456 \times 10^{-14}$
CHEBYQUAD	(8, 8)	$1 \times 10^{-1}$	69	$2.683\times10^{-2}$	79	$6.164\times10^{-2}$
CHEBYQUAD	(8, 8)	$1 \times 10^{-3}$	93	$1.261\times 10^{-2}$	96	$1.182\times10^{-3}$
CHEBYQUAD	(8, 8)	$1 \times 10^{-5}$	128	$1.039\times 10^{-3}$	199	$6.120\times10^{-7}$
CHEBYQUAD	(8, 8)	$1 \times 10^{-7}$	240	$6.147\times10^{-7}$	208	$1.288\times 10^{-9}$
CHEBYQUAD	(9,9)	$1 \times 10^{-1}$	$54^{**}$	$2.888\times 10^{-2}$	86	$2.597\times10^{-2}$
CHEBYQUAD	(9,9)	$1 \times 10^{-3}$	100	$2.013\times10^{-2}$	100	$4.981\times10^{-5}$
CHEBYQUAD	(9,9)	$1 \times 10^{-5}$	134	$2.326\times10^{-9}$	102	$6.629 \times 10^{-10}$
CHEBYQUAD	(9, 9)	$1 \times 10^{-7}$	161	$3.856 \times 10^{-14}$	104	$3.012 \times 10^{-14}$
CHEBYQUAD	(10, 10)	$1 \times 10^{-1}$	94	$4.026\times 10^{-2}$	89	$2.091\times 10^{-2}$
CHEBYQUAD	(10, 10)	$1 \times 10^{-3}$	153	$9.237  imes 10^{-4}$	87	$4.596\times10^{-3}$
CHEBYQUAD	(10, 10)	$1 \times 10^{-5}$	163	$5.114\times10^{-6}$	113	$4.622\times 10^{-5}$
CHEBYQUAD	(10, 10)	$1 \times 10^{-7}$	207	$1.732 \times 10^{-3}$	163	$5.151\times10^{-8}$
CHEBYQUAD	(11, 11)	$1 \times 10^{-1}$	78	$2.983\times10^{-2}$	87	$1.920\times 10^{-2}$
CHEBYQUAD	(11, 11)	$1 \times 10^{-3}$	104	$9.225\times10^{-3}$	93	$6.273\times10^{-3}$
CHEBYQUAD	(11, 11)	$1 \times 10^{-5}$	213	$4.205\times10^{-4}$	235	$8.670\times10^{-7}$
CHEBYQUAD	(11, 11)	$1 \times 10^{-7}$	449	$1.141\times 10^{-6}$	246	$1.473\times 10^{-8}$

Table A.23. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

			]	LMDER		DFOLS
Problem	(n,m)	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
BROWNAL	(10, 10)	$1 \times 10^{-1}$	127	1.419	89	$1.655\times 10^2$
BROWNAL	(10, 10)	$1 \times 10^{-3}$	147	$5.220 \times 10^{-7}$	99	$3.284\times 10^{-5}$
BROWNAL	(10, 10)	$1 \times 10^{-5}$	155	$4.102 \times 10^{-10}$	89	$9.707 \times 10^{-10}$
BROWNAL	(10, 10)	$1 \times 10^{-7}$	131	$1.084\times10^{-13}$	75	$1.308\times10^{-14}$
OSBORNE1	(5, 33)	$1 \times 10^{-1}$	53	1.279	53	5.737
OSBORNE1	(5, 33)	$1 \times 10^{-3}$	75	$3.424 \times 10^{-3}$	86	$4.512\times10^{-1}$
OSBORNE1	(5, 33)	$1 \times 10^{-5}$	86	$1.010\times 10^{-2}$	225	$2.446\times 10^{-2}$
OSBORNE1	(5, 33)	$1 \times 10^{-7}$	376	$4.951\times 10^{-8}$	185	$2.446\times 10^{-2}$
OSBORNE2	(11, 65)	$1 \times 10^{-1}$	113	1.297	104	2.226
OSBORNE2	(11, 65)	$1 \times 10^{-3}$	426	$6.798  imes 10^{-3}$	153	$2.216\times 10^{-4}$
OSBORNE2	(11, 65)	$1 \times 10^{-5}$	308	$9.726\times10^{-5}$	174	$6.715\times10^{-7}$
OSBORNE2	(11, 65)	$1 \times 10^{-7}$	235	$1.654\times 10^{-7}$	148	$8.159\times10^{-10}$
OSBORNE2	(11, 65)	$1 \times 10^{-1}$	124	$2.799  imes 10^1$	111	$2.293  imes 10^1$
OSBORNE2	(11, 65)	$1 \times 10^{-3}$	204	1.754	132	1.750
OSBORNE2	(11, 65)	$1 \times 10^{-5}$	277	1.750	134	1.750
OSBORNE2	(11, 65)	$1 \times 10^{-7}$	223	1.750	139	1.750
BDQRTIC	(8,8)	$1 \times 10^{-1}$	124	5.023	89	4.566
BDQRTIC	(8, 8)	$1 \times 10^{-3}$	128	2.317	103	$2.289\times10^{-1}$
BDQRTIC	(8, 8)	$1 \times 10^{-5}$	265	$2.905\times10^{-3}$	122	$1.864\times10^{-2}$
BDQRTIC	(8, 8)	$1 \times 10^{-7}$	437	$2.902\times10^{-5}$	173	$8.502\times10^{-5}$
BDQRTIC	(10, 12)	$1 \times 10^{-1}$	128	$1.221 \times 10^1$	97	$1.447 \times 10^2$
BDQRTIC	(10, 12)	$1 \times 10^{-3}$	197	$1.021 \times 10^{-1}$	124	1.141
BDQRTIC	(10, 12)	$1 \times 10^{-5}$	340	$5.836 \times 10^{-3}$	153	$3.739\times10^{-2}$
BDQRTIC	(10, 12)	$1 \times 10^{-7}$	574	$6.811\times10^{-5}$	262	$7.248\times10^{-4}$
BDQRTIC	(11, 14)	$1 \times 10^{-1}$	139	7.006	106	$9.772 \times 10^1$
BDQRTIC	(11, 14)	$1 \times 10^{-3}$	188	$1.299\times 10^{-1}$	127	3.141
BDQRTIC	(11, 14)	$1 \times 10^{-5}$	396	$6.226 \times 10^{-3}$	190	$7.543\times10^{-2}$
BDQRTIC	(11, 14)	$1 \times 10^{-7}$	725	$4.880\times 10^{-5}$	333	$2.513\times10^{-3}$
BDQRTIC	(12, 16)	$1 \times 10^{-1}$	137	5.551	125	$1.355  imes 10^2$
BDQRTIC	(12, 16)	$1 \times 10^{-3}$	204	$2.591\times 10^{-1}$	135	2.972
BDQRTIC	(12, 16)	$1 \times 10^{-5}$	283	$1.565\times10^{-2}$	191	$5.335\times10^{-2}$
BDQRTIC	(12, 16)	$1 \times 10^{-7}$	626	$3.785\times10^{-4}$	323	$3.455\times10^{-3}$
CUBE	(5,5)	$1 \times 10^{-1}$	96	$2.386\times10^{-2}$	54	$6.129\times 10^{-2}$
CUBE	(5, 5)	$1 \times 10^{-3}$	112	$3.481 \times 10^{-3}$	90	$1.519\times10^{-4}$
CUBE	(5, 5)	$1 \times 10^{-5}$	250	$1.971 \times 10^{-10}$	367	$4.920 \times 10^{-12}$
CUBE	(5, 5)	$1 \times 10^{-7}$	415	$9.699\times10^{-14}$	272	$1.508\times10^{-13}$
CUBE	(6, 6)	$1 \times 10^{-1}$	127	$1.312\times 10^{-2}$	56	$5.761\times10^{-2}$
CUBE	(6, 6)	$1 \times 10^{-3}$	162	$1.853  imes 10^{-4}$	82	$1.329\times 10^{-4}$
CUBE	(6, 6)	$1 \times 10^{-5}$	832	$8.835\times10^{-7}$	460	$2.319\times10^{-6}$
CUBE	(6, 6)	$1 \times 10^{-7}$	958	$9.762 \times 10^{-14}$	758	$4.059\times10^{-14}$
CUBE	(8, 8)	$1 \times 10^{-1}$	168	$2.402\times10^{-1}$	63	$9.869  imes 10^1$
CUBE	(8, 8)	$1 \times 10^{-3}$	185	$8.640\times 10^{-6}$	87	$7.661\times 10^{-4}$
CUBE	(8, 8)	$1 \times 10^{-5}$	621	$1.251\times 10^{-5}$	263	$1.030\times 10^{-4}$
CUBE	(8, 8)	$1 \times 10^{-7}$	3431	$6.217\times 10^{-8}$	1763	$2.522\times 10^{-7}$

Table A.24. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

			LMDER		DFOLS	
Problem	(n,m)	$\sigma_{f}$	#feval	$\phi(x) - \phi^*$	#feval	$\phi(x) - \phi^*$
MANCINO	(5, 5)	$1 \times 10^{-1}$	47	$7.329\times10^{-2}$	48	$1.334\times 10^{-2}$
MANCINO	(5, 5)	$1 \times 10^{-3}$	45	$7.341\times10^{-6}$	44	$5.405\times10^{-8}$
MANCINO	(5, 5)	$1 \times 10^{-5}$	39	$7.341 \times 10^{-10}$	33	$1.657 \times 10^{-12}$
MANCINO	(5, 5)	$1 \times 10^{-7}$	39	$7.342\times10^{-14}$	25	$3.807\times10^{-15}$
MANCINO	(5, 5)	$1 \times 10^{-1}$	47	$7.435\times10^{-2}$	58	$2.132\times 10^{-2}$
MANCINO	(5, 5)	$1 \times 10^{-3}$	45	$7.511\times10^{-6}$	44	$1.457\times 10^{-5}$
MANCINO	(5, 5)	$1 \times 10^{-5}$	45	$1.247\times 10^{-9}$	27	$4.630 \times 10^{-10}$
MANCINO	(5, 5)	$1 \times 10^{-7}$	45	$8.751 \times 10^{-14}$	27	$4.120 \times 10^{-14}$
MANCINO	(8, 8)	$1 \times 10^{-1}$	72	$1.550\times10^{-1}$	82	$1.228\times10^{-1}$
MANCINO	(8, 8)	$1 \times 10^{-3}$	78	$1.130 \times 10^{-7}$	57	$8.045\times10^{-6}$
MANCINO	(8,8)	$1 \times 10^{-5}$	60	$1.549 \times 10^{-9}$	46	$2.898 \times 10^{-10}$
MANCINO	(8, 8)	$1 \times 10^{-7}$	60	$1.549 \times 10^{-13}$	23	$7.222 \times 10^{-13}$
MANCINO	(10, 10)	$1 \times 10^{-1}$	109	$8.335 \times 10^{-5}$	82	$7.748 \times 10^{-3}$
MANCINO	(10, 10)	$1 \times 10^{-3}$	107	$8.191 \times 10^{-9}$	59	$8.847\times10^{-6}$
MANCINO	(10, 10)	$1 \times 10^{-5}$	74	$6.522 \times 10^{-10}$	45	$1.015 \times 10^{-10}$
MANCINO	(10, 10)	$1 \times 10^{-7}$	74	$6.526 \times 10^{-14}$	36	$5.372 \times 10^{-15}$
MANCINO	(12, 12)	$1 \times 10^{-1}$	114	$3.374 \times 10^{-4}$	107	$3.447\times10^{-2}$
MANCINO	(12, 12)	$1 \times 10^{-3}$	101	$3.354 imes10^{-8}$	77	$8.807 \times 10^{-6}$
MANCINO	(12, 12)	$1 \times 10^{-5}$	101	$3.353 \times 10^{-12}$	60	$9.663 \times 10^{-10}$
MANCINO	(12, 12)	$1 \times 10^{-7}$	101	$3.357 \times 10^{-16}$	32	$9.721 \times 10^{-14}$
MANCINO	(12, 12)	$1 \times 10^{-1}$	114	$3.374 \times 10^{-4}$	126	$9.117\times10^{-2}$
MANCINO	(12, 12)	$1 \times 10^{-3}$	114	$3.348 \times 10^{-8}$	72	$9.369 \times 10^{-6}$
MANCINO	(12, 12)	$1 \times 10^{-5}$	101	$3.357 \times 10^{-12}$	35	$4.632 \times 10^{-9}$
MANCINO	(12, 12)	$1 \times 10^{-7}$	101	$3.452 \times 10^{-16}$	35	$2.607 \times 10^{-9}$
HEART8LS	(8,8)	$1 \times 10^{-1}$	113	$6.289 \times 10^{-2}$	71	$1.356 \times 10^{-1}$
HEART8LS	(8,8)	$1 \times 10^{-3}$	83	$1.244 \times 10^{-6}$	66	$3.717 \times 10^{-6}$
HEART8LS	(8,8)	$1 \times 10^{-5}$	82	$1.508 \times 10^{-11}$	66	$3.928 \times 10^{-12}$
HEART8LS	(8,8)	$1 \times 10^{-7}$	79	$1.167 \times 10^{-15}$	67	$1.275 \times 10^{-14}$
HEART8LS	(8,8)	$1 \times 10^{-1}$	156	5.632	86	5.270
HEART8LS	(8, 8)	$1 \times 10^{-3}$	566	$7.222 \times 10^{-6}$	77	4.911
HEART8LS	(8,8)	$1 \times 10^{-5}$	164	$1.102 \times 10^{-9}$	172	4.794
HEART8LS	(8,8)	$1 \times 10^{-7}$	161	$5.788 \times 10^{-14}$	376	4.428

Table A.25. Benchmarking Problems Tested. n is the number of variables and m is the dimension of the residual vector.

		Number of			Number of		
	Objective	Variables	One-sided	Two-sided	Constraints	Linear	Nonlinear
Problem	Type	(n)	Bounds	Bounds	(m)	Ineq.	Ineq.
BURKEHAN	quadratic	1	1	0	1	0	1
CANTILVR	linear	5	5	0	1	0	1
CB2	linear	3	0	0	3	0	3
CB3	linear	3	0	0	3	0	3
CHACONN1	linear	3	Ő	Ő	3	Ő	3
CHACONN2	linear	3	Ő	ů 0	3	Õ	3
CRESC4	general	6	3	1	8	0	8
CRESC50	general	6	3	1	100	0	100
DEMBO7	quadratic	16	0	16	20	1	19
DIPICRI	general	7	0	10	4	0	10
CICOMEZ2	linoar	3	0	0	3	0	
CICOME73	linear	3	0	0	3	0	J 2
HALDMADS	linear	5	0	0	42	0	12
HALDWADS	roporol	0 7	0	0	42	0	42
HS100 HS100MOD	general	7	0	0	4	0	4
IIST00MOD	general	7	0	0 7	4 F	0	4
П5101 ЦС109	general		0	1	5	0	5 F
H5102	general		0	1	5	0	5 F
HS103	general	7	0	7	5	0	5
HS104	general	8	0	8	5	0	5
HS13	quadratic	2	2	0	1	0	1
HS34	linear	3	0	3	2	0	2
HS64	general	3	3	0	1	0	1
HS66	linear	3	0	3	2	0	2
HS67	general	3	0	3	14	0	14
HS72	linear	4	0	4	2	0	2
HS85	general	5	0	5	21	1	20
HS88	quadratic	2	0	0	1	0	1
HS89	quadratic	3	0	0	1	0	1
HS90	quadratic	4	0	0	1	0	1
HS91	quadratic	5	0	0	1	0	1
HS92	quadratic	6	0	0	1	0	1
HS93	general	6	6	0	2	0	2
MADSEN	linear	3	0	0	6	0	6
MATRIX2	quadratic	6	4	0	2	0	2
MINMAXBD	linear	5	0	0	20	0	20
POLAK1	linear	3	0	0	2	0	2
POLAK2	linear	11	0	0	2	0	2
POLAK3	linear	12	Ő	Ő	10	Ő	10
POLAK5	linear	3	Ő	ů 0	2	Õ	2
POLAK6	linear	5	ů 0	0 0	4	0	- 4
S365	quadratic	7	4	0	5	0	5
S365MOD	quadratic	7		0	5	0	5
SNAKE	linoar	9	- - -	0	9	0	0
SPIRAT	linoar	2 2	0	0	2	0	2
SF INAL GVNTHEG1	mear	3 C	0	0		0	2
SINITESI TWODADS	general	0	0	0	0	4	2
TWOBARS	general	2	0	2	2	0	2
WOMFLET	linear	3	0	0	3	0	3

Table A.26. Properties of small-scale CUTEst problems.

Problem	SIF Parameter	Testing Name	d	m
SVANBERG	N = 10	SVANBERGN10	10	10
SVANBERG	N = 50	SVANBERGN50	50	50
SVANBERG	N = 100	SVANBERGN100	100	100
SVANBERG	N = 500	SVANBERGN500	500	500
READING4	N=2	READING4N2	2	2
READING4	N = 50	READING4N50	50	50
READING4	N = 100	READING4N100	100	100
READING4	N = 500	READING4N500	500	500
COSHFUN	M = 3	COSHFUNM3	10	3
COSHFUN	M = 8	COSHFUNM8	25	8
COSHFUN	M = 14	COSHFUNM14	43	14
COSHFUN	M = 20	COSHFUNM20	61	20

Table A.27. Instances of variable-size CUTEst problems.
		KNIT	TRO			COBY	ΊLΑ	
Problem	$\phi(x)$	#feval	CPU	feaserr	$\phi(x)$	#fevals	CPU	feaserr
CB3	2.0000	40	0.023	2.22e-16	2.0000	44	0.157	0.00e+00
CHACONN2	2.0000	40	0.021	0.00e+00	2.0000	48	0.156	0.00e + 00
GIGOMEZ3	2.0000	40	0.022	0.00e+00	2.0000	45	0.160	0.00e + 00
HS100	680.6301	360	0.160	0.00e+00	680.6300	454	0.272	0.00e + 00
DIPIGRI	680.6301	389	0.152	2.84e-14	680.6300	458	0.288	0.00e + 00
HS93	135.0760	2247	1.040	4.44e-16	135.0760	3000*	0.873	0.00e+00
HS64	6299.8424	144	0.093	0.00e+00	6299.8400	409	0.244	0.00e + 00
POLAK3	5.9330	902	0.414	3.89e-16	5.9330	3139	0.071	0.00e + 00
POLAK1	2.7183	60	0.034	0.00e+00	2.7183	326	0.232	0.00e + 00
HS104	3.9512	718	0.282	0.00e+00	3.9512	3977	0.131	0.00e + 00
HS100MOD	678.6796	477	0.207	0.00e+00	678.6790	385	0.240	0.00e + 00
TWOBARS	1.5087	63	0.037	0.00e+00	1.5087	83	0.172	0.00e + 00
HS85	-2.2156	117	0.078	7.11e-15	-2.2156	299	0.242	0.00e + 00
CB2	1.9522	72	0.041	1.11e-16	1.9522	111	0.173	0.00e + 00
CHACONN1	1.9522	55	0.030	0.00e+00	1.9522	120	0.175	0.00e + 00
MADSEN	0.6164	73	0.042	0.00e+00	0.6164	112	0.175	0.00e + 00
HS66	0.5182	43	0.026	4.44e-16	0.5182	101	0.170	0.00e + 00
MINMAXBD	115.7064	420	0.335	6.17e-14	115.7060	924	0.395	0.00e + 00
CANTILVR	1.3400	170	0.078	3.61e-16	1.3400	281	0.227	0.00e + 00
HS92	1.3627	976	0.988	4.60e-17	1.3627	3000*	0.993	0.00e + 00
HALDMADS	0.0001	113	0.073	4.44e-16	0.0001	669	0.341	0.00e + 00
HS34	-0.8340	59	0.033	4.44e-16	-0.8340	47	0.149	0.00e + 00
HS90	1.3627	570	0.475	0.00e+00	1.3629	2000*	0.066	0.00e + 00
SNAKE	0.0000	25	0.029	1.02e-16	0.0000	72	0.165	0.00e + 00
HS72	727.6794	108	0.103	1.73e-18	727.6794	1002	0.424	4.77e-18
GIGOMEZ2	1.9522	79	0.048	2.22e-16	1.9522	111	0.169	1.11e-16
SYNTHES1	0.7593	96	0.041	0.00e+00	0.7593	174	0.186	1.54e-23
HS88	1.3627	209	0.156	0.00e+00	1.3627	167	0.203	2.09e-17
HS13	0.9999	140	0.155	2.61e-13	1.0000	86	0.168	1.07e-27
MATRIX2	0.0000	275	0.131	3.60e-17	0.0000	216	0.199	5.20e-18
WOMFLET	0.0000	103	0.070	5.89e-10	0.0000	117	0.181	8.99e-18
S365	0.0000	63	0.032	1.72e-07	0.0000	210	0.215	4.35e-17
HS67	-1162.1187	285	0.198	0.00e+00	-980.1600	7000*	0.885	0.00e+00
CRESC50	0.5952	16704	17.912	0.00e+00	7.9917	50000*	4.863	0.00e + 00
CRESC4	0.8719	688	0.363	0.00e+00	44.5901	4000*	0.318	0.00e+00
HS91	1.3627	1229	1.057	9.96e-18	1.1199	2500*	0.583	6.50e-04
DEMBO7	174.7870	3225	1.594	3.61e-16	250.0720	10000*	0.238	7.29e-02
POLAK5	50.0000	65	0.043	0.00e+00	34.4929	1500*	0.571	$1.55e \pm 01$
HS101	1809.7691	3500*	1.908	1.38e-13	3000.1900	3500*	0.077	1.99e-01
HS89	1 3627	315	0.257	2 17e-17	0.6520	1500*	0.742	5 12e-02
HS102	911 8816	3500*	1 965	4 67e-12	3000 3500	3500*	0.096	3.60e-01
HS102	543.6680	$3500^{*}$	1.858	2.57e-16	3000.1700	$3500^{*}$	0.000 0.091	1.76e-01
SPIRAL	0.0195	1500*	0.796	1.34e-05	0.0958	1500*	0.508	0.00e+00
POLAK6	-78.6745	474	1.311	8.55e+01	0.0000	2500*	0.765	0.00e+00
S365MOD	0.2500	162	0.204	$1.25e \pm 00$	0.0301	247	0.211	3.26e-01
BURKEHAN	10.0000	13	0.007	1.01e+02	0.0000	54	0.175	1.00e+00
POLAK2	29.9570	2641	5.119	5.39e+01	-259.2500	5500*	0.794	3.14e+02
	1		0.210			2000		JJ

Table A.28. Summary of the results for small-scale noiseless constrained CUTEst problems. The horizontal bars divide cases (i), (ii), (iii), and (iv).

$TOL = 10^{-6}$							TC	$L = 10^{-3}$			$TOL = 10^{-1}$					
			KNIT	RO	COBY	LA		KNIT	RO	COBY	LA		KNIT	RO	COBY	LA
problem	$\phi^*$	target $\phi$	$\phi(x)$	#fevals	$\phi(x)$	#fevals	target $\phi$	$\phi(x)$	#fevals	$\phi(x)$	#fevals	target $\phi$	$\phi(x)$	#fevals	$\phi(x)$	#fevals
CB3	2.000000	2.00000	2.00000	30	2.00000	25	2.00200	2.00000	30	2.00000	25	2.20000	2.00000	30	2.00351	18
CHACONN2	2.000000	2.00000	2.00000	30	2.00000	32	2.00200	2.00000	30	2.00021	24	2.20000	2.00000	30	2.06187	19
GIGOMEZ3	2.000000	2.00000	2.00000	30	2.00000	39	2.00200	2.00000	30	2.00000	39	2.20000	2.00000	30	2.00000	39
HS100	680.630057	680.63074	680.63013	124	680.63000	247	681.31069	680.75511	75	680.63000	247	748.69306	714.00000	1	714.00000	1
DIPIGRI	680.630057	680.63074	680.63013	124	680.63000	410	681.31069	680.75511	75	680.63100	141	748.69306	714.00000	1	714.00000	1
HS93	135.075964	135.07610	135.07610	871	135.07600	2416	135.21104	135.14058	229	135.21000	270	148.58356	137.06640	1	137.06600	1
HS64	6299.842428	6299.84873	6299.73981	64	6299.84000	335	6306.14227	6302.48265	54	6305.80000	316	6929.82667	6687.88225	30	6684.54000	239
POLAK3	5.933003	5.93301	5.93301	469	5.93300	2637	5.93894	5.93353	349	5.93644	2350	6.52630	5.93353	349	5.94016	2328
POLAK1	2.718282	2.71828	2.71828	50	2.71828	265	2.72100	2.71828	50	2.71834	253	2.99011	2.71828	50	2.83811	250
HS104	3.951163	3.95117	3.95117	290	3.95116	1764	3.95511	3.95148	231	3.95381	730	4.34628	3.95946	193	4.15766	173
HS100MOD	678.679638	678.68032	678.67969	173	678.68000	169	679.35832	679.08174	63	679.18500	50	746.54760	714.00000	1	714.00000	1
TWOBARS	1.508652	1.50865	1.50865	50	1.50865	33	1.51016	1.50865	50	1.50892	21	1.65952	1.50865	50	1.65447	10
HS85	-2.215605	-2.21560	-2.21560	110	-2.21560	299	-2.21339	-2.21560	110	-2.21500	247	-1.99404	-2.21560	110	-2.02180	158
CB2	1.952224	1.95223	1.95222	52	1.95222	57	1.95418	1.95222	52	1.95230	41	2.14745	1.95222	52	1.95230	41
CHACONN1	1.952224	1.95223	1.95222	40	1.95222	55	1.95418	1.95222	40	1.95222	55	2.14745	1.95222	40	2.00282	13
MADSEN	0.616432	0.61643	0.61643	58	0.61643	65	0.61743	0.61643	58	0.61643	65	0.71643	0.61643	58	0.65433	22
HS66	0.518163	0.51816	0.51816	33	0.51816	37	0.51916	0.51866	21	0.51817	28	0.61816	0.58000	1	0.58000	1
MINMAXBD	115.706440	115.70656	115.70644	343	115.70600	800	115.82215	115.70644	343	115.72100	766	127.27708	118.74188	246	115.72100	766
CANTILVR	1.339956	1.33996	1.33994	107	1.33995	189	1.34130	1.33994	107	1.33997	130	1.47395	1.33994	107	1.34758	107
HS92	1.362657	1.36266	1.36266	844	1.36265	2121	1.36402	1.36327	405	1.36310	203	1.49892	1.36327	405	1.36310	203
HALDMADS	0.000122	0.00012	0.00012	97	0.00013	669	0.00112	0.00012	97	0.00057	51	0.10012	0.00012	97	0.00057	51
HS34	-0.834032	-0.83403	-0.83403	54	-0.83400	40	-0.83303	-0.83403	54	-0.83400	25	-0.73403	-0.83186	44	-0.83400	25
HS90	1.362657	1.36266	1.36185	259	1.36287	2000	1.36402	1.36185	259	1.36311	159	1.49892	1.36185	259	1.36311	159
SNAKE	0.000000	0.00000	0.00000	17	0.00000	67	0.00100	0.00000	17	0.00009	61	0.10000	0.00000	17	0.00009	61
HS72	727.679358	727.68009	727.66249	72	727.67973	877	728.40704	727.66249	72	728.36304	856	800.44729	727.66249	72	738.77714	832
GIGOMEZ2	1.952224	1.95223	1.95222	64	1.95223	47	1.95418	1.95222	64	1.95266	32	2.14745	1.95222	64	2.08642	13
SYNTHES1	0.759284	0.75929	0.75928	72	0.75928	79	0.76028	0.75928	72	0.75928	79	0.85928	0.75928	72	0.79837	17
HS88	1.362657	1.36266	1.36249	189	1.36266	138	1.36402	1.36249	189	1.36387	123	1.49892	1.36249	189	1.36387	123
HS13	0.999872	0.99987	0.99987	116	1.00000	86	1.00087	1.00060	88	1.00085	41	1.09987	1.07956	40	1.07745	22
MATRIX2	0.000000	0.00000	0.00000	131	0.00000	187	0.00100	0.00001	107	0.00000	187	0.10000	0.00001	107	0.09622	46
WOMFLET	0.000000	0.00000	0.00000	78	0.00000	104	0.00100	0.00000	78	0.00000	104	0.10000	0.00000	78	0.00000	104
S365	0.000000	0.00000	0.00000	63	0.00000	179	0.00100	0.00000	63	0.00000	179	0.10000	0.00000	63	0.00000	179

Table A.29. Function evaluations to obtain  $\phi_k \leq \phi^* + TOL \cdot \max\{1.0, |\phi^*|\}$  for small scale noiseless constrained problems.



Figure A.14. Efficiency, Noiseless Case. Log-ratio profiles comparing KNI-TRO with an L-BFGS Hessian approximation of memory one and memory 10 when  $\epsilon(x) = 0$ . The figures measure number of function evaluations to satisfy (2.2.7) for  $\tau = 10^{-1}$  (left),  $10^{-3}$  (middle),  $10^{-6}$  (right).

A.3.3.3. Noisy Functions. In this section, we list the final objective value( $\phi(x)$ ), feasibility error(feaserr) and the number of function evaluations (#feval) needed to achieve this for each problem instance, varying the noise level  $\sigma_f \in \{10^{-1}, 10^{-3}, 10^{-5}, 10^{-7}\}$ , in Tables A.30-A.35. Function evaluations marked with a \* denote cases where the algorithm reached the maximum number of function evaluations. An 'f' marks the cases where the feasibility error is large compared to the noise level, i.e.,  $\|\max\{\psi(x), 0\}\|_{\infty} \geq \sqrt{3\sigma_f}$ .

			KNITRO					COBYLA		
$\sigma_f$	0.0	1e-07	1e-05	1e-03	1e-01	0.0	1.00e-07	1.00e-05	1.00e-03	1.00e-01
CB3	4.92e-12	1.00e-08	2.00e-06	1.35e-04	1.51e-01	8.36e-12	1.00e-08	6.00e-06	1.10e-03	2.97e-01
CHACONN2	1.72e-12	1.00e-08	1.60e-05	3.80e-04	8.27e-02	1.72e-12	1.00e-08	3.00e-06	1.01e-03	f
GIGOMEZ3	9.80e-13	1.00e-08	2.00e-06	1.35e-04	1.62e-01	9.80e-13	1.00e-08	1.00e-05	9.30e-05	f
HS100	2.05e-12	2.37e-06	2.14e-04	1.34e-02	2.16e+00	5.87e-07	1.63e-06	9.62e-04	3.13e-02	3.65e-02
DIPIGRI	4.09e-12	3.24e-05	1.87e-04	1.25e-02	1.98e + 00	2.05e-12	1.63e-06	4.23e-04	3.21e-02	9.14e-01
HS93	3.90e-09	4.34e-03	6.21e-02	2.49e-01	2.69e + 00	3.96e-05	1.52e-01	4.88e-01	1.79e + 00	1.99e+00
HS64	2.28e-09	4.29e-05	3.10e-03	4.90e+00	8.78e + 02	2.28e-09	2.62e-04	7.11e-02	f	f
POLAK3	1.00e-12	6.51e-07	2.13e-04	4.39e-03	f	1.01e-12	5.65e-06	3.58e-04	4.13e-04	f
POLAK1	1.00e-12	1.72e-07	6.17e-06	1.53e-03	2.62e + 00	9.90e-13	1.72e-07	1.17e-06	3.40e-02	5.15e+00
HS104	1.05e-11	5.36e-05	1.30e-02	2.99e-02	6.63e-01	1.05e-11	1.55e-03	1.38e-02	1.37e-01	2.74e+00
HS100MOD	1.02e-12	2.01e-05	2.48e-04	2.06e-02	1.34e+00	1.02e-12	1.11e-07	2.13e-04	8.63e-03	1.20e+00
TWOBARS	1.51e-12	4.18e-07	1.14e-05	1.39e-04	9.27e-02	1.51e-12	4.18e-07	1.54e-04	1.56e-03	5.86e-02
HS85	6.00e-14	3.12e-07	3.12e-07	5.72e-03	9.27e-01	4.51e-06	3.12e-07	6.88e-07	4.82e-01	9.59e-01
CB2	3.70e-13	5.06e-07	4.49e-06	5.47e-04	2.05e-01	3.70e-13	4.94e-07	3.49e-06	8.49e-06	f
CHACONN1	8.80e-13	4.94e-07	6.49e-06	3.96e-03	1.34e-01	1.07e-12	4.94e-07	1.51e-06	6.81e-04	2.44e-02
MADSEN	1.00e-12	5.64e-07	7.56e-06	1.11e-04	2.58e-01	1.01e-12	5.64 e-07	6.44e-06	2.90e-04	2.66e-01
HS66	8.64e-13	2.74e-07	1.37e-05	6.56e-04	2.53e-02	8.64e-13	2.74e-07	4.27e-06	5.23e-04	9.11e-03
MINMAXBD	1.01e-12	3.21e-07	5.63e-05	3.75e-03	f	2.52e-10	3.21e-07	2.03e-05	1.65e-04	$1.31e{+}01$
CANTILVR	4.40e-13	3.61e-07	1.64e-06	1.72e-03	1.91e-01	4.40e-13	6.39e-07	7.16e-05	8.57e-03	f
HS92	1.06e-09	1.16e-02	1.18e-01	f	9.97e-01	1.40e-07	5.02e-03	2.51e-01	2.14e+00	1.07e+00
HALDMADS	9.94e-13	3.71e-07	1.06e-05	7.51e-03	2.53e-01	9.52e-06	3.66e-05	5.66e-05	4.94e-04	f
HS34	2.89e-13	4.45e-07	2.55e-06	4.24e-04	6.93e-01	3.92e-09	4.45e-07	1.24e-05	1.82e-03	8.34e-01
HS90	1.06e-09	1.09e-04	6.99e-01	f	5.73e-01	2.18e-04	7.24e-03	3.13e-01	1.30e+00	1.03e+00
SNAKE	7.96e-13	f	2.28e + 00	4.51e+01	8.35e+00	6.74e-07	2.12e-02	3.24e-02	3.56e-02	1.28e+00
HS72	0.00e+00	6.83e-02	2.53e+01	3.78e + 02	f	0.00e+00	1.63e-02	f	f	f
GIGOMEZ2	0.00e+00	4.94e-07	1.15e-05	3.21e-04	7.78e-02	0.00e+00	4.94e-07	2.95e-05	9.14e-04	1.04e-01
SYNTHES1	0.00e+00	2.39e-06	7.39e-06	6.80e-03	1.22e + 00	4.00e-15	2.61e-06	7.98e-02	8.49e-02	1.10e+00
HS88	1.10e-13	1.29e-04	3.85e-02	4.04e-01	1.32e + 00	4.08e-12	1.76e-04	1.82e-01	7.87e-01	1.13e+00
HS13	2.97e-08	1.08e-02	4.54e-02	1.43e-01	6.42e-01	1.28e-04	6.70e-03	2.95e-02	1.15e-01	1.38e+00
MATRIX2	7.93e-16	1.00e-08	1.10e-05	3.63e-03	7.53e-02	3.82e-17	1.00e-08	3.80e-05	2.52e-03	1.19e+00
WOMFLET	7.89e-13	2.90e-05	1.39e-03	8.39e-03	6.25e + 00	7.89e-13	1.00e-08	2.00e-06	3.12e-03	3.34e + 00
S365	0.00e+00	1.00e-08	1.00e-08	1.00e-08	1.00e-08	8.89e-35	1.00e-08	1.00e-08	1.00e-08	1.00e-08

Table A.30. Optimality gap  $|\phi(x_k) - \phi^*|$  for small scale noisy constrained CUTEst problems.

			KNITRO					COBYLA		
$\sigma_{f}$	0.0	1e-07	1e-05	1e-03	1e-01	0.0	1.00e-07	1.00e-05	1.00e-03	1.00e-01
CB3	2.22e-16	0.00e+00	0.00e+00	7.50e-05	0.00e+00	0.00e+00	0.00e+00	1.20e-05	0.00e+00	3.71e-02
CHACONN2	0.00e+00	0.00e+00	0.00e+00	1.22e-04	2.54e-01	0.00e+00	0.00e+00	5.00e-06	0.00e+00	1.53e+00
GIGOMEZ3	0.00e+00	0.00e+00	0.00e+00	7.50e-05	0.00e+00	0.00e+00	0.00e+00	0.00e+00	3.32e-04	5.89e-01
HS100	0.00e+00	5.00e-06	2.10e-05	6.18e-04	1.01e-01	0.00e+00	0.00e+00	0.00e+00	0.00e+00	1.67e-01
DIPIGRI	2.84e-14	3.10e-05	3.00e-06	5.53e-04	1.38e-01	0.00e+00	0.00e+00	0.00e+00	0.00e+00	0.00e+00
HS93	4.44e-16	0.00e+00	1.00e-05	1.03e-03	1.13e-01	0.00e+00	0.00e+00	0.00e+00	0.00e+00	0.00e+00
HS64	0.00e+00	0.00e+00	1.00e-05	2.18e-03	4.96e-01	0.00e+00	0.00e+00	0.00e+00	1.82e-01	4.97e + 00
POLAK3	3.89e-16	1.00e-06	8.00e-06	1.34e-03	5.56e-01	0.00e+00	0.00e+00	0.00e+00	8.60e-03	8.72e + 00
POLAK1	0.00e+00	0.00e+00	0.00e+00	2.54e-04	6.74e-02	0.00e+00	0.00e+00	1.10e-05	9.15e-04	9.40e-02
HS104	0.00e+00	0.00e+00	0.00e+00	1.04e-03	2.81e-01	0.00e+00	0.00e+00	3.00e-06	2.80e-02	3.70e-01
HS100MOD	0.00e+00	0.00e+00	2.80e-05	3.91e-04	1.77e-01	0.00e+00	0.00e+00	0.00e+00	0.00e+00	7.18e-02
TWOBARS	0.00e+00	0.00e+00	9.00e-06	1.28e-03	1.86e-02	0.00e+00	0.00e+00	0.00e+00	1.04e-03	1.40e-01
HS85	7.11e-15	1.00e-06	7.00e-06	1.52e-04	0.00e+00	0.00e+00	0.00e+00	0.00e+00	1.41e-03	0.00e+00
CB2	1.11e-16	0.00e+00	1.40e-05	0.00e+00	1.15e-02	0.00e+00	0.00e+00	6.00e-06	1.02e-03	1.10e+00
CHACONN1	0.00e+00	1.00e-06	1.40e-05	0.00e+00	2.43e-03	0.00e+00	0.00e+00	2.00e-06	7.34e-04	2.38e-02
MADSEN	0.00e+00	0.00e+00	0.00e+00	4.89e-04	0.00e+00	0.00e+00	0.00e+00	1.00e-05	3.03e-04	0.00e+00
HS66	4.44e-16	0.00e+00	3.00e-06	0.00e+00	0.00e+00	0.00e+00	0.00e+00	5.00e-06	1.21e-03	4.57e-02
MINMAXBD	6.17e-14	2.20e-05	1.90e-05	2.96e-04	4.65e + 02	0.00e+00	0.00e+00	0.00e+00	1.53e-03	0.00e+00
CANTILVR	3.61e-16	0.00e+00	8.00e-06	0.00e+00	5.74e-02	0.00e+00	0.00e+00	0.00e+00	0.00e+00	2.47e+00
HS92	4.60e-17	0.00e+00	2.58e-04	1.33e-01	1.35e-01	0.00e+00	0.00e+00	6.73e-04	1.04e-03	7.11e-02
HALDMADS	4.44e-16	2.00e-06	1.40e-05	0.00e+00	0.00e+00	0.00e+00	0.00e+00	1.50e-05	9.94e-04	1.82e + 01
HS34	4.44e-16	0.00e+00	6.00e-06	0.00e+00	1.15e-01	0.00e+00	0.00e+00	0.00e+00	0.00e+00	0.00e+00
HS90	0.00e+00	0.00e+00	1.20e-05	1.33e-01	1.40e-01	0.00e+00	0.00e+00	1.07e-03	6.30e-05	1.20e-01
SNAKE	1.02e-16	1.70e-03	2.12e-04	2.90e-03	1.87e-01	0.00e+00	1.00e-06	3.00e-06	3.27e-04	3.65e-01
HS72	1.73e-18	0.00e+00	7.31e-04	4.62e-02	5.76e-01	4.77e-18	0.00e+00	7.94e-03	3.47e-01	3.26e + 00
GIGOMEZ2	2.22e-16	2.00e-06	0.00e+00	6.68e-04	1.78e-01	1.11e-16	0.00e+00	0.00e+00	1.29e-03	0.00e+00
SYNTHES1	0.00e+00	1.00e-06	5.00e-06	0.00e+00	2.09e-01	1.54e-23	0.00e+00	0.00e+00	0.00e+00	0.00e+00
HS88	0.00e+00	0.00e+00	4.40e-05	2.05e-03	1.26e-01	2.09e-17	0.00e+00	3.57e-04	1.82e-02	1.47e-01
HS13	2.61e-13	0.00e+00	1.20e-05	4.07e-04	6.52e-02	1.07e-27	0.00e+00	3.00e-06	0.00e+00	0.00e+00
MATRIX2	3.60e-17	0.00e+00	9.00e-06	7.30e-04	2.95e-02	5.20e-18	0.00e+00	6.00e-06	3.17e-04	0.00e+00
WOMFLET	5.89e-10	1.50e-05	1.00e-06	4.53e-04	5.05e-02	8.99e-18	0.00e+00	5.00e-06	0.00e+00	0.00e+00
S365	1.72e-07	0.00e+00	2.00e-06	6.43e-04	6.98e-03	4.35e-17	$0.00\mathrm{e}{+00}$	7.00e-06	1.70e-04	1.72e-02

Table A.31. Feasibility error  $\|\max\{\psi(x_k), 0\}\|_{\infty}$  for small scale noisy constrained CUTEst problems.

			$\tau =$	$10^{-6}$			$\tau = 10^{-2}$					
	F	KNITRO		0	COBYLA		К	NITRO			COBYLA	
problem	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals
CB3	2.000000	0.00e+00	35	2.000000	0.00e+00	30	1.9999980	6.00e-06	30	2.000030	1.20e-05	20
CHACONN2	1.999998	6.00e-06	30	2.000000	0.00e+00	29	1.9999980	6.00e-06	30	1.999631	5.01e-04	22
GIGOMEZ3	2.000000	0.00e+00	35	2.000000	0.00e+00	24	2.0000000	0.00e+00	35	2.000000	0.00e+00	24
HS100	680.630087	8.00e-06	173	680.630055	1.15e-04	125	680.7482460	0.00e+00	75	680.707367	0.00e+00	87
DIPIGRI	680.630025	3.10e-05	189	680.630055	2.80e-05	205	680.7513980	0.00e+00	75	680.638226	2.42e-04	114
HS93	135.080301	1.00e-06	650	135.092558	2.00e-06	f	135.0937190	1.00e-06	477	135.080622	5.37e-04	535
HS64	6299.795253	2.10e-05	77	6300.065790	1.20e-05	312	6688.2748180	5.70e-05	30	6704.980195	1.24e-04	238
POLAK3	5.933006	1.00e-06	357	5.933005	1.19e-04	1333	5.9332460	1.55e-04	268	5.939929	0.00e+00	1220
POLAK1	2.718281	1.00e-06	50	2.718279	9.00e-06	251	2.7182600	5.50e-05	45	2.743571	7.00e-05	216
HS104	3.951218	0.00e+00	305	3.952662	1.00e-06	f	3.9535870	1.20e-05	198	3.953450	4.63e-04	353
HS100MOD	678.679616	5.42e-04	146	678.679644	7.50e-05	179	678.7677270	4.00e-06	110	678.849653	0.00e+00	83
TWOBARS	1.508653	0.00e+00	55	1.508652	0.00e+00	40	1.5086550	3.00e-06	45	1.508930	0.00e+00	21
HS85	-2.215605	0.00e+00	152	-2.215605	9.00e-06	246	-2.2156220	5.41e-04	145	-2.206122	0.00e+00	233
CB2	1.952224	2.00e-06	57	1.952225	1.00e-06	48	1.9522210	4.00e-06	52	1.960966	0.00e+00	22
CHACONN1	1.952225	1.00e-06	40	1.952225	1.00e-06	37	1.9521680	8.20e-05	35	1.955056	0.00e+00	20
MADSEN	0.616433	0.00e+00	58	0.616432	4.74e-04	32	0.6164330	0.00e+00	58	0.616264	5.11e-04	31
HS66	0.518164	0.00e+00	33	0.518163	0.00e+00	37	0.5186740	0.00e+00	21	0.518026	3.71e-04	26
MINMAXBD	115.706388	1.17e-04	485	115.706134	4.03e-04	783	115.7063880	1.17e-04	485	115.919506	0.00e+00	741
CANTILVR	1.339957	0.00e+00	211	1.339958	3.80e-05	131	1.3415130	0.00e+00	187	1.347590	0.00e+00	107
HS92	1.374297	0.00e+00	f	1.367672	0.00e+00	139	1.3742970	0.00e+00	f	1.366663	2.00e-06	117
HALDMADS	0.000121	6.00e-06	78	0.000159	1.00e-06	f	0.0001210	6.00e-06	78	0.000159	1.00e-06	f
HS34	-0.834033	1.00e-06	54	-0.834032	0.00e+00	34	-0.8320120	0.00e+00	44	-0.834016	0.00e+00	25
HS90	1.362766	0.00e+00	326	1.368155	1.00e-06	f	1.3608930	2.00e-06	206	1.360847	1.10e-05	156
SNAKE	-22.242365	1.70e-03	f	-0.021173	2.00e-06	f	-22.2423650	1.70e-03	f	-0.021173	2.00e-06	f
HS72	727.747619	0.00e+00	f	727.695652	0.00e+00	870	721.1362080	3.33e-04	66	720.999346	3.14e-04	815
GIGOMEZ2	1.952224	1.00e-06	64	1.952225	1.80e-05	46	1.9520790	2.15e-04	59	1.952432	0.00e+00	31
SYNTHES1	0.759272	1.60e-05	64	0.759285	1.00e-06	53	0.7591430	1.95e-04	56	0.785862	1.00e-06	30
HS88	1.362528	0.00e+00	331	1.362528	0.00e+00	114	1.3581270	5.00e-06	240	1.354190	8.00e-06	93
HS13	0.989068	0.00e+00	97	1.006571	0.00e+00	f	1.1183300	0.00e+00	36	1.138030	0.00e+00	19
MATRIX2	0.000000	0.00e+00	153	0.000001	0.00e+00	113	0.0000000	0.00e+00	153	0.000001	0.00e+00	113
WOMFLET	0.000029	1.50e-05	f	0.000006	6.00e-05	121	0.0230460	4.21e-04	81	0.000468	3.02e-04	68
S365	0.000000	5.41e-04	54	0.000000	3.13e-04	48	0.0000000	5.41e-04	54	0.000000	3.13e-04	48

Table A.32. Number of function evaluations to obtain (2.4.5) for noise level  $\sigma_f = 10^{-7}$ .

			au = 1	$10^{-6}$			$\tau = 10^{-2}$					
	ŀŀ	KNITRO		C	OBYLA		1	KNITRO		0	OBYLA	
problem	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals
CB3	2.000003	0.00e+00	47	2.000000	8.00e-06	32	1.999981	2.50e-05	30	2.000031	1.10e-05	20
CHACONN2	2.000006	4.00e-06	45	2.000001	2.79e-04	31	1.999999	1.90e-05	30	1.999624	5.08e-04	22
GIGOMEZ3	2.000003	0.00e + 00	47	2.000003	1.26e-03	25	2.000003	0.00e+00	47	2.000003	1.26e-03	25
HS100	680.630284	9.00e-06	160	680.630275	1.47e-04	152	680.849165	0.00e+00	46	680.757842	1.82e-03	56
DIPIGRI	680.630221	7.00e-06	149	680.630243	5.00e-05	164	680.817892	1.54e-03	56	680.718994	3.36e-04	58
HS93	135.138021	1.00e-05	229	135.416906	1.60e-05	f	135.133971	1.16e-03	66	135.416906	1.60e-05	f
HS64	6299.871889	0.00e + 00	66	6299.950263	2.61e-03	308	6708.036271	2.43e-03	25	6669.894009	2.09e-03	246
POLAK3	5.933216	8.00e-06	313	5.933217	7.40e-05	2303	5.932429	3.09e-03	270	5.943411	4.15e-03	2186
POLAK1	2.718282	6.00e-06	50	2.718280	1.20e-05	269	2.718182	9.00e-04	45	2.743588	4.30e-05	230
HS104	3.964210	0.00e + 00	170	3.963888	1.04e-04	f	3.964996	3.71e-04	109	3.965500	9.61e-04	113
HS100MOD	678.679887	1.30e-05	193	678.679885	5.49e-04	122	678.947334	1.28e-03	73	678.953589	0.00e+00	52
TWOBARS	1.508641	9.00e-06	155	1.508665	1.90e-05	f	1.508463	1.28e-04	41	1.508969	0.00e+00	21
HS85	-2.215605	1.70e-05	106	-2.215605	3.30e-05	243	-2.215628	6.85e-04	99	-2.206170	8.10e-05	235
CB2	1.952220	1.40e-05	60	1.952219	7.00e-06	64	1.957493	6.23e-04	44	1.960991	0.00e+00	22
CHACONN1	1.952218	1.40e-05	58	1.952220	1.40e-05	42	1.952172	6.10e-05	35	1.955070	0.00e+00	20
MADSEN	0.616425	9.00e-06	146	0.616425	1.17e-04	40	0.616422	2.61e-04	49	0.613403	3.27e-03	20
HS66	0.518177	3.00e-06	f	0.518158	1.00e-05	37	0.518164	3.50e-05	32	0.518055	1.87e-03	19
MINMAXBD	115.706237	5.45e-04	926	115.705937	1.72e-03	779	115.705316	2.33e-03	919	122.679517	0.00e+00	719
CANTILVR	1.339958	8.00e-06	149	1.339959	3.30e-05	151	1.338856	2.67e-03	91	1.342286	4.59e-03	111
HS92	1.244708	2.58e-04	f	1.112153	6.73e-04	140	1.244708	2.58e-04	f	1.109888	7.50e-04	68
HALDMADS	0.000133	1.40e-05	101	0.000179	1.40e-05	f	0.000133	1.40e-05	101	0.000135	5.10e-05	77
HS34	-0.834035	6.00e-06	52	-0.834035	3.20e-05	32	-0.833901	0.00e+00	36	-0.834014	0.00e+00	25
HS90	2.061544	1.20e-05	f	1.049278	1.07e-03	118	2.061544	1.20e-05	f	1.049423	1.07e-03	72
SNAKE	-2.275095	2.12e-04	87	-0.032419	3.00e-06	f	-2.269176	2.34e-04	69	-0.032419	3.00e-06	f
HS72	702.418590	7.31e-04	f	528.392603	7.94e-03	f	702.418590	7.31e-04	f	528.392603	7.94e-03	f
GIGOMeZ2	1.952237	0.00e + 00	81	1.952209	1.80e-05	f	1.952214	1.60e-03	49	1.952439	0.00e + 00	31
SYNTHES1	0.759275	6.00e-06	122	0.759278	2.40e-05	113	0.758501	1.34e-04	56	0.839094	0.00e+00	29
HS88	1.324205	4.40e-05	f	1.180569	3.57e-04	67	1.324205	4.40e-05	f	1.177325	3.68e-04	47
HS13	0.954433	1.20e-05	331	1.029338	5.00e-06	f	1.089502	0.00e+00	36	1.138525	0.00e+00	19
MATRIX2	0.000012	9.00e-06	135	0.000038	6.00e-06	f	0.000012	9.00e-06	135	0.000038	6.00e-06	f
WOMFLET	0.001388	1.00e-06	f	-0.000003	2.33e-03	80	0.017147	1.63e-03	90	0.001884	1.73e-04	68
S365	0.000000	6.32e-04	54	0.000000	5.30e-03	36	0.000000	6.32e-04	54	0.000000	5.30e-03	36

Table A.33. Number of function evaluations to obtain (2.4.5) for noise level  $\sigma_f = 10^{-5}$ 

			$\tau =$	$10^{-6}$			$\tau = 10^{-2}$					
	ŀ	KNITRO		0	COBYLA		ł	KNITRO		0	OBYLA	
problem	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals
CB3	2.000135	7.50e-05	57	2.001105	0.00e+00	f	1.991529	1.86e-02	25	1.999922	9.37e-04	17
CHACONN2	2.000382	1.20e-04	30	2.000360	1.84e-03	f	1.989816	2.59e-02	25	2.007665	3.71e-02	19
GIGOMEZ3	2.000135	7.50e-05	f	1.999908	3.32e-04	35	2.000135	7.50e-05	f	1.999908	3.32e-04	35
HS100	680.643454	6.18e-04	118	680.653293	2.80e-03	f	680.862278	2.56e-03	47	680.751163	1.19e-02	46
DIPIGRI	680.642602	5.53e-04	146	680.653238	2.48e-03	f	680.858601	0.00e + 00	47	680.938844	1.37e-02	46
HS93	135.324499	1.03e-03	126	136.540776	7.65e-03	f	135.323447	3.39e-03	35	136.540776	7.65e-03	f
HS64	6294.941500	2.18e-03	74	6307.172876	1.82e-01	f	6746.496044	5.01e-02	20	6307.172876	1.82e-01	f
POLAK3	5.937398	1.34e-03	f	5.932593	8.61e-03	1645	5.954694	7.66e-03	252	5.909504	3.84e-02	1510
POLAK1	2.719811	2.51e-04	86	2.752280	9.16e-04	f	2.723341	7.78e-04	36	2.736065	3.33e-02	231
HS104	3.981084	1.04e-03	174	4.087923	2.80e-02	f	3.981888	9.32e-04	154	4.087923	2.80e-02	f
HS100MOD	678.700188	3.91e-04	f	678.688238	0.00e+00	99	678.964451	3.27e-03	73	678.817053	2.85e-02	61
TWOBARS	1.508513	1.28e-03	f	1.507091	1.04e-03	27	1.508513	1.28e-03	f	1.507577	1.34e-02	21
HS85	-2.209889	1.52e-04	271	-1.733123	1.41e-03	f	-2.211306	1.94e-02	127	-1.733123	1.41e-03	f
CB2	1.952771	0.00e+00	f	1.952216	1.02e-03	45	1.943877	1.91e-02	51	1.956972	0.00e+00	23
CHACONN1	1.956186	0.00e+00	f	1.952905	7.34e-04	23	1.949855	4.28e-02	46	1.934334	2.62e-02	15
MADSEN	0.616321	4.89e-04	f	0.616142	3.03e-04	35	0.619826	0.00e+00	68	0.612382	1.21e-02	26
HS66	0.518819	0.00e + 00	f	0.517640	1.21e-03	48	0.518819	0.00e+00	f	0.518165	1.35e-02	13
MINMAXBD	115.710191	2.96e-04	f	115.706307	4.03e-02	756	115.703385	1.16e-02	538	121.478546	2.68e-02	717
CANTILVR	1.341675	0.00e + 00	141	1.347345	3.04e-03	f	1.343901	5.05e-03	84	1.336533	4.03e-02	145
HS92	0.003678	1.33e-01	f	0.507997	2.58e-02	f	0.003678	1.33e-01	f	0.507997	2.58e-02	f
HALDMADS	0.007637	0.00e + 00	f	0.000616	9.93e-04	71	0.007637	0.00e+00	f	0.000616	9.93e-04	71
HS34	-0.833608	0.00e + 00	71	-0.834809	1.73e-03	f	-0.832197	0.00e + 00	36	-0.832210	0.00e + 00	22
HS90	0.001271	1.33e-01	f	0.683684	1.11e-02	f	0.001271	1.33e-01	f	0.683684	1.11e-02	f
SNAKE	-45.074608	2.90e-03	130	-0.035633	3.27e-04	f	-44.926682	3.22e-02	60	-0.035633	3.27e-04	f
HS72	349.372500	4.62e-02	f	66.439486	3.47e-01	f	349.372500	4.62e-02	f	66.439486	3.47e-01	f
GIGOMeZ2	1.951903	6.68e-04	f	1.951310	1.29e-03	48	1.951903	6.68e-04	f	1.951561	7.74e-04	36
SYNTHES1	0.766084	0.00e + 00	76	0.755256	1.12e-03	f	0.790862	1.53e-03	56	0.844177	0.00e + 00	29
HS88	0.958819	2.05e-03	f	0.575949	1.82e-02	45	0.958819	2.05e-03	f	0.576448	1.81e-02	26
HS13	0.857313	4.07e-04	48	1.114815	0.00e + 00	f	0.857380	4.06e-04	38	1.114815	0.00e + 00	f
MATRIX2	0.003633	7.30e-04	f	0.002521	3.17e-04	68	0.003633	7.30e-04	f	0.002521	3.17e-04	68
WOMFLET	0.008386	4.53e-04	f	0.003122	0.00e + 00	108	0.007706	5.15e-02	69	-0.051536	5.38e-02	87
S365	0.000000	4.03e-02	45	0.000000	5.11e-02	32	0.000000	4.03e-02	45	0.000016	4.64 e- 02	25

Table A.34. Number of function evaluations to obtain (2.4.5) for noise level  $\sigma_f = 10^{-3}$ .

			$\tau =$	$10^{-6}$			$\tau = 10^{-2}$						
	l ł	KNITRO		C	OBYLA		l ł	KNITRO		C	OBYLA		
problem	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	$\phi(x)$	feaserr	#fevals	
CB3	2.151366	0.00e+00	200	2.297095	3.71e-02	f	2.151799	0.00e+00	183	2.297095	3.71e-02	f	
CHACONN2	1.917274	2.54e-01	f	0.544480	1.53e + 00	f	1.917274	2.54e-01	f	0.544480	1.53e + 00	f	
GIGOMEZ3	2.162155	0.00e + 00	f	1.558092	5.89e-01	f	2.162155	0.00e+00	f	1.554036	5.32e-01	20	
HS100	682.785696	1.01e-01	f	680.545842	2.09e-01	f	682.785696	1.01e-01	f	680.679716	3.84e-01	25	
DIPIGRI	682.612679	1.38e-01	f	681.544064	0.00e+00	49	682.612679	1.38e-01	f	681.282921	7.16e-02	39	
HS93	132.388506	1.13e-01	126	135.053154	4.58e-02	f	132.388506	1.13e-01	126	135.053154	4.58e-02	f	
HS64	7178.300921	4.96e-01	33	13921.561135	4.97e + 00	f	6221.151364	5.05e-01	27	13921.561135	4.97e + 00	f	
POLAK3	5.768461	5.56e-01	f	-2.341289	8.72e + 00	f	5.768461	5.56e-01	f	-2.341289	8.72e + 00	f	
POLAK1	5.342043	6.74e-02	86	7.870717	9.39e-02	f	5.381476	1.35e-01	62	7.870717	9.39e-02	f	
HS104	3.287811	2.81e-01	f	1.208522	3.70e-01	30	3.287811	2.81e-01	f	1.193347	3.97e-01	10	
HS100MOD	680.016339	1.77e-01	f	679.878287	7.18e-02	45	680.111489	8.61e-02	53	679.599540	2.77e-01	34	
TWOBARS	1.601395	1.86e-02	f	1.450060	1.40e-01	30	1.601395	1.86e-02	f	1.450155	1.40e-01	12	
HS85	-1.288983	0.00e+00	30	-1.256810	0.00e+00	f	-1.288983	0.00e+00	30	-1.256810	0.00e+00	f	
CB2	2.157319	1.15e-02	f	0.871901	1.10e+00	f	2.157319	1.15e-02	f	0.871901	1.10e+00	f	
CHACONN1	2.085730	2.43e-03	f	1.976579	2.38e-02	27	2.085730	2.43e-03	f	1.979692	2.45e-02	12	
MADSEN	0.874476	0.00e+00	70	0.882487	0.00e+00	f	0.874516	0.00e+00	33	0.882487	0.00e+00	f	
HS66	0.543480	0.00e+00	f	0.509050	4.56e-02	37	0.543480	0.00e+00	f	0.508822	4.63e-02	20	
MINMAXBD	-33.388429	4.65e + 02	f	112.975518	2.81e + 00	f	-33.388429	4.65e + 02	f	112.975518	$2.81e{+}00$	f	
CANTILVR	1.531159	5.74e-02	f	0.934009	2.47e + 00	f	1.531159	5.74e-02	f	0.934009	2.47e + 00	f	
HS92	0.366123	1.35e-01	f	0.293398	7.11e-02	48	0.366123	1.35e-01	f	0.297797	7.09e-02	35	
HALDMADS	0.252628	0.00e+00	f	-3.027998	1.82e + 01	f	0.252628	0.00e+00	f	-3.027998	1.82e + 01	f	
HS34	-0.141457	1.15e-01	49	-0.750941	2.99e-01	f	-0.140985	1.15e-01	25	-0.750941	2.99e-01	f	
HS90	0.790027	1.40e-01	f	0.331080	1.20e-01	48	0.790027	1.40e-01	f	0.331906	1.20e-01	13	
SNAKE	-8.350669	1.87e-01	93	1.284550	3.65e-01	f	-8.350669	1.87e-01	93	1.284550	3.65e-01	f	
HS72	67.968573	5.76e-01	f	8.946566	3.26e + 00	f	67.968573	5.76e-01	f	8.946566	3.26e + 00	f	
GIGOMeZ2	1.874401	1.78e-01	78	2.056211	0.00e+00	f	1.874484	1.90e-01	70	2.056211	0.00e+00	f	
SYNTHES1	1.977343	2.09e-01	f	1.855036	0.00e+00	11	1.841763	3.65e-01	69	1.855036	0.00e+00	11	
HS88	0.047086	1.26e-01	16	0.235414	1.47e-01	f	0.047086	1.26e-01	16	0.235414	1.47e-01	f	
HS13	0.357703	6.52e-02	48	2.381364	0.00e+00	f	0.551282	1.71e-02	19	2.381364	0.00e+00	f	
MATRIX2	0.075268	2.95e-02	194	1.186858	0.00e+00	f	0.075268	2.95e-02	194	1.186858	0.00e+00	f	
WOMFLET	6.254269	5.05e-02	f	3.338088	0.00e+00	48	6.254269	5.05e-02	f	3.336520	0.00e+00	35	
S365	0.000000	9.22e-02	45	0.000000	5.72e-02	24	0.000000	9.22e-02	45	-0.001477	1.53e-01	23	

Table A.35. Number of function evaluations to obtain (2.4.5) for noise level  $\sigma_f = 10^{-1}$ .

#### APPENDIX B

## Adaptive Finite-Difference Interval Estimation for Noisy Derivative-Free Optimization

#### **B.1.** Finite-Difference Formula Derivation and Tables

We summarize the different standard finite-difference schemes with equidistant points, their theoretical error, optimal steplength, and optimal error in terms of the noise level  $\epsilon_f$ and local bound on the *q*-th derivative  $L_q$  for a smooth univariate function  $\phi : \mathbb{R} \to \mathbb{R}$  in Tables B.1 and B.2. For completeness, we provide a complete derivation of the errors for a generic finite-difference approximation to the *d*-th order derivative below.

We will use  $f : \mathbb{R} \to \mathbb{R}$  to denote the noisy function evaluations  $f(t) = \phi(t) + \epsilon(t)$ . We will consider two settings for  $\epsilon(t)$ : (1) we will assume that  $\epsilon(t)$  is bounded, i.e.,  $|\epsilon(t)| \leq \epsilon_f$ for all t; (2) we will assume that  $\epsilon(t)$  is a random variable with  $\mathbb{E}[\epsilon(t)] = 0$  and  $\mathbb{E}[\epsilon(t)^2] = \sigma_f^2$ for all t. The tables vary the number of evaluated points m and is dependent on the local Lipschitz constant  $L_q \geq 0$  which bounds the q-th derivative

$$|\phi^{(q)}(t+h_0s)| \le L_q$$

for all  $s \in [s_1, s_m]$ , where q is the order of the remainder term in the Taylor expansion.

In the most general case, given distinct shifts  $\{s_j\}_{j=1}^m$  and points  $\{t_1, ..., t_m\} = \{t + hs_1, t + hs_2, ..., t + hs_m\}$ , one can derive a generic finite-difference method to approximate

the d-th derivative of the form:

$$\phi^{(d)}(t) \approx \frac{\sum_{j=1}^{m} w_j f(t+s_j h)}{h^d} = f^{(d)}(t;h).$$

We will assume without loss of generality that  $s_1 < s_2 < ... < s_m$ . First, note that  $f^{(d)}$  can be decomposed into a noiseless finite-difference formula and its corresponding error:

$$f^{(d)}(t;h) = \frac{\sum_{j=1}^{m} w_j \phi(t+s_j h)}{h^d} + \frac{\sum_{j=1}^{m} w_j \epsilon(t+s_j h)}{h^d}$$

Considering the noiseless finite-difference term, since the function is smooth, one can write the Lagrange remainder form of the Taylor series expansions for each function evaluation without noise as:

$$\phi(t+hs_j) = \sum_{l=0}^{q-1} \frac{1}{l!} \phi^{(l)}(t) s_j^l + \frac{1}{q!} \phi^{(q)}(\xi_j) s_j^q$$

for  $\xi_j \in [t, t + hs_j]$  for j = 1, ..., m. Therefore, if the weights w satisfy

$$\frac{1}{d!} \sum_{j=1}^{m} w_j s_j^l = \begin{cases} 0 & \text{for } l \neq d, \ l = 0, 1, ..., q-1 \\ 1 & \text{for } l = d \end{cases}$$

then

$$\frac{\sum_{j=1}^{m} w_j \phi(t+s_j h)}{h^d} = \phi^{(d)}(t) + \frac{h^{q-d}}{q!} \sum_{j=1}^{m} w_j \phi^{(q)}(\xi_j) s_j^q.$$

This can be written compactly by the linear system of equations:

$$V(s)^T w = d! \cdot e_{p-d}$$

where  $V(s) \in \mathbb{R}^{m \times q}$  is the Vandermonde matrix defined as

$$V(s) = \begin{bmatrix} s_1^{q-1} & s_1^{q-2} & \dots & s_1^0 \\ s_2^{q-1} & s_2^{q-2} & \dots & s_2^0 \\ \vdots & \vdots & \ddots & \vdots \\ s_m^{q-1} & s_m^{q-2} & \dots & s_m^0 \end{bmatrix}$$

and  $e_{p-d} \in \mathbb{R}^p$  is the (p-d)-th coordinate vector.

To derive a reasonable bound on the total error, suppose we are given  $h_0 > 0$  and a bound on  $\phi^{(q)}$ 

$$|\phi^{(q)}(t+sh_0)| \le L_q$$

for all  $s \in [s_1, s_m]$ . If we assume that the error is bounded, i.e.,  $|\epsilon(t)| \leq \epsilon_f$ , then one can then bound the error in the approximation by:

$$|f^{(d)}(t;h) - \phi^{(d)}(t)| \le \frac{L_q h^{q-d}}{q!} \sum_{j=1}^m |w_j s_j^q| + \frac{\|w\|_1 \epsilon_f}{h^d} = \epsilon_g(h)$$

for all  $0 < h \le h_0$ . If we assume instead that  $Var(\epsilon(t)) = \sigma_f^2$ , then we can similarly show

$$\mathbb{E}[(f^{(d)}(t;h) - \phi^{(d)}(t))^2] \le \frac{L_q^2 h^{2(q-d)}}{(q!)^2} \sum_{j=1}^m w_j^2 s_j^{2q} + \frac{\|w\|_2^2 \sigma_f^2}{h^{2d}} = \sigma_g^2(h)$$

for all  $0 < h \le h_0$ .

The above Taylor series analysis is pessimistic in that it requires multiple  $\xi_j$  points, and therefore yields a loose bound when applying the triangle inequality. Instead, one can consider the derivation of finite-difference schemes for approximating the *first* derivative at an interpolation point using Lagrange polynomials, which yields a tighter bound on the error.

As above, suppose we are given distinct points  $\{t_1, ..., t_m\} = \{t + hs_1, ..., t + hs_m\}$  and we are interested in approximating  $\phi^{(1)}(t)$ . Recall that the Lagrange basis polynomials are defined as:

$$\psi_{p,j}(\tilde{t}) = \frac{\prod_{k \neq j} (\tilde{t} - t_k)}{\prod_{k \neq j} (t_j - t_k)} = \frac{\omega_m(\tilde{t})}{\omega_m^{(1)}(t_j)(\tilde{t} - t_j)}, \ \omega_m(\tilde{t}) = \prod_{j=1}^m (\tilde{t} - t_j).$$

Then the Lagrange interpolation is defined as:

$$\ell(\tilde{t}) = \sum_{j=1}^{m} \psi_{m,j}(\tilde{t})\phi(t_j).$$

It is well-known that the remainder is

$$\phi(\tilde{t}) - \ell(\tilde{t}) = \frac{\omega_m(\tilde{t})}{m!} \phi^{(m)}(\xi)$$

for some  $\xi \in [t_1, t_m]$ . Note that the finite-difference formula can simply be obtained by differentiating the Lagrange polynomial

$$\ell^{(1)}(\tilde{t}) = \sum_{j=1}^{m} \psi_{m,j}^{(1)}(\tilde{t})\phi(t_j).$$

Therefore, the finite-difference coefficients are obtained by evaluating  $\psi_{m,j}^{(1)}(\tilde{t})$ . The error is also obtained by noting

$$\phi^{(1)}(\tilde{t}) = \ell^{(1)}(\tilde{t}) + \frac{\omega_m^{(1)}(\tilde{t})}{m!}\phi^{(m)}(\xi) + \frac{\omega_m(\tilde{t})}{m!}\phi^{(m)}(\xi)\frac{d\xi}{dx}.$$

Since

$$\omega_m^{(1)}(\tilde{t}) = \sum_{j=1}^m \prod_{k \neq j} (\tilde{t} - t_k),$$

plugging in  $\tilde{t} = t_i$  for any i = 1, ..., m, we get the following equality

$$\phi^{(1)}(t_i) = \ell^{(1)}(t_i) + \frac{\omega_m^{(1)}(t_i)}{m!}\phi^{(m)}(\xi) = \ell^{(1)}(t_i) + \prod_{j \neq i} (t_i - t_j)\frac{\phi^{(m)}(\xi)}{m!}.$$

Given  $h_0 > 0$  and a bound on  $\phi^{(m)}$ 

$$|\phi^{(m)}(t+h_0s)| \le L_m$$

for all  $s \in [s_1, s_m]$  and assuming  $t = t_i$  is one of the interpolation points, we obtain the bound

$$|\phi^{(1)}(t) - \ell^{(1)}(t)| \le \frac{L_m h^{m-1}}{m!} \left| \prod_{j \ne i} s_j \right|$$

and if we incorporate the error in the function evaluations, we obtain a error and variance bounds of

$$|f^{(1)}(t;h) - \phi^{(1)}(t)| \le \frac{L_m h^{m-1}}{m!} \left| \prod_{j \ne i} s_j \right| + \frac{\|w\|_1 \epsilon_f}{h} = \epsilon_g(h)$$
$$\mathbb{E}[(f^{(1)}(t;h) - \phi^{(1)}(t))^2] \le \frac{L_m^2 h^{2(m-1)}}{(m!)^2} \prod_{j \ne i} s_j^2 + \frac{\|w\|_2^2 \sigma_f^2}{h^2} = \sigma_g^2(h)$$

for all  $0 < h \leq h_0$ .

Table B.1. Table containing the finite-difference formula, deterministic error bound  $|f^{(1)}(t;h) - \phi^{(1)}(t)| \le \epsilon_g(h)$  for generic h, optimal  $h^*$ , and optimal error  $\epsilon_g(h^*)$  for forward-difference schemes with number of points  $m \in \{2, 3, 4, 5\}$ .

m	$f^{(1)}(t;h)$	$\epsilon_g(h)$	$h^*$	$\epsilon_g(h^*)$
2	$rac{f(t+h)-f(t)}{h}$	$\frac{L_2h}{2} + \frac{2\epsilon_f}{h}$	$2\sqrt{\frac{\epsilon_f}{L_2}}$	$2\sqrt{L_2\epsilon_f}$
3	$\frac{-3f(t)+4f(t+h)-f(t+2h)}{2h}$	$\frac{L_3h^2}{3} + \frac{4\epsilon_f}{h}$	$\sqrt[3]{\frac{6\epsilon_f}{L_3}}$	$6^{2/3}L_3^{1/3}\epsilon_f^{2/3}$
4	$\frac{-11f(t)+18f(t+h)-9f(t+2h)+2f(t+3h)}{6h}$	$\frac{L_4h^3}{4} + \frac{20\epsilon_f}{3h}$	$\sqrt[4]{\frac{80\epsilon_f}{9L_4}}$	$\frac{8\cdot 5^{3/4}}{3\sqrt{3}}L_4^{1/4}\epsilon_f^{3/4}$
5	$\frac{-25f(t)+48f(t+h)-36f(t+2h)+16f(t+3h)-3f(t+4h)}{12h}$	$\frac{L_5h^4}{5} + \frac{32\epsilon_f}{3h}$	$\sqrt[5]{\frac{40\epsilon_f}{3L_5}}$	$4\left(\frac{5}{3}\right)^{4/5} 2^{2/5} L_5^{1/5} \epsilon_f^{4/5}$

Table B.2. Table containing the finite-difference formula, deterministic error bound  $|f^{(1)}(t;h) - \phi^{(1)}(t)| \le \epsilon_g(h)$  for generic h, optimal  $h^*$ , and optimal error  $\epsilon_g(h^*)$  for central-difference schemes with number of points  $m \in \{2, 4, 6\}$ .

$\overline{m}$	$f^{(1)}(t;h)$	$\epsilon_g(h)$	$h^*$	$\epsilon_g(h^*)$
2	$rac{f(t+h)-f(t-h)}{2h}$	$\frac{L_3h^2}{6} + \frac{\epsilon_f}{h}$	$\sqrt[3]{\frac{3\epsilon_f}{L_3}}$	$\frac{3^{2/3}}{2}L_3^{1/3}\epsilon_f^{2/3}$
4	$\frac{f(t-2h) - 8f(t-h) + 8f(t+h) - f(t+2h)}{12h}$	$\frac{L_5h^4}{30} + \frac{3\epsilon_f}{2h}$	$\sqrt[5]{\frac{45\epsilon_f}{4L_5}}$	$\frac{1}{4} \left(\frac{3}{2}\right)^{4/5} 5^{4/5} L_5^{1/5} \epsilon_f^{4/5}$
6	$\frac{-f(t-3h)+9f(t-2h)-45f(t-h)+45f(t+h)-9f(t+2h)+f(t+3h)}{60h}$	$\frac{L_7h^6}{140} + \frac{11\epsilon_f}{6h}$	$\sqrt[7]{\frac{385\epsilon_f}{9L_7}}$	$\frac{77^{6/7}}{12\cdot 3^{5/7}\cdot \sqrt[7]{5}}L_7^{1/7}\epsilon_f^{6/7}$

Table B.3. Table containing the finite-difference formula, MSE error bound  $\mathbb{E}[(f^{(1)}(t;h) - \phi^{(1)}(t))^2] \leq \sigma_g^2(h)$  for generic h, optimal  $h^*$ , and optimal error  $\sigma_g(h^*)$  for forward-difference schemes with number of points  $m \in \{2, 3, 4, 5\}$ .

p	$f^{(1)}(t;h)$	$\sigma_g^2(h)$	$h^*$	$\sigma_g(h^*)$
1	$rac{f(t+h)-f(t)}{h}$	$\frac{L_2^2h^2}{4} + \frac{2\epsilon_f^2}{h^2}$	$8^{1/4}\sqrt{\frac{\epsilon_f}{L_2}}$	$2^{1/4}\sqrt{L_2\epsilon_f}$
2	$\frac{-3f(t)+4f(t+h)-f(t+2h)}{2h}$	$\frac{L_3^2 h^4}{9} + \frac{13\epsilon_f^2}{2h^2}$	$\left(\frac{3}{2}\right)^{1/3} 13^{1/6} \sqrt[3]{\frac{\epsilon_f}{L_3}}$	$\frac{\sqrt[6]{3}\sqrt[3]{13}}{2^{2/3}}L_3^{1/3}\epsilon_f^{2/3}$
3	$\frac{-11f(t)\!+\!18f(t\!+\!h)\!-\!9f(t\!+\!2h)\!+\!2f(t\!+\!3h)}{6h}$	$\frac{L_4^2 h^6}{16} + \frac{265 \epsilon_f^2}{18 h^2}$	$\left(\frac{2}{3}\right)^{3/8} 265^{1/8} \sqrt[4]{\frac{\epsilon_f}{L_4}}$	$\frac{1}{3}\sqrt[8]{\frac{2}{3}}265^{3/8}L_4^{1/4}\epsilon_f^{3/4}$
4	$\frac{-25f(t)\!+\!48f(t\!+\!h)\!-\!36f(t\!+\!2h)\!+\!16f(t\!+\!3h)\!-\!3f(t\!+\!4h)}{12h}$	$\frac{L_5^2 h^8}{25} + \frac{2245\epsilon_f^2}{72h^2}$	$\frac{5^{3/10}449^{1/10}}{\sqrt{25/3}} \sqrt[5]{\frac{\epsilon_f}{L_5}}$	$\frac{5^{7/10}449^{2/5}}{4\cdot 3^{4/5}}L_5^{1/5}\epsilon_f^{4/5}$

Table B.4. Table containing the finite-difference formula, MSE error bound  $\mathbb{E}[(f^{(1)}(t;h) - \phi^{(1)}(t))^2] \leq \sigma_g^2(h)$  for generic h, optimal  $h^*$ , and optimal error  $\sigma_g(h^*)$  for central-difference schemes with number of points  $m \in \{2, 4, 6\}$ .

p	$f^{(1)}(t;h)$	$\sigma_g^2(h)$	$h^*$	$\sigma_g(h^*)$
2	$rac{f(t+h)-f(t-h)}{2h}$	$\frac{L_3^2 h^4}{36} + \frac{\epsilon_f^2}{2h^2}$	$\sqrt[3]{3} \sqrt[3]{\frac{\epsilon_f}{L_3}}$	$\frac{\sqrt[6]{3}}{2}L_3^{1/3}\epsilon_f^{2/3}$
4	$\frac{f(t-2h) - 8f(t-h) + 8f(t+h) - f(t+2h)}{12h}$	$\frac{L_5^2 h^8}{900} + \frac{65\epsilon_f^2}{72h^2}$	$\left(\frac{5}{2}\right)^{3/10} 13^{1/10} \sqrt[5]{\frac{\epsilon_f}{L_5}}$	$\frac{5^{7/10} \cdot 13^{2/5}}{12 \cdot \sqrt[5]{2}} L_5^{1/5} \epsilon_f^{4/5}$
6	$\frac{-f(t-3h)+9f(t-2h)-45f(t-h)+45f(t+h)-9f(t+2h)+f(t+3h)}{60h}$	$\tfrac{L_7^2 h^{12}}{140^2} + \tfrac{2107 \epsilon_f^2}{1800 h^2}$	$\frac{7^{2/7}43^{1/14}}{3^{3/14}}\sqrt[7]{\frac{\epsilon_f}{L_7}}$	$\frac{7^{17/14}43^{3/7}}{60\cdot 3^{2/7}}L_7^{1/7}\epsilon_f^{6/7}$

#### **B.2.** Complete Experimental Results

Here, we present the complete experimental results from Section 3.4.

#### **B.2.1.** Robustness to Different Noise Levels

We test our procedure on a simple function  $\phi(t) = \cos(t)$  for different noise levels using different schemes listed in Table 3.1. These are shown in Figure B.1. Detailed numerical results, including the number of iterations and relative error, are listed in Table B.5.

Observe that our method is able to consistently achieve low relative error using a similar number of function evaluations across all tested noise levels. This is a desirable property, as it demonstrates that our initial choice of the interval h and our method is consistent over different noise levels.

#### **B.2.2.** Affine Invariance

One advantage of our proposed method is that the testing ratio remains unchanged under affine transformations of the function. It is particularly obvious that our procedure is invariant when adding a constant to the function. Hence, we focus on transformations of the form  $\phi(t) \rightarrow a \cdot \phi(b \cdot t)$  for some  $a, b \neq 0$ .

To do this, we test Algorithm 3.2 on the function  $\phi(t) = a \cdot \sin(b \cdot t)$  at t = 0 for various a and b. We fix the noise level to be  $\epsilon_f = 10^{-3}$ . The results are shown in Figure B.2. Detailed results can be found in Table B.6 and B.7. As seen in Figure B.2, our method is affine-invariant and can output consistently correct results for different a and b.

Table B.5. Detailed results for  $\phi(t) = \cos(t)$  with different noise levels; r represents the final testing ratio;  $h^*$  is the h that minimizes  $\delta_S(h; \phi, t, \epsilon_f)$  reported by minimize\_scalar function in scipy.optimize and could be unreliable.

scheme	$h_{\dagger}$	$h^*$	r	#iters	#Evals	relative error	$\epsilon$
FD	2.00e-04	2.72e-04	2.08	1	3	1.86e-05	1.00e-08
FD	6.32e-04	8.59e-04	2.00	1	3	4.21e-05	1.00e-0'
FD	3.50e-03	2.73e-03	4.79	4	8	1.17e-03	1.00e-06
FD	6.32e-03	8.64e-03	1.77	1	3	1.72e-03	1.00e-05
FD	2.00e-02	2.76e-02	2.00	1	3	1.69e-03	1.00e-04
FD	6.32e-02	9.05e-02	1.73	1	3	5.07 e-04	1.00e-0
FD	5.00e-01	1.73e + 00	3.89	3	6	9.86e-02	1.00e-0
FD	6.32e-01	8.26e + 00	1.52	1	3	2.97e-01	1.00e-0
CD	3.11e-03	3.29e-03	2.40	1	4	1.89e-06	1.00e-0
CD	6.69e-03	7.09e-03	2.66	1	4	5.39e-06	1.00e-0
CD	1.44e-02	1.53e-02	2.72	1	4	9.35e-06	1.00e-0
CD	3.11e-02	3.29e-02	2.05	1	4	3.34e-04	1.00e-0
CD	6.69e-02	7.09e-02	2.18	1	4	1.33e-03	1.00e-0
CD	1.44e-01	1.53e-01	2.55	1	4	3.32e-03	1.00e-0
CD	3.11e-01	3.30e-01	1.89	1	4	3.84e-02	1.00e-0
CD	6.69e-01	7.74e + 00	2.01	1	4	5.71e-02	1.00e-0
FD_3P	3.91e-03	4.14e-03	2.88	1	5	1.01e-05	1.00e-0
FD_3P	8.43e-03	8.92e-03	3.24	1	5	2.01e-05	1.00e-0
FD_3P	1.82e-02	1.92e-02	2.76	1	$\tilde{5}$	2.05e-04	1.00e-0
FD_3P	3.91e-02	4.11e-02	3.28	1	$\tilde{5}$	5.82e-04	1.00e-0
FD_3P	8.43e-02	8.77e-02	2.86	1	$\tilde{5}$	5.65e-03	1.00e-0
FD_3P	1.82e-01	1.86e-01	3.10	1	5	2.29e-02	1.00e-0
FD_3P	3.91e-01	2.99e + 00	2.88	1	5	9.96e-02	1.00e-0
FD 3P	2.53e+00	2.12e+01	5.75	2	7	4.17e-01	1.00e-0
FD 4P	2.16e-02	2.04e-02	9.36	5	18	2.14e-06	1.00e-0
FD 4P	4.61e-02	3.67e-02	16.51	4	13	1.14e-05	1.00e-0
FD_4P	8.19e-02	4.76e-01	10.80	4	13	9.43e-05	1.00e-0
FD_4P	1.70e-01	1.25e-01	4.40	5	18	6.15e-04	1.00e-0
FD 4P	2.59e-01	$3.28e \pm 00$	14.62	4	13	7.80e-04	1.00e-0
FD_4P	3.07e-01	3.28e + 00	4.23	1	6	5.46e-03	1.00e-0
FD_4P	5.46e-01	$8.78e \pm 00$	6.44	1	6	3.19e-02	1.00e-0
FD 4P	2.91e+00	$8.79e \pm 00$	4.28	2	8	9.55e-01	1.00e-0
CD_4P	4.08e-02	4.22e-02	2.52	1	6	1.16e-07	1.00e-0
CD_4P	6.46e-02	6.69e-02	2.04	1	6	8.34e-07	1.00e-0
CD_4P	1.02e-01	1.06e-01	1.95	1	6	8.81e-06	1.00e-0
CD_4P	1.62e-01	1.68e-01	1.79	1	6	4.29e-05	1.00e-0
CD_4P	2.57e-01	2.67e-01	1.71	1	6	4.00e-04	1.00e-0
CD 4P	4.08e-01	4.25e-01	1.89	1	6	8.32e-04	1.00e-0
CD_4P	8.07e-01	7.97e+00	4.83	4	20	5.87e-03	1.00e-0
CD_4P	1.54e + 00	2.06e + 01	4.25	3	_0 14	2.34e-01	1.00e-0
L2_CD	3.29e-02	3.07e-02	3.78	4	15	1.00e-04	1.00e-0
L2_CD	4.68e-02	5.46e-02	1.89	1	-0	8.55e-05	1.00e-0
L2 CD	1.04e-01	9.71e-02	4.22	4	15	7.28e-04	1.00e-0
L2 CD	1.48e-01	1.73e-01	1.90	1	5	1.06e-03	1.00e-0
L2 CD	3.29e-01	3.07e-01	4 03	4	15	8.28e-03	1.00e-0
L2 CD	5.85e-01	5.49e-01	3.81	4	15	2.84e-02	1.00e-0
	$1.04e \pm 00$	9.87e-01	345	4	15	7.95e-02	1.00e-0
			0.10		т. <del>С</del>	1.000 04	

abscheme  $h_{\dagger}$ h' #iters #Evals relative error r0.10 0.10 2.53e + 003.94e + 00FD 1.89 $\mathbf{5}$ 8 1.75e-020.101.00FD 3.48e-013.94e-014.346 115.29e-0210.002.77e-022.682.01e-050.10FD 3.94e-0248  $1.39\mathrm{e}{+00}$ 1.82e + 002.617121.000.10FD 1.30e-024.913  $\mathbf{6}$ 1.001.00FD 1.58e-011.82e-013.01e-03 1.0010.00FD 1.58e-021.82e-024.58 $\mathbf{2}$ 4 6.97e-03 10.00FD 6.32e-01 8.44e-01 3.234 74.76e-040.1010.00 1.00FD 6.32e-028.44e-02 3.381 3 2.83e-049 10.0010.00FD 6.92e-038.44e-033.34 $\mathbf{5}$ 2.92e-03 10 0.100.10CD  $3.89\mathrm{e}{+00}$ 3.12e + 005.474 2.51e-02100.101.00CD 2.88e-013.12e-012.303 1.38e-020.1010.00CD 3.20e-023.12e-023.134 141.70e-021.000.10CD 1.30e + 001.44e + 002.173 8 2.81e-03 1.44e-01 1.44e-01 2.971.001.00CD 1 4 3.46e-031.0010.00CD 1.60e-021.44e-024.063 104.27e-0310.00CD 6.49e-016.70e-012.73160.107.02e-04510.00CD 6.70e-023.74141.007.21e-024 8.66e-04 1210.0010.00CD 5.34e-036.70e-03 1.524 4.75e-040.100.10FD\_3P 4.91e + 008.14e + 012.934 112.13e-021.00FD\_3P 5.45e-018.14e + 002.45276.48e-020.100.1010.00FD\_3P 7.57e-028.14e-01 4.905151.61e-01 1.000.10 FD\_3P 1.64e + 002.14e+012.253 9 8.79e-03 51.001.00FD\_3P 1.82e-012.14e+003.631 1.86e-043 9 1.0010.00FD\_3P 2.02e-021.83e-024.431.15e-02FD\_3P 8.18e-01 1310.000.108.45e-01 3.48 $\mathbf{5}$ 1.71e-03 10.001.00FD\_3P 9.09e-028.45e-024.584 11 1.77e-0310.0010.00FD\_3P 1.01e-028.45e-03 6.676 151.00e-03 0.100.10FD\_4P 9.33e + 008.44e + 014.959 311.63e-011.00FD\_4P 9.79e-018.44e + 009.64 8 321.50e-010.10 $1.96e{+}00$ 7 2910.00FD\_4P 1.78e+017.719.60e-01 0.103 111.000.10FD\_4P 2.76e + 003.59e + 004.473.59e-031.001.00FD\_4P 3.07e-013.59e-016.551 68.16e-03 1.0010.00FD\_4P 6.62e-01 5.88e + 0010.158 359.49e-01 10.00 0.10FD\_4P 1.84e + 002.25e + 006.574 14 3.68e-0410.001.00FD\_4P 2.05e-012.71e+0010.683 10 2.90e-0410.00 FD\_4P 3.07e-014.62e + 0016.476 10.001 8.38e-01 6.52e + 002.120.100.10 $CD_4P$ 7.97e + 015145.73e-03 3 0.101.00CD\_4P 6.11e-01 7.97e + 001.57144.45e-030.1010.00CD\_4P 7.64e-027.97e-01 4.32 $\mathbf{5}$ 181.06e-021.000.10 $CD_4P$ 4.08e + 004.10e + 002.307269.02e-04 1.001.00CD\_4P 4.08e-014.10e-012.301 6 9.02e-04 1.0010.00 CD\_4P 3.82e-02 4.10e-021.686 206.98e-042.45e + 002.58e+0010.000.10CD\_4P 181.8951.18e-0410.001.00 $CD_4P$ 2.55e-012.58e-012.314 201.39e-0410.0010.00 $CD_4P$ 2.55e-022.58e-022.315141.39e-04

Table B.6. Detailed results for  $\phi(t) = a \cdot \sin(b \cdot t)$  with  $\epsilon_f = 1E-3$ ; r represents the final testing ratio;  $h^*$  is the h that minimizes  $\delta_S(h; \phi, t, \epsilon_f)$  reported by minimize\_scalar function in scipy.optimize and could be unreliable.



Figure B.1. Worst case relative error  $\delta_S(h; \phi, t, \epsilon_f)$  against h on function  $\phi(t) = \cos(t)$  with different noise levels; the vertical dashed line represents the  $h_{\dagger}$  output by Algorithm 3.2.



Figure B.2. Worst case relative error  $\delta_S(h; \phi, t, \epsilon_f)$  against h on function  $\phi(t) = a \cdot \sin(b \cdot t)$  for different a and b; the vertical dashed line represents the  $h_{\dagger}$  output by Algorithm 3.2.

#### **B.2.3.** Difficult and Special Examples

Here, we present the full table of results for the examples listed in Section 3.4 in Table B.8 with  $\epsilon_f = 10^{-3}$ . For reference, the considered problems are:

Table B.7. Detailed results for  $\phi(t) = a \cdot \sin(b \cdot t)$  with  $\epsilon_f = 1E-3$ ; r represents the final testing ratio;  $h^*$  is the h that minimizes  $\delta_S(h; \phi, t, \epsilon_f)$  reported by minimize\_scalar function in scipy.optimize and could be unreliable.

a	b	scheme	$h_\dagger$	$h^*$	r	#iters	#Evals	relative error
0.10	0.10	L2_CD	9.36e + 00	8.42e + 00	4.39	8	23	5.97e-02
0.10	1.00	L2_CD	9.36e-01	8.42e-01	4.79	2	7	4.28e-02
0.10	10.00	L2_CD	8.78e-02	8.99e-01	3.28	5	15	5.32e-02
1.00	0.10	L2_CD	3.74e + 00	4.70e + 00	1.99	4	11	9.49e-03
1.00	1.00	$L2_CD$	4.68e-01	4.70e-01	2.57	1	5	2.13e-02
1.00	10.00	$L2_CD$	4.39e-02	4.70e-02	2.15	6	17	1.70e-02
10.00	0.10	L2_CD	$2.81e{+}00$	2.64e + 00	3.59	5	15	7.63e-03
10.00	1.00	L2_CD	2.34e-01	2.64e-01	2.62	2	7	6.94e-04
10.00	10.00	L2_CD	2.93e-02	2.64e-02	5.02	5	13	3.94e-03

(1) 
$$\phi(t) = (e^t - 1)^2$$
, at  $t = -8$ .

- (2)  $\phi(t) = e^{100t}$ , at t = 0.01.
- (3)  $\phi(t) = t^4 + 3t^2 10t$ , at t = 0.999999.
- (4)  $\phi(t) = 10000t^3 + 0.01t^2 + 5t$ , at  $t = 10^{-9}$ .

#### B.2.4. Comparison with Moré-Wild Heuristic

We compare our adaptive forward-difference procedure against the Moré-Wild heuristic [71], as described in Section 3.2.1.

First, observe that if function  $\phi$  has (near) central symmetry at t, then Moré-Wild heuristic is very likely to fail. To demonstrate this, we test on  $\phi(t) = \sin(t)$  with various value of t close to 0 and different noise levels  $\epsilon_f$ . The results are summarized in Table B.9.

Next, we test our adaptive procedure and the Moré-Wild heuristic on  $\phi(t) = a \cdot (\exp(b \cdot t) - 1)$  at t = 0, with a fixed noise level:  $\epsilon_f = 1E-3$ . We summarize our result in Table B.10. Notice that Moré-Wild heuristic may not be able to find a suitable estimation for h, in which case a failure is declared. In such cases, we will report the result as "--".

$\phi(t)$	scheme	$h_{\dagger}$	$h^*$	r	#iters	#Evals	relative error
$(e^t - 1.0)^2$	FD	1.01e+00	1.46e+00	4.49	3	5	1.02e+00
$(e^t - 1.0)^2$	CD	$1.30e{+}00$	$1.53e{+}00$	3.38	3	8	5.73e-02
$(e^t - 1.0)^2$	FD_3P	8.18e-01	$3.82e{+}02$	2.28	5	13	6.63e-01
$(e^t - 1.0)^2$	$FD_4P$	9.21e-01	3.82e + 02	4.14	2	8	$4.15e{+}00$
$(e^t - 1.0)^2$	$CD_4P$	$1.43e{+}00$	3.82e + 02	1.84	5	22	$1.11e{+}00$
$(e^t - 1.0)^2$	L2_CD	$2.34e{+}00$	8.68e + 00	3.03	6	19	3.90e-01
$e^{100t}$	FD	4.32e-04	3.79e-04	3.72	7	11	2.74e-02
$e^{100t}$	CD	1.19e-03	1.03e-03	4.29	7	20	3.09e-03
$e^{100t}$	FD_3P	1.12e-03	3.82e + 02	3.08	8	19	6.81e-03
$e^{100t}$	FD_4P	1.90e-03	3.82e + 02	6.54	8	23	4.97e-03
$e^{100t}$	CD_4P	3.18e-03	3.82e + 02	2.15	8	20	4.60e-04
$e^{100t}$	L2_CD	3.66e-03	3.64e-03	3.01	8	19	1.18e-02
$t^4 + 3t^2 - 10t$	FD	1.58e-02	1.48e-02	3.55	2	4	7.97e + 02
$t^4 + 3t^2 - 10t$	CD	4.81e-02	5.00e-02	2.94	2	6	2.15e+01
$t^4 + 3t^2 - 10t$	FD_3P	6.06e-02	6.16e-02	3.64	2	7	2.38e+02
$t^4 + 3t^2 - 10t$	FD_4P	1.54e-01	1.39e-01	11.90	4	13	1.93e+02
$t^4 + 3t^2 - 10t$	$CD_4P$	9.39e + 02	4.87e + 03	1.62	16	48	2.71e-03
$t^4 + 3t^2 - 10t$	L2_CD	2.34e-01	2.11e-01	4.53	2	7	5.38e-03
$10000t^3 + 0.01t^2 + 5t$	FD	3.95e-03	4.64 e- 03	4.36	3	5	5.46e-02
$10000t^3 + 0.01t^2 + 5t$	CD	3.56e-03	3.68e-03	2.63	6	18	4.02e-02
$10000t^3 + 0.01t^2 + 5t$	FD_3P	4.49e-03	4.64 e- 03	3.65	6	15	1.17e-02
$10000t^3 + 0.01t^2 + 5t$	FD_4P	6.72e + 02	3.20e + 03	11.72	8	22	2.37e-06
$10000t^3 + 0.01t^2 + 5t$	CD_4P	8.35e + 02	$1.03e{+}04$	1.95	12	28	1.99e-07
$10000t^3 + 0.01t^2 + 5t$	L2_CD	$9.59\mathrm{e}{+02}$	$2.84\mathrm{e}{+03}$	1.95	12	27	7.49e-08

Table B.8. Detailed results for special examples, with  $\epsilon_f = 1E-3$ ; r represents the final testing ratio;  $h^*$  is the h that minimizes  $\delta_S(h; \phi, t, \epsilon_f)$  reported by minimize\_scalar function in scipy.optimize and could be unreliable.

We can see that when the Moré-Wild heuristic does not declare a failure, it usually outputs an interval h that is quite close to our procedure and produces similar relative error as ours. However, there are many cases where Moré-Wild heuristic fails, while our procedure works very robustly in all cases.

#### **B.2.5.** Finite-Difference L-BFGS

We present the total number of function evaluations and final optimality gap  $\phi(x) - \phi^*$ used by each method in Tables B.11–B.18.

Table B.9. Comparison between the Moré-Wild heuristic against our adaptive procedure on function  $\phi(t) = \sin(t)$  with various  $\epsilon_f$  and t. We use "--" to report the cases where Moré-Wild heuristic fails. Subscript "MW" labels the results corresponding to Moré-Wild heuristic, and subscript "ada" labels the results corresponding to our adaptive procedure;  $\delta$  is the relative error, and  $\overline{\delta}$  is the worst-case relative error.

$\epsilon_{f}$	t	$h_{\rm MW}$	$\delta_{ m MW}$	$\overline{\delta}_{\mathrm{MW}}$	$h_{\rm ada}$	$\delta_{ m ada}$	$\overline{\delta}_{\mathrm{ada}}$
1.00e-08	1.00e-08				3.20e-03	0.000	0.000
1.00e-08	1.00e-06				3.20e-03	0.000	0.000
1.00e-08	1.00e-04	1.68e-02	0.000	0.000	3.20e-03	0.000	0.000
1.00e-08	1.00e-02	1.68e-03	0.000	0.000	2.00e-03	0.000	0.000
1.00e-08	0.00e+00				3.20e-03	0.000	0.000
1.00e-06	1.00e-08				1.40e-02	0.000	0.000
1.00e-06	1.00e-06				1.40e-02	0.000	0.000
1.00e-06	1.00e-04				1.40e-02	0.000	0.000
1.00e-06	1.00e-02	1.69e-02	0.000	0.000	1.40e-02	0.000	0.000
1.00e-06	0.00e+00				1.40e-02	0.000	0.000
1.00e-04	1.00e-08				5.00e-02	0.001	0.004
1.00e-04	1.00e-06				6.50e-02	0.003	0.004
1.00e-04	1.00e-04				5.00e-02	0.001	0.004
1.00e-04	1.00e-02				5.00e-02	0.001	0.005
1.00e-04	0.00e+00				6.50e-02	0.000	0.004
1.00e-02	1.00e-08				2.00e-01	0.058	0.107
1.00e-02	1.00e-06	5.20e-01	0.067	0.083	3.50e-01	0.018	0.077
1.00e-02	1.00e-04				3.50e-01	0.019	0.077
1.00e-02	1.00e-02				3.50e-01	0.029	0.079
1.00e-02	0.00e+00	6.10e-01	0.085	0.094	3.50e-01	0.032	0.077

In general, our adaptive procedure is more robust to different noise levels. Our method only fails when the initial choice of h is not sufficiently small to initially identify the local behavior of the function. This can be seen, for example, with the BOX2 example. On the other hand, Moré and Wild's heuristic frequently fails when the noise level is large (for example, with  $\epsilon_f = 10^{-1}$ ). This is due both to the case where  $\phi(x) \approx 0$  and hence (3.2.11) fails, as well as the case where two iterations are insufficient to find an h that satisfies their conditions. In both cases, we denote a failure case with \*. As expected, using a fixed interval is always efficient, but may perform poorly when the Hessian in the function changes, as described in Section 3.4.

Table B.10. Comparison between the Moré-Wild heuristic against our adaptive procedure on function  $\phi(t) = a \cdot (\exp(b \cdot t) - 1)$  with  $\epsilon_f = 1\text{E-3}$  at t = 0. We use "--" to report the cases where Moré-Wild heuristic fails. Subscript "MW" labels the results corresponding to Moré-Wild heuristic, and subscript "ada" labels the results corresponding to our adaptive procedure;  $\delta$  is the relative error, and  $\overline{\delta}$  is the worst-case relative error.

a	b	$h_{\rm MW}$	$\delta_{\mathrm{MW}}$	$\overline{\delta}_{ m MW}$	$h_{\rm ada}$	$\delta_{ m ada}$	$\overline{\delta}_{ m ada}$
0.01	0.01				4.05e+01	0.135	0.727
0.01	0.10				4.05e+00	0.145	0.727
0.01	1.00				4.43e-01	0.275	0.710
0.01	10.00	4.68e-02	0.177	0.702	3.95e-02	0.084	0.732
0.01	100.00				3.95e-03	0.123	0.732
0.10	0.01				1.62e + 01	0.032	0.209
0.10	0.10				1.77e+00	0.074	0.207
0.10	1.00	1.64e-01	0.072	0.209	1.58e-01	0.095	0.210
0.10	10.00	1.62e-02	0.096	0.209	1.58e-02	0.055	0.210
0.10	100.00				1.73e-03	0.078	0.207
1.00	0.01				7.08e+00	0.041	0.065
1.00	0.10				6.32e-01	0.034	0.064
1.00	1.00	5.26e-02	0.029	0.065	6.32e-02	0.036	0.064
1.00	10.00	5.28e-03	0.012	0.065	6.92e-03	0.014	0.064
1.00	100.00				6.18e-04	0.009	0.064
10.00	0.01				2.53e+00	0.012	0.021
10.00	0.10	1.76e-01	0.011	0.020	2.53e-01	0.009	0.021
10.00	1.00	1.68e-02	0.006	0.020	1.58e-02	0.005	0.021
10.00	10.00	1.68e-03	0.002	0.020	2.47e-03	0.014	0.021
10.00	100.00				2.47e-04	0.010	0.021
100.00	0.01				6.32e-01	0.003	0.006
100.00	0.10	5.39e-02	0.002	0.006	6.32e-02	0.004	0.006
100.00	1.00	5.31e-03	0.001	0.006	6.92e-03	0.001	0.006
100.00	10.00	5.31e-04	0.000	0.006	6.18e-04	0.001	0.006
100.00	100.00				6.18e-05	0.001	0.006

Table B.11. Total number of function evaluations used and final accuracy achieved by forward-difference L-BFGS method with different choices of the finite-difference interval.

			Fixe	d Interval	Mo	oré-Wild	Adaptive	
Problem	n	$\epsilon_f$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$
AIRCRFTB	5	$1 \times 10^{-1}$	1373	7.651	1457	$4.777 \times 10^{-1}$	559	$4.950 \times 10^{-1}$
ALLINITU	4	$1 \times 10^{-1}$	656	2.767	972	$3.708 \times 10^{-1}$	686	$1.003 \times 10^{-1}$
ARWHEAD	100	$1 \times 10^{-1}$	2513	$2.027\times10^{-1}$	8784	$3.040 \times 10^{-1}$	2811	$2.458\times10^{-1}$
BARD	3	$1 \times 10^{-1}$	650	$8.311\times10^{-1}$	602	$7.581\times10^{-2}$	670	$3.942 \times 10^{-1}$
BDQRTIC	100	$1 \times 10^{-1}$	2601	6.278	9213	5.902	5246	4.438
BIGGS3	3	$1 \times 10^{-1}$	650	1.195	726	1.348	674	1.545
BIGGS5	5	$1 \times 10^{-1}$	662	1.294	1457	1.319	696	1.359
BIGGS6	6	$1 \times 10^{-1}$	668	$4.026 \times 10^{-1}$	853	$5.483 \times 10^{-1}$	693	$6.766 \times 10^{-1}$
BOX2	$^{2}$	$1 \times 10^{-1}$	644	$4.345 \times 10^{-2}$	709	$4.095 \times 10^{-2}$	1975	$3.406^{*}$
BOX3	3	$1 \times 10^{-1}$	650	$4.843 \times 10^{-2}$	1385	$1.887 \times 10^{-1}$	676	$7.527 \times 10^{-2}$
BRKMCC	$^{2}$	$1 \times 10^{-1}$	644	$4.391 \times 10^{-2}$	709	$5.850 \times 10^{-3}$	658	$1.674 \times 10^{-1}$
BROWNAL	100	$1 \times 10^{-1}$	2513	$5.097\times10^{-2}$	10800	$4.318\times10^{-2}$	3293	$1.250 \times 10^{-2}$
BROWNDEN	4	$1 \times 10^{-1}$	656	$1.159\times10^{-1}$	1364	$1.568\times10^{-1}$	1328	$1.465\times10^{-1}$
CLIFF	2	$1 \times 10^{-1}$	644	$2.902 \times 10^2$	861	$1.180 \times 10^1$	1593	$3.385 \times 10^{-1}$
CRAGGLVY	100	$1 \times 10^{-1}$	8011	$1.850 \times 10^2$	207136	$5.048 \times 10^{2*}$	11498	$1.906 \times 10^{1}$
CUBE	2	$1 \times 10^{-1}$	644	$4.352 \times 10^{-2}$	6340	$4.227^{*}$	659	$8.670 \times 10^{-2}$
DENSCHND	3	$1 \times 10^{-1}$	650	$2.254 \times 10^2$	2062	$2.627 \times 10^{-1}$	1018	$9.226 \times 10^{-3}$
DENSCHNE	3	$1 \times 10^{-1}$	650	1.151	2062	1.059	542	1.018
DIXMAANH	300	$1 \times 10^{-1}$	3713	$1.382 \times 10^{1}$	22922	$1.392 \times 10^{1}$	7809	$1.280 \times 10^{1}$
DQRTIC	100	$1 \times 10^{-1}$	2513	$1.743 \times 10^2$	26299	4.930	10912	$4.944\times10^{-1}$
EDENSCH	36	$1 \times 10^{-1}$	2129	1.674	4995	1.642	4624	3.352
EIGENALS	110	$1 \times 10^{-1}$	2609	$5.105 \times 10^1$	63118	$1.496 \times 10^1$	11630	$1.236 \times 10^1$
EIGENBLS	110	$1 \times 10^{-1}$	5103	6.717	8306	$1.292 \times 10^1$	5098	6.533
EIGENCLS	30	$1 \times 10^{-1}$	1452	1.401	10899	$1.218 \times 10^1$	4088	1.907
ENGVAL1	100	$1 \times 10^{-1}$	2295	$7.316\times10^{-1}$	610646	$2.114\times10^{2*}$	4956	1.186
EXPFIT	2	$1 \times 10^{-1}$	644	$6.682 \times 10^{-1}$	31080	$1.270 \times 10^{2*}$	659	$6.881 \times 10^{-2}$
<b>FLETCBV3</b>	100	$1 \times 10^{-1}$	2549	$1.785 \times 10^{5}$	9488	$1.785 \times 10^{5}$	21950	$8.454 \times 10^{3}$
FLETCHBV	100	$1 \times 10^{-1}$	40856	$-1.175 \times 10^{9}$	135927	$1.154 \times 10^{9}$	166981	$5.696 \times 10^{9}$
FREUROTH	100	$1 \times 10^{-1}$	2295	$7.097 \times 10^1$	9658	$5.756 \times 10^1$	4008	$5.912 \times 10^1$
GENROSE	100	$1 \times 10^{-1}$	4326	$1.421 \times 10^2$	36745	$1.321 \times 10^2$	5347	$1.364 \times 10^2$
GULF	3	$1 \times 10^{-1}$	650	6.713	1385	6.820	670	6.664
HAIRY	$^{2}$	$1 \times 10^{-1}$	644	$9.325 \times 10^1$	1328	$4.889\times 10^{2*}$	1008	$7.985\times10^{-2}$
HELIX	3	$1 \times 10^{-1}$	650	7.601	745	7.591	671	7.656
NCB20B	100	$1 \times 10^{-1}$	2295	$4.793 \times 10^{-1}$	5531	$2.000 \times 10^{-1}$	5393	$1.280 \times 10^{-1}$
NONDIA	100	$1 \times 10^{-1}$	2513	$5.546 \times 10^{-1}$	10800	$4.717 \times 10^{-1}$	5076	$3.770 \times 10^{-1}$
NONDQUAR	100	$1 \times 10^{-1}$	2613	$8.483 \times 10^{-1}$	33658	7.457	5707	$3.571 \times 10^{-1}$
OSBORNEA	5	$1 \times 10^{-1}$	662	$1.748 \times 10^{-1}$	52825	$1.142 \times 10^{2*}$	2869	$1.722 \times 10^{-1}$
OSBORNEB	11	$1 \times 10^{-1}$	698	1.178	1014	3.021	769	1.851
PENALTY1	100	$1 \times 10^{-1}$	8895	$1.008 \times 10^2$	20136	1.131	12532	1.268
PFIT1LS	3	$1 \times 10^{-1}$	17316	$1.346 \times 10^2$	745	4.139	675	7.796
PFIT2LS	3	$1 \times 10^{-1}$	20616	$2.684 \times 10^2$	1385	$2.350 \times 10^1$	1991	4.006
PFIT3LS	3	$1 \times 10^{-1}$	20088	$9.024 \times 10^2$	2126	2.933	1988	2.171
PFIT4LS	3	$1 \times 10^{-1}$	21702	$2.771 \times 10^{3}$	3839	2.704	1989	9.426
QUARTC	100	$1 \times 10^{-1}$	2513	$1.743 \times 10^{2}$	26299	4.930	10912	$4.944 \times 10^{-1}$
SINEVAL	$^{2}$	$1 \times 10^{-1}$	644	5.576	566	4.712	664	5.554
SINQUAD	100	$1 \times 10^{-1}$	4326	$1.254 \times 10^1$	28209	$3.331 \times 10^1$	5281	9.760
SISSER	$^{2}$	$1 \times 10^{-1}$	644	$3.393\times10^{-2}$	2026	$1.611\times 10^{-1}$	664	$7.452\times10^{-3}$
SPARSQUR	100	$1 \times 10^{-1}$	4326	1.365	506948	$6.928\times10^{1*}$	5592	$1.209\times 10^{-1}$
TOINTGSS	100	$1 \times 10^{-1}$	2513	$2.071\times 10^1$	8784	$1.467\times10^{1}$	4133	9.678
TQUARTIC	100	$1 \times 10^{-1}$	2513	1.108	62958	$3.128^{*}$	3391	$7.331\times10^{-1}$
TRIDIA	100	$1 \times 10^{-1}$	2513	$8.478 \times 10^1$	18300	$1.691 \times 10^1$	5677	$3.935 \times 10^1$
WATSON	31	$1 \times 10^{-1}$	1881	$7.753 \times 10^{-1}$	4702	1.993	2174	2.873
WOODS	100	$1 \times 10^{-1}$	10414	$8.618 \times 10^1$	16987	$2.907 \times 10^2$	5255	$5.533 \times 10^1$
ZANGWIL2	2	$1 \times 10^{-1}$	644	$2.686 \times 10^{-2}$	1349	$3.307 \times 10^{-2}$	657	$3.852 \times 10^{-2}$

Table B.12. Total number of function evaluations used and final accuracy achieved by forward-difference L-BFGS method with different choices of the finite-difference interval.

Fixed Int			d Interval	Mo	oré-Wild	Adaptive		
Problem	n	$\epsilon_{f}$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$
AIRCRFTB	5	$1 \times 10^{-3}$	1008	$3.772 \times 10^{-1}$	817	$3.313 \times 10^{-1}$	560	$3.286 \times 10^{-1}$
ALLINITU	4	$1 \times 10^{-3}$	656	$2.354 \times 10^{-3}$	563	$1.858 \times 10^{-3}$	686	$1.172 \times 10^{-3}$
ARWHEAD	100	$1 \times 10^{-3}$	2513	$4.510 \times 10^{-2}$	7331	$4.407 \times 10^{-2}$	4951	$4.160 \times 10^{-2}$
BARD	3	$1 \times 10^{-3}$	650	$2.093 \times 10^{-3}$	602	$3.467 \times 10^{-3}$	538	$3.273 \times 10^{-3}$
BDQRTIC	100	$1 \times 10^{-3}$	6017	$9.588 \times 10^{-2}$	8733	$1.151 \times 10^{-1}$	8551	$9.804 \times 10^{-2}$
BIGGS3	3	$1 \times 10^{-3}$	650	$8.010 \times 10^{-4}$	712	$1.375 \times 10^{-2}$	672	$2.297 \times 10^{-2}$
BIGGS5	5	$1 \times 10^{-3}$	662	$1.813 \times 10^{-1}$	1457	$1.261 \times 10^{-1}$	680	$1.453 \times 10^{-1}$
BIGGS6	6	$1 \times 10^{-3}$	668	$2.904 \times 10^{-1}$	853	$2.907 \times 10^{-1}$	563	$2.909 \times 10^{-1}$
BOX2	2	$1 \times 10^{-3}$	644	$1.964 \times 10^{-4}$	709	$2.624 \times 10^{-6}$	661	$1.882 \times 10^{-5}$
BOX3	3	$1 \times 10^{-3}$	650	$9.556 \times 10^{-4}$	726	$4.567 \times 10^{-4}$	672	$3.568 \times 10^{-4}$
BRKMCC	2	$1 \times 10^{-3}$	644	$3.552 \times 10^{-3}$	566	$3.872 \times 10^{-3}$	661	$8.029 \times 10^{-4}$
BROWNAL	100	$1 \times 10^{-3}$	2513	$6.869 \times 10^{-4}$	10800	$4.648 \times 10^{-4}$	3293	$1.670 \times 10^{-4}$
BROWNDEN	4	$1 \times 10^{-3}$	656	$1.183 \times 10^{-3}$	706	$1.912 \times 10^{-3}$	1327	$8.187 \times 10^{-4}$
CLIFF	2	$1 \times 10^{-3}$	644	$2.902 \times 10^2$	28980	$8.220 \times 10^{-1}$	1008	$3.059 \times 10^{-4}$
CRAGGLVY	100	$1 \times 10^{-3}$	7993	4.852	8619	1.046	7272	$9.025 \times 10^{-1}$
CUBE	2	$1 \times 10^{-3}$	644	$4.201 \times 10^{-2}$	709	$4.951 \times 10^{-2}$	659	$4.381 \times 10^{-2}$
DENSCHND	3	$1 \times 10^{-3}$	631	4.109	1385	$1.170 \times 10^{-3}$	1018	$9.852 \times 10^{-4}$
DENSCHNE	3	$1 \times 10^{-3}$	650	$9.994 \times 10^{-1}$	745	1.000	672	1.002
DIXMAANH	300	$1 \times 10^{-3}$	7774	1.041	37580	$7.456 \times 10^{-2}$	15794	$9.375 \times 10^{-2}$
DQRTIC	100	$1 \times 10^{-3}$	6609	4.832	21149	$9.935 \times 10^{-3}$	11600	$1.713 \times 10^{-2}$
EDENSCH	36	$1 \times 10^{-3}$	1194	$5.382 \times 10^{-2}$	6085	$2.188 \times 10^{-2}$	3234	$3.748 \times 10^{-2}$
EIGENALS	110	$1 \times 10^{-3}$	4850	$3.394 \times 10^{-1}$	22676	$8.848 \times 10^{-2}$	8753	$7.456 \times 10^{-2}$
EIGENBLS	110	$1 \times 10^{-3}$	1932	2.084	21095	1.631	6825	1.740
EIGENCLS	30	$1 \times 10^{-3}$	1875	$7.554 \times 10^{-2}$	8449	$7.032 \times 10^{-2}$	3187	$8.686 \times 10^{-2}$
ENGVAL1	100	$1 \times 10^{-3}$	2295	$8.024 \times 10^{-2}$	9658	$2.733 \times 10^{-2}$	4841	$4.366 \times 10^{-2}$
EXPFIT	2	$1 \times 10^{-3}$	644	$6.733 \times 10^{-3}$	566	$3.801 \times 10^{-4}$	663	$3.592 \times 10^{-4}$
FLETCBV3	100	$1 \times 10^{-3}$	2549	$1.785 \times 10^{5}$	9488	$1.785 \times 10^{5}$	29889	$-1.181 \times 10^2$
FLETCHBV	100	$1 \times 10^{-3}$	33825	$-7.812 \times 10^8$	141125	$6.878 \times 10^7$	62282	$-6.875 \times 10^{8}$
FREUROTH	100	$1 \times 10^{-3}$	4861	$2.366 \times 10^{-1}$	9725	$2.041 \times 10^{-2}$	8226	$4.143 \times 10^{-2}$
GENROSE	100	$1 \times 10^{-3}$	5779	$1.109 \times 10^{2}$	17379	$1.108 \times 10^{2}$	7561	$1.110 \times 10^{2}$
GULF	3	$1 \times 10^{-3}$	650	6.626	599	6.622	672	6.622
HAIRY	2	$1 \times 10^{-3}$	644	$7.979 \times 10^{-3}$	566	$7.030 \times 10^{-4}$	853	$3.304 \times 10^{-3}$
HELIX	3	$1 \times 10^{-3}$	650	$4.079 \times 10^{-3}$	1091	$2.495 \times 10^{-4}$	671	$3.746 \times 10^{-4}$
NCB20B	100	$1 \times 10^{-3}$	2295	$1.400 \times 10^{-2}$	8984	$5.856 \times 10^{-3}$	6650	$4.462 \times 10^{-3}$
NONDIA	100	$1 \times 10^{-3}$	2513	$4.911 \times 10^{-1}$	13295	$4.862 \times 10^{-1}$	8550	$4.612 \times 10^{-1}$
NONDQUAR	100	$1 \times 10^{-3}$	5534	$9.503 \times 10^{-2}$	29253	$1.247 \times 10^{-2}$	8653	$3.602 \times 10^{-2}$
OSBORNEA	5	$1 \times 10^{-3}$	662	$1.525 \times 10^{-1}$	54694	2.093	553	$1.567 \times 10^{-1}$
OSBORNEB	11	$1 \times 10^{-3}$	555	$3.697 \times 10^{-1}$	2314	$3.170 \times 10^{-1}$	766	$3.115 \times 10^{-1}$
PENALTY1	100	$1 \times 10^{-3}$	4326	$2.200 \times 10^{1}$	14928	$1.345 \times 10^{-3}$	12554	$2.741 \times 10^{-4}$
PFIT1LS	3	$1 \times 10^{-3}$	17244	7.787	2062	$6.076 \times 10^{-2}$	1312	$4.829\times10^{-2}$
PFIT2LS	3	$1 \times 10^{-3}$	19169	$1.147 \times 10^{2}$	2062	$5.424 \times 10^{-2}$	1310	$8.222 \times 10^{-2}$
PFIT3LS	3	$1 \times 10^{-3}$	21645	$6.633 \times 10^{2}$	2062	$2.615 \times 10^{-2}$	1993	$4.635 \times 10^{-2}$
PFIT4LS	3	$1 \times 10^{-3}$	22928	$2.305 \times 10^{3}$	2631	$2.308 \times 10^{-1}$	1992	$2.053 \times 10^{-1}$
QUARTC	100	$1 \times 10^{-3}$	6609	4.832	21149	$9.935 \times 10^{-3}$	11600	$1.713 \times 10^{-2}$
SINEVAL	2	$1 \times 10^{-3}$	644	1.557	2026	$2.780 \times 10^{-1}$	1299	$1.695 \times 10^{-1}$
SINQUAD	100	$1 \times 10^{-3}$	4326	$3.852 \times 10^{-2}$	19928	$1.067 \times 10^{-1}$	5543	$7.637 \times 10^{-2}$
SISSER	2	$1 \times 10^{-3}$	644	$2.069 \times 10^{-3}$	566	$1.051 \times 10^{-3}$	664	$2.094 \times 10^{-4}$
SPARSQUR	100	$1 \times 10^{-3}$	6609	$3.112 \times 10^{-2}$	14021	$8.064 \times 10^{-2}$	6625	$3.022 \times 10^{-3}$
TOINTGSS	100	$1 \times 10^{-3}$	2513	$2.198 \times 10^{-1}$	5071	$1.026 \times 10^{-1}$	3730	$6.133 \times 10^{-2}$
TQUARTIC	100	$1 \times 10^{-3}$	2513	$7.227 \times 10^{-1}$	7529	$7.241 \times 10^{-1}$	4090	$7.220 \times 10^{-1}$
TRIDIA	100	$1 \times 10^{-3}$	7921	$3.607 \times 10^{-1}$	28195	$3.539 \times 10^{-1}$	9887	$4.283 \times 10^{-1}$
WATSON	31	$1 \times 10^{-3}$	1881	$1.862 \times 10^{-1}$	10162	$2.201 \times 10^{-2}$	2919	$1.754 \times 10^{-1}$
WOODS	100	$1 \times 10^{-3}$	11923	$7.268 \times 10^{-1}$	10536	6.643	5493	6.472
ZANGWIL2	2	$1 \times 10^{-3}$	644	$8.620\times 10^{-4}$	566	$8.797 \times 10^{-4}$	657	$6.881\times10^{-4}$

Table B.13. Total number of function evaluations used and final accuracy achieved by forward-difference L-BFGS method with different choices of the finite-difference interval.

			Fixe	d Interval	Mo	oré-Wild	A	daptive
Problem	n	$\epsilon_{f}$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$
AIRCRFTB	5	$1 \times 10^{-5}$	662	$3.023 \times 10^{-1}$	1880	$2.674 \times 10^{-4}$	2014	$4.084 \times 10^{-4}$
ALLINITU	4	$1 \times 10^{-5}$	656	$1.835 \times 10^{-5}$	706	$1.146 \times 10^{-5}$	682	$6.468 \times 10^{-6}$
ARWHEAD	100	$1 \times 10^{-5}$	3832	$1.827 \times 10^{-4}$	11900	$3.262 \times 10^{-4}$	7570	$5.117\times10^{-4}$
BARD	3	$1 \times 10^{-5}$	650	$1.894 \times 10^{-3}$	602	$1.878 \times 10^{-3}$	674	$1.903 \times 10^{-3}$
BDQRTIC	100	$1 \times 10^{-5}$	10225	$3.227 \times 10^{-3}$	14968	$2.305\times10^{-4}$	9860	$5.248 \times 10^{-4}$
BIGGS3	3	$1 \times 10^{-5}$	631	$2.826 \times 10^{-4}$	712	$2.404 \times 10^{-4}$	670	$1.581 \times 10^{-4}$
BIGGS5	5	$1 \times 10^{-5}$	1008	$2.170 \times 10^{-2}$	3911	$1.064 \times 10^{-4}$	2030	$9.355 \times 10^{-3}$
BIGGS6	6	$1 \times 10^{-5}$	1014	$1.863 \times 10^{-3}$	2170	$9.419 \times 10^{-5}$	1348	$4.434 \times 10^{-4}$
BOX2	2	$1 \times 10^{-5}$	644	$3.493 \times 10^{-6}$	1349	$8.092 \times 10^{-6}$	658	$2.480 \times 10^{-6}$
BOX3	3	$1 \times 10^{-5}$	650	$7.493 \times 10^{-6}$	726	$4.612 \times 10^{-6}$	672	$1.114 \times 10^{-5}$
BRKMCC	2	$1 \times 10^{-5}$	644	$3.927 \times 10^{-7}$	566	$1.244 \times 10^{-5}$	661	$3.079 \times 10^{-5}$
BROWNAL.	100	$1 \times 10^{-5}$	2513	$1.963 \times 10^{-5}$	9567	$1.629 \times 10^{-5}$	4851	$1.487 \times 10^{-5}$
BROWNDEN	4	$1 \times 10^{-5}$	656	$3.971 \times 10^{-6}$	1145	$1.020 \times 10^{-5}$ 1.268 × 10 <sup>-5</sup>	1324	$1.107 \times 10^{-5}$ $1.062 \times 10^{-5}$
CLIFF	2	$1 \times 10^{-5}$	644	$2.902 \times 10^2$	1349	$2.275 \times 10^{-4}$	1008	$2.238 \times 10^{-4}$
CRACCIVY	100	$1 \times 10^{-5}$	7003	$8.012 \times 10^{-2}$	13711	$2.270 \times 10$ 8.641 × 10 <sup>-3</sup>	10035	$1.884 \times 10^{-2}$
CURE	100	$1 \times 10^{-5}$	644	$4.069 \times 10^{-2}$	700	$1.148 \times 10^{-2}$	662	$4.524 \times 10^{-3}$
DENGGRND	2	$1 \times 10^{-5}$	0044	$1.000 \times 10^{-1}$ $1.720 \times 10^{-1}$	036	$1.140 \times 10^{-5}$	1016	$2.024 \times 10^{-5}$
DENSCHID	2	$1 \times 10^{-5}$	990 650	$1.720 \times 10^{-1}$	562	$1.141 \times 10$ $0.002 \times 10^{-1}$	672	$2.311 \times 10^{-1}$
	300 200	$1 \times 10$ $1 \times 10^{-5}$	10244	$3.335 \times 10^{-2}$	28056	$5.335 \times 10^{-3}$	15709	$3.333 \times 10^{-3}$
DODTTO	100	$1 \times 10$ $1 \times 10^{-5}$	2400	$1.055 \times 10$ $2.687 \times 10^{-1}$	16494	$1.667 \times 10^{-4}$	1/100	$1.864 \times 10^{-4}$
DAVITC	100 96	$1 \times 10^{-5}$ $1 \times 10^{-5}$	3480 1104	$4.007 \times 10^{-4}$	10434	$1.007 \times 10^{-1}$	14109	$1.004 \times 10^{-4}$
EDENSCH	110	$1 \times 10^{-5}$	1194	$4.513 \times 10$ 1.020 × 10 <sup>-2</sup>	4840	$3.302 \times 10^{-2}$	4024	$1.831 \times 10$ $1.252 \times 10^{-2}$
EIGENALS	110	$1 \times 10^{-5}$	4921	1.929 × 10	11960	$1.337 \times 10$	10030	$1.333 \times 10$ $1.010 \times 10^{-2}$
EIGENBLS	110	$1 \times 10^{-5}$	3178	1.550	11308	1.553	46254	$1.018 \times 10^{-4}$
EIGENCLS	30	$1 \times 10^{-5}$	4370	$5.560 \times 10^{-4}$	9380	$8.302 \times 10^{-4}$	0373	$5.415 \times 10^{-4}$
ENGVAL1	100	$1 \times 10^{-5}$	2601	$2.285 \times 10^{-5}$	8619	$2.302 \times 10^{-6}$	(134	$2.310 \times 10^{-6}$
EXPETT	2	$1 \times 10^{-5}$	644	$4.595 \times 10^{-5}$	563	$6.982 \times 10^{-5}$	666	$6.814 \times 10^{-3}$
FLETCBV3	100	$1 \times 10^{-5}$	4326	$1.785 \times 10^{\circ}$	44623	$1.785 \times 10^{\circ}$	100756	$-1.304 \times 10^{-1}$
FLETCHBV	100	$1 \times 10^{-5}$	113423	$3.562 \times 10^{\circ}$	315886	$5.660 \times 10^{-4}$	213807	$1.765 \times 10^{\circ}$
FREUROTH	100	$1 \times 10^{-3}$	4326	$2.559 \times 10^{-3}$	10097	$5.248 \times 10^{-4}$	7166	$6.820 \times 10^{-4}$
GENROSE	100	$1 \times 10^{-3}$	27714	$8.446 \times 10^{-3}$	86415	$5.674 \times 10^{-3}$	56271	$1.067 \times 10^{-2}$
GULF	3	$1 \times 10^{-3}$	650	$6.596 \times 10^{-3}$	1091	$2.275 \times 10^{-3}$	675	$3.771 \times 10^{-3}$
HAIRY	2	$1 \times 10^{-3}$	644	$1.129 \times 10^{-4}$	1328	$4.688 \times 10^{-6}$	1008	$8.137 \times 10^{-6}$
HELIX	3	$1 \times 10^{-3}$	650	$7.422 \times 10^{-4}$	1385	$8.376 \times 10^{-3}$	650	$4.881 \times 10^{-3}$
NCB20B	100	$1 \times 10^{-3}$	4790	$4.231 \times 10^{-4}$	13270	$3.565 \times 10^{-4}$	9691	$3.061 \times 10^{-4}$
NONDIA	100	$1 \times 10^{-5}$	4326	$2.174 \times 10^{-4}$	21377	$1.572 \times 10^{-2}$	8423	$1.396 \times 10^{-2}$
NONDQUAR	100	$1 \times 10^{-5}$	11923	$9.225 \times 10^{-3}$	33549	$1.510 \times 10^{-3}$	14104	$1.836 \times 10^{-3}$
OSBORNEA	5	$1 \times 10^{-5}$	853	$7.139 \times 10^{-4}$	62478	$5.603 \times 10^{-1}$	1334	$1.215 \times 10^{-3}$
OSBORNEB	11	$1 \times 10^{-5}$	1137	$1.051 \times 10^{-1}$	4506	$2.381 \times 10^{-3}$	2936	$3.325 \times 10^{-3}$
PENALTY1	100	$1 \times 10^{-5}$	4326	$1.113 \times 10^{-1}$	20475	$1.893 \times 10^{-4}$	12804	$1.877 \times 10^{-4}$
PFIT1LS	3	$1 \times 10^{-5}$	17338	$1.609 \times 10^{-2}$	2062	$2.921 \times 10^{-6}$	1312	$2.320 \times 10^{-5}$
PFIT2LS	3	$1 \times 10^{-5}$	18947	$3.652 \times 10^{-2}$	2062	$3.200 \times 10^{-3}$	1310	$2.463 \times 10^{-3}$
PFIT3LS	3	$1 \times 10^{-5}$	21038	$2.034 \times 10^{-1}$	2062	$2.364 \times 10^{-2}$	1999	$2.822 \times 10^{-2}$
PFIT4LS	3	$1 \times 10^{-5}$	22064	$2.451 \times 10^{-1}$	2631	$1.004 \times 10^{-1}$	1990	$1.177 \times 10^{-1}$
QUARTC	100	$1 \times 10^{-5}$	3486	$2.687 \times 10^{-1}$	16434	$1.667 \times 10^{-4}$	14109	$1.864 \times 10^{-4}$
SINEVAL	2	$1 \times 10^{-5}$	625	$4.640 \times 10^{-3}$	2026	$3.676 \times 10^{-3}$	1977	$1.667 \times 10^{-4}$
SINQUAD	100	$1 \times 10^{-5}$	2613	$1.417\times 10^{-3}$	11616	$1.581\times10^{-4}$	5545	$2.677\times 10^{-4}$
SISSER	2	$1 \times 10^{-5}$	644	$7.894\times10^{-6}$	709	$2.995\times10^{-7}$	658	$3.218\times 10^{-6}$
SPARSQUR	100	$1 \times 10^{-5}$	8365	$9.873\times10^{-4}$	14478	$7.713\times10^{-5}$	8320	$3.739\times10^{-5}$
TOINTGSS	100	$1 \times 10^{-5}$	2513	$2.193\times10^{-3}$	10274	$6.054\times10^{-4}$	3729	$1.167\times 10^{-3}$
TQUARTIC	100	$1 \times 10^{-5}$	4571	$2.622\times10^{-1}$	9184	$1.534 \times 10^{-1}$	10072	$1.199\times10^{-1}$
TRIDIA	100	$1 \times 10^{-5}$	9428	$1.900\times 10^{-3}$	45233	$3.442 \times 10^{-3}$	19815	$2.348\times10^{-3}$
WATSON	31	$1 \times 10^{-5}$	3072	$3.088 \times 10^{-3}$	14287	$8.549 \times 10^{-3}$	6393	$6.386 \times 10^{-3}$
WOODS	100	$1 \times 10^{-5}$	7964	$8.185 \times 10^{-3}$	25179	$5.822 \times 10^{-3}$	13779	$3.840 \times 10^{-3}$
ZANGWIL2	2	$1 \times 10^{-5}$	644	$8.737 \times 10^{-6}$	670	$1.685 \times 10^{-5}$	657	$7.597 \times 10^{-6}$

Table B.14. Total number of function evaluations used and final accuracy achieved by forward-difference L-BFGS method with different choices of the finite-difference interval.

			<b>F:</b>	dIntomal	ъл	oné Wild	Α	dantivo
Problem	m	<i>c</i> -	#Evole	$\phi(x) = \phi^*$	#Evole	$\phi(x) = \phi^*$	A #Evole	$\phi(x) = \phi^*$
	<i>n</i>	$\frac{\epsilon_f}{1 \times 10^{-7}}$	#15vais	$\frac{\varphi(x) - \varphi}{0.768 \times 10^{-5}}$	#15vais	$\frac{\psi(x) - \phi}{1.100 \times 10^{-5}}$	#15vais	$\frac{\psi(x) - \psi}{1.001 \times 10^{-5}}$
ALLINTTU	⊿	$1 \times 10^{-7}$	656	$2.172 \times 10^{-7}$	2900 706	$1.133 \times 10^{-7}$ 1.061 × 10 <sup>-7</sup>	1004 543	$7.839 \times 10^{-8}$
ARWHEAD	100	$1 \times 10^{-7}$	4326	$1.643 \times 10^{-6}$	11900	$5.117 \times 10^{-7}$	5842	$1.039 \times 10^{-6}$
RARD	100	$1 \times 10^{-7}$	4520	$1.043 \times 10^{-5}$ 2 376 × 10 <sup>-5</sup>	706	$7.175 \times 10^{-7}$	660	$1.043 \times 10^{-7}$ $1.273 \times 10^{-7}$
BDORTIC	100	$1 \times 10^{-7}$	8011	$1.802 \times 10^{-5}$	1/968	$5.277 \times 10^{-6}$	12535	$4.346 \times 10^{-6}$
BIGGSS	100	$1 \times 10^{-7}$	650	$1.302 \times 10^{-7}$ 2.348 × 10 <sup>-7</sup>	726	$3.277 \times 10^{-7}$	1018	$1.386 \times 10^{-6}$
BIGGS5	5	$1 \times 10^{-7}$	1008	$6.378 \times 10^{-6}$	1614	$1.637 \times 10^{-6}$	2860	$3.801 \times 10^{-5}$
BIGGS6	6	$1 \times 10^{-7}$	1014	$-5.616 \times 10^{-3}$	2724	$-5.595 \times 10^{-3}$	2003	$-5.647 \times 10^{-3}$
BUAS	2	$1 \times 10^{-7}$	644	$-5.010 \times 10^{-6}$ 1 726 × 10 <sup>-6</sup>	690	$-5.535 \times 10^{-7}$ 1 547 $\times 10^{-7}$	658	$-5.047 \times 10$ 1 865 × 10 <sup>-6</sup>
BOX3	2	$1 \times 10^{-7}$	650	$7.011 \times 10^{-7}$	602	$6.639 \times 10^{-7}$	862	$5.213 \times 10^{-7}$
BBKMCC	2	$1 \times 10^{-7}$	644	$2.465 \times 10^{-7}$	670	$4.240 \times 10^{-7}$	661	$9.632 \times 10^{-7}$
BROWNAT	100	$1 \times 10^{-7}$	4326	$2.403 \times 10^{-7}$ $2.571 \times 10^{-7}$	0567	$4.240 \times 10$ $1.334 \times 10^{-5}$	7250	$5.032 \times 10^{-7}$
BROWNAL	100	$1 \times 10^{-7}$	4520	$2.571 \times 10^{-7}$ $3.507 \times 10^{-7}$	1145	$1.354 \times 10^{-7}$ $1.214 \times 10^{-7}$	1324	$1.468 \times 10^{-7}$
CITEE	-4 -9	$1 \times 10^{-7}$	644	$2.007 \times 10^{2}$	861	$1.214 \times 10$ $2.108 \times 10^{-4}$	1324	$1.400 \times 10$ $2.102 \times 10^{-4}$
CRACCIVY	100	$1 \times 10^{-7}$	4678	$1.570 \times 10^{-2}$	18164	$1.153 \times 10^{-4}$	14194	$7.253 \times 10^{-5}$
CURE	100	$1 \times 10^{-7}$	4070	$1.570 \times 10^{-4}$ $1.708 \times 10^{-4}$	13/0	$1.103 \times 10$ $1.188 \times 10^{-4}$	662	$1.203 \times 10^{-4}$
DENSCHND	2	$1 \times 10^{-7}$	650	$1.750 \times 10^{-2}$ $3.018 \times 10^{-2}$	1345	$3.805 \times 10^{-8}$	1311	$1.404 \times 10$ $1.200 \times 10^{-7}$
DENSCHNE	2	$1 \times 10^{-7}$	650	$0.003 \times 10^{-1}$	1001	$1.427 \times 10^{-7}$	1015	$1.203 \times 10^{-8}$ 8.808 × 10 <sup>-8</sup>
DIYMAANU	300	$1 \times 10^{-7}$	62344	$3.333 \times 10^{-4}$ 1 773 × 10 <sup>-4</sup>	63812	$1.427 \times 10$ $2.104 \times 10^{-5}$	37104	$2.753 \times 10^{-5}$
DORTIC	100	$1 \times 10^{-7}$	4571	$5.110 \times 10^{-3}$	20136	$2.194 \times 10$ $3.500 \times 10^{-7}$	17191	$4.377 \times 10^{-6}$
EDENSCH	36	$1 \times 10^{-7}$	2165	$2.262 \times 10^{-6}$	20130	$3.309 \times 10^{-6}$ $3.776 \times 10^{-6}$	4624	$4.577 \times 10$ 3 756 × 10 <sup>-6</sup>
EDENSON	110	$1 \times 10^{-7}$	2100	$1.202 \times 10^{-3}$	132863	$3.770 \times 10$ $2.204 \times 10^{-4}$	4024	$3.750 \times 10^{-4}$ $2.803 \times 10^{-4}$
EIGENRES	110	$1 \times 10^{-7}$	28055	$1.291 \times 10^{-3}$ $1.160 \times 10^{-3}$	182840	$2.204 \times 10^{-4}$ $0.604 \times 10^{-4}$	02870	$2.633 \times 10^{-4}$ $0.672 \times 10^{-4}$
EIGENDLS	30	$1 \times 10^{-7}$	3006	$1.103 \times 10^{-5}$ $1.173 \times 10^{-5}$	12/08	$9.004 \times 10$ $0.130 \times 10^{-6}$	6345	$3.072 \times 10^{-6}$ 8.630 × 10 <sup>-6</sup>
EIGENCLS ENCUAL 1	100	$1 \times 10^{-7}$	4326	$1.173 \times 10^{-6}$ $1.734 \times 10^{-6}$	8007	$3.133 \times 10^{-6}$ $3.737 \times 10^{-6}$	6527	$2.550 \times 10^{-6}$
ENGVALI	100	$1 \times 10^{-7}$	4520	$1.754 \times 10$ $3.707 \times 10^{-7}$	700	$5.737 \times 10^{-8}$ $5.737 \times 10^{-8}$	666	$4.705 \times 10^{-8}$
FIFTCBV3	100	$1 \times 10^{-7}$	30622	$4.451 \times 10^2$	281142	$1.666 \times 10^2$	208557	$-5.161 \times 10^{1}$
FIFTCHEV	100	$1 \times 10^{-7}$	116172	$-1.303 \times 10^9$	311603	$2.365 \times 10^9$	212007	$5.101 \times 10^{9}$ $5.073 \times 10^{9}$
FREIROTH	100	$1 \times 10^{-7}$	4678	$2.155 \times 10^{-5}$	10097	$2.500 \times 10^{-6}$ 8.099 × 10 <sup>-6</sup>	7165	$6.707 \times 10^{-6}$
GENROSE	100	$1 \times 10^{-7}$	27731	$1.016 \times 10^{-4}$	83521	$6.033 \times 10^{-5}$	56975	$8.924 \times 10^{-5}$
GILF	3	$1 \times 10^{-7}$	650	$4.478 \times 10^{-3}$	1163	$3.909 \times 10^{-3}$	1992	$1.893 \times 10^{-5}$
HATRY	2	$1 \times 10^{-7}$	644	$1.478 \times 10^{-6}$	670	$1.107 \times 10^{-7}$	1008	$2.080 \times 10^{-7}$
HELTX	3	$1 \times 10^{-7}$	650	$7.990 \times 10^{-6}$	1091	$1.355 \times 10^{-5}$	674	$1.967 \times 10^{-6}$
NCB20B	100	$1 \times 10^{-7}$	12717	$2.174 \times 10^{-5}$	34691	$2.432 \times 10^{-5}$	21290	$2.303 \times 10^{-5}$
NONDIA	100	$1 \times 10^{-7}$	5779	$2.111 \times 10^{-6}$ 2.912 × 10 <sup>-6</sup>	24808	$1.134 \times 10^{-5}$	5826	$8.625 \times 10^{-4}$
NONDOLLAR	100	$1 \times 10^{-7}$	8/38	$1.256 \times 10^{-3}$	102555	$1.104 \times 10$ $1.348 \times 10^{-4}$	44583	$1.100 \times 10^{-4}$
OSBORNEA	5	$1 \times 10^{-7}$	662	$1.200 \times 10^{-5}$ 8 747 × 10 <sup>-5</sup>	47678	$1.340 \times 10^{-1}$ $1.721 \times 10^{-1}$	2014	$2.200 \times 10^{-5}$
OSBORNER	11	$1 \times 10^{-7}$	2015	$1.342 \times 10^{-5}$	4/08	$1.721 \times 10$ $1.260 \times 10^{-5}$	3626	$2.250 \times 10^{-5}$ 2.063 × 10 <sup>-5</sup>
DENALTY1	100	$1 \times 10^{-7}$	4326	$6.722 \times 10^{-4}$	15820	$1.200 \times 10^{-4}$ $1.860 \times 10^{-4}$	1308/	$1.870 \times 10^{-4}$
DETT1IS	100	$1 \times 10^{-7}$	15/100	$5.585 \times 10^{-4}$	2062	$1.000 \times 10^{-6}$ $1.250 \times 10^{-6}$	1021	$7.981 \times 10^{-6}$
PETTOIS	3	$1 \times 10^{-7}$	18475	$3.104 \times 10^{-2}$	2002	$1.203 \times 10^{-3}$ $1.834 \times 10^{-3}$	1314	$1.001 \times 10^{-3}$
PETTRIS	3	$1 \times 10^{-7}$	10410	$4.062 \times 10^{-2}$	2062	$2.782 \times 10^{-2}$	1080	$3.161 \times 10^{-2}$
DETTAIS	2	$1 \times 10^{-7}$	22220	$4.002 \times 10^{-1}$ 2.817 $\times 10^{-1}$	2631	$1.171 \times 10^{-1}$	1909	$1.668 \times 10^{-1}$
OUARTC	100	$1 \times 10^{-7}$	4571	$5.110 \times 10^{-3}$	2001	$3509 \times 10^{-7}$	17191	$4.377 \times 10^{-6}$
SINEVAL	200	$1 \times 10^{-7}$	625	$7.770 \times 10^{-4}$	20100	$9.603 \times 10^{-7}$	1300	$5.182 \times 10^{-4}$
SINEVAL	100	$1 \times 10^{-7}$	3400	$1.620 \times 10^{-5}$	11616	$5.023 \times 10^{-6}$	5543	$6.333 \times 10^{-6}$
SINGORD	200	$1 \times 10^{-7}$	644	$2.627 \times 10^{-7}$	700	$1.365 \times 10^{-9}$	661	$6.235 \times 10^{-8}$
STREET	100	$1 \times 10^{-7}$	/861	$2.067 \times 10^{-5}$	26138	$1.300 \times 10^{-7}$	12562	$2.200 \times 10^{-7}$
TOINTGO	100	$1 \times 10^{-7}$	2512	$2.000 \times 10^{-5}$ 2.102 $\times 10^{-5}$	10651	$6.014 \times 10^{-6}$	3720	$1.022 \times 10^{-5}$
TUINDALG	100	$1 \times 10^{-7}$	2010 8805	$1.645 \times 10^{-2}$	1/091	$8.050 \times 10^{-4}$	10069	$1.100 \times 10^{-3}$
TRIDIA	100	$1 \times 10$ $1 \times 10^{-7}$	10610	$1.040 \times 10^{-5}$ $2.400 \times 10^{-5}$	14021	$2.000 \times 10^{-5}$	210/2	$3.486 \times 10^{-5}$
MATCOM	21	$1 \times 10$ $1 \times 10^{-7}$	10019	$2.455 \times 10^{-3}$ 1 161 $\sim 10^{-3}$	12659	$2.035 \times 10$ 1 179 $\sim 10^{-3}$	4595 4595	$1.854 \sim 10^{-3}$
WOODG	100	$1 \times 10^{-7}$	6600	$5.611 \times 10^{-5}$	25170	$4.354 \times 10^{-5}$	10111	$9.748 \times 10^{-5}$
ZANGWII 2	200	$1 \times 10^{-7}$	644	$8.738 \times 10^{-8}$	670	$1.683 \times 10^{-7}$	657	$7.832 \times 10^{-8}$
2011 Q W 1 L L Z	4	T \\ TO	011	0.100 / 10	010	T:000 V TO	001	1.005 / 10

Table B.15. Total number of function evaluations used and final accuracy achieved by central-difference L-BFGS method with different choices of the finite-difference interval.

			Five	d Interval		Adaptive
Problem	n	E s	#Evals	$\phi(x) = \phi^*$	#Evals	$\phi(x) = \phi^*$
ATRCRETE	5	$\frac{c_f}{1 \times 10^{-1}}$	#12 (all)	$\frac{\psi(w)}{4455 \times 10^{-1}}$	3074	$\frac{\psi(x)}{4451 \times 10^{-1}}$
ALLINTTU	4	$1 \times 10^{-1}$	848	$7.698 \times 10^{-1}$	1371	$3.378 \times 10^{-3}$
ABWHFAD	100	$1 \times 10^{-1}$	4001	$4.137 \times 10^{-2}$	5021	$5.576 \times 10^{-2}$
BARD	3	$1 \times 10^{-1}$	660	$2.616 \times 10^{-3}$	0/7	$1.535 \times 10^{-2}$
BDORTIC	100	$1 \times 10$ $1 \times 10^{-1}$	8504	$2.010 \times 10$ $2.047 \times 10^{-1}$	11555	$1.035 \times 10^{-1}$ $1.018 \times 10^{-1}$
DEGICITO	100	$1 \times 10^{-1}$	660	$2.047 \times 10$ $2.607 \times 10^{-1}$	1150	$1.510 \times 10^{-2}$ $1.154 \times 10^{-2}$
DIGG55	5	$1 \times 10^{-1}$	603	$2.007 \times 10$ 8 820 × 10 <sup>-2</sup>	1195	$1.104 \times 10$ $3.740 \times 10^{-2}$
DIGGSS	6	$1 \times 10^{-1}$	705	$2.629 \times 10^{-1}$	1/10	$3.140 \times 10^{-1}$
DIGGOO	0 9	$1 \times 10^{-1}$	657	$2.300 \times 10$ 1.012 × 10 <sup>-2</sup>	560	$2.307 \times 10^{-3}$
DUAZ	2	$1 \times 10$ $1 \times 10^{-1}$	660	$1.013 \times 10$ $4.700 \times 10^{-2}$	500	$0.435 \times 10^{-3}$
BUAS	ა ი	$1 \times 10$ $1 \times 10^{-1}$	009 657	$4.700 \times 10$ $4.767 \times 10^{-4}$	094	$1.444 \times 10$ $1.082 \times 10^{-4}$
BRAMCC	100	$1 \times 10$ $1 \times 10^{-1}$	007 7175	$4.707 \times 10^{-5}$	2410	$1.965 \times 10$ $1.221 \times 10^{-7}$
BRUWNAL	100	$1 \times 10$ $1 \times 10^{-1}$	(1/0	$1.321 \times 10^{-4}$	12140	$1.331 \times 10^{-3}$
BRUWNDEN	4	$1 \times 10^{-1}$	681	$4.778 \times 10^{-10}$	2247	$1.814 \times 10^{-2}$
CLIFF	2	$1 \times 10^{-1}$	657	$2.902 \times 10^{2}$	4819	$2.413 \times 10^{-2}$
CRAGGLVY	100	$1 \times 10^{-1}$	11144	2.765	23987	$6.089 \times 10^{-2}$
CUBE	2	$1 \times 10^{-1}$	657	$4.421 \times 10^{-2}$	3101	$4.462 \times 10^{-2}$
DENSCHND	3	$1 \times 10^{-1}$	669	$5.418 \times 10^{-3}$	2310	$8.228 \times 10^{-3}$
DENSCHNE	3	$1 \times 10^{-1}$	669	1.015	576	1.025
DIXMAANH	300	$1 \times 10^{-1}$	32019	$7.301 \times 10^{-1}$	41575	$4.225 \times 10^{-2}$
DQRTIC	100	$1 \times 10^{-1}$	10891	$4.910 \times 10^{-1}$	28135	$7.032 \times 10^{-4}$
EDENSCH	36	$1 \times 10^{-1}$	4217	$5.638 \times 10^{-2}$	6835	$7.687 \times 10^{-2}$
EIGENALS	110	$1 \times 10^{-1}$	7295	2.049	20145	$1.099 \times 10^{-1}$
EIGENBLS	110	$1 \times 10^{-1}$	12096	2.326	13672	1.801
EIGENCLS	30	$1 \times 10^{-1}$	3161	$6.304 \times 10^{-1}$	8124	$3.930 \times 10^{-1}$
ENGVAL1	100	$1 \times 10^{-1}$	6704	$1.473 \times 10^{-1}$	11258	$2.414 \times 10^{-1}$
EXPFIT	2	$1 \times 10^{-1}$	7418	$2.456 \times 10^{1}$	5915	$2.093 \times 10^{-2}$
<b>FLETCBV3</b>	100	$1 \times 10^{-1}$	2546	$1.785 \times 10^{5}$	136921	$-1.561 \times 10^{2}$
FLETCHBV	100	$1 \times 10^{-1}$	138282	$6.330 \times 10^{9}$	380580	$-1.079 \times 10^{9}$
FREUROTH	100	$1 \times 10^{-1}$	6216	$2.650 \times 10^{-1}$	14495	$4.930 \times 10^{-2}$
GENROSE	100	$1 \times 10^{-1}$	7084	$1.116 \times 10^2$	28786	$1.117 \times 10^2$
GULF	3	$1 \times 10^{-1}$	669	6.621	2124	6.728
HAIRY	2	$1 \times 10^{-1}$	657	$1.261 \times 10^{-4}$	1483	$6.896 \times 10^{-4}$
HELIX	3	$1 \times 10^{-1}$	505	$8.544 \times 10^{2}$	4900	$2.476 \times 10^{1}$
NCB20B	100	$1 \times 10^{-1}$	3028	2.230	12309	$4.906 \times 10^{-2}$
NONDIA	100	$1 \times 10^{-1}$	4782	$4.941 \times 10^{-1}$	9226	$4.929 \times 10^{-1}$
NONDQUAR	100	$1 \times 10^{-1}$	7210	$1.974 \times 10^{-1}$	14490	$7.337 \times 10^{-2}$
OSBORNEA	5	$1 \times 10^{-1}$	693	$1.530 \times 10^{-1}$	2345	$8.790 \times 10^{-1}$
OSBORNEB	11	$1 \times 10^{-1}$	746	$6.496 \times 10^{-1}$	3119	1.366
PENALTY1	100	$1 \times 10^{-1}$	7210	$3.404 \times 10^{-2}$	21785	$1.847 \times 10^{-4}$
PFTT1LS	3	$1 \times 10^{-1}$	17843	$1.276 \times 10^{1}$	3706	$8.023 \times 10^{-1}$
PFIT2LS	3	$1 \times 10^{-1}$	3466	$1.270 \times 10^{2}$ $1.453 \times 10^{2}$	2827	$5.360 \times 10^{-1}$
PETTALS	3	$1 \times 10^{-1}$	26780	$8.283 \times 10^2$	3807	$1.767 \times 10^{-1}$
PFITAI S	3	$1 \times 10^{-1}$	27602	$2.416 \times 10^3$	3727	$3.979 \times 10^{-1}$
OUARTC	100	$1 \times 10$ $1 \times 10^{-1}$	10801	$4.910 \times 10^{-1}$	28135	$5.575 \times 10^{-4}$ $7.032 \times 10^{-4}$
QUARTO C	100	$1 \times 10^{-1}$	501	4.310 × 10	20100	$7.052 \times 10$ 7.018
SINGVAL	100	$1 \times 10^{-1}$	501	0.047 0.047 0.047	11990	2 802
SINGOAD	100	$1 \times 10^{-1}$	657	$2.813 \times 10^{-3}$	401	$2.000 \times 10^{-5}$
STOPEC	4 100	$1 \times 10$ $1 \times 10^{-1}$	E945	$9.207 \times 10$ 6 181 × 10 <sup>-2</sup>	19405	$2.900 \times 10^{-2}$
SPARSUUK	100	$1 \times 10^{-1}$	0340 E 201	$0.101 \times 10^{-2}$	16490	$1.010 \times 10^{-9}$
TUINTGSS	100	$1 \times 10^{-1}$	0391	$1.021 \times 10^{-1}$	8324	$4.000 \times 10^{-1}$
TQUARTIC	100	$1 \times 10^{-1}$	2605	$0.330 \times 10^{-1}$	12297	$(.320 \times 10^{-13})$
TRIDIA	100	$1 \times 10^{-1}$	10830	$1.095 \times 10^{-1}$	88893	$5.839 \times 10^{-10}$
WATSON	31	$1 \times 10^{-1}$	2880	$2.177 \times 10^{-11}$	5954	$1.693 \times 10^{-1}$
WOODS	100	$1 \times 10^{-1}$	5279	6.439	21136	5.545
ZANGWIL2	2	$1 \times 10^{-1}$	657	$4.902 \times 10^{-4}$	1351	$-9.999 \times 10^{-11}$

Table B.16. Total number of function evaluations used and final accuracy achieved by central-difference L-BFGS method with different choices of the finite-difference interval.

Fixed Interval Ada	ptive
Problem $n  \epsilon_f  \# \text{Evals}  \phi(x) - \phi^*  \# \text{Evals}$	$\phi(x) - \phi^*$
AIRCRFTB 5 $1 \times 10^{-3}$ 594 $4.451 \times 10^{-1}$ 2306	$4.451 \times 10^{-1}$
ALLINITU 4 $1 \times 10^{-3}$ 574 $3.325 \times 10^{-3}$ 1377	$1.847 \times 10^{-5}$
ARWHEAD $100 \ 1 \times 10^{-3}$ 7175 $5.590 \times 10^{-4}$ 8379	$6.125 \times 10^{-4}$
BARD 3 $1 \times 10^{-3}$ 669 $1.905 \times 10^{-3}$ 713	$1.883 \times 10^{-3}$
<b>BDQRTIC</b> 100 $1 \times 10^{-3}$ 11220 $3.667 \times 10^{-3}$ 31078	$4.314 \times 10^{-4}$
<b>BIGGS3</b> 3 $1 \times 10^{-3}$ 669 $1.009 \times 10^{-3}$ 1363	$1.314 \times 10^{-4}$
<b>BIGGS5</b> 5 $1 \times 10^{-3}$ 1333 $1.649 \times 10^{-2}$ 1921	$1.418 \times 10^{-2}$
<b>BIGGS6</b> 6 $1 \times 10^{-3}$ 705 $5.181 \times 10^{-2}$ 1844 -	$3.749 \times 10^{-3}$
<b>BOX2</b> 2 $1 \times 10^{-3}$ 657 $1.287 \times 10^{-5}$ 560	$1.740 \times 10^{-4}$
<b>BOX3</b> 3 $1 \times 10^{-3}$ 669 $3.555 \times 10^{-5}$ 1351	$5.017 \times 10^{-5}$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$3.658 \times 10^{-7}$
<b>BROWNAL</b> 100 $1 \times 10^{-3}$ 4782 $8.022 \times 10^{-6}$ 12030	$2.446 \times 10^{-8}$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$1.803 \times 10^{-7}$
CLIEF $2.1 \times 10^{-3}$ $657$ $2.002 \times 10^{2}$ $3500$	$2.407 \times 10^{-4}$
CELEFF 2 1 × 10 057 2.502 × 10 5550 CPACCUV 100 1 × 10 <sup>-3</sup> 11076 1.745 × 10 <sup>-2</sup> 20115	$2.437 \times 10^{-3}$
CINE 2 1 $\times$ 10 <sup>-3</sup> 657 4 220 $\times$ 10 <sup>-2</sup> 2418	$4.580 \times 10^{-2}$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$4.000 \times 10$ 6 470 × 10 <sup>-4</sup>
DENGCINE $2 + 1 \times 10^{-3}$ CC0 0.002 $\times 10^{-1}$ 1000	$0.472 \times 10^{-1}$
DENSCHINE 3 1 × 10 $009 9.993 \times 10^{-1} 1002$	$9.994 \times 10^{-3}$
<b>DIAMAANH</b> 300 $1 \times 10^{-5}$ 28496 $1.428 \times 10^{-5}$ 48610	$3.093 \times 10^{-5}$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$2.425 \times 10^{-4}$
EDENSCH 36 $1 \times 10^{-3}$ 2128 $2.700 \times 10^{-4}$ 4599	$2.101 \times 10^{-4}$
<b>EIGENALS</b> 110 $1 \times 10^{-3}$ 8516 $4.539 \times 10^{-2}$ 27468	$1.303 \times 10^{-2}$
<b>EIGENBLS</b> 110 $1 \times 10^{-3}$ 47792 $3.108 \times 10^{-2}$ 25441	1.551
<b>EIGENCLS</b> 30 $1 \times 10^{-3}$ 4087 $2.953 \times 10^{-3}$ 10419	$1.545 \times 10^{-3}$
<b>ENGVAL1</b> 100 $1 \times 10^{-3}$ 7210 $1.658 \times 10^{-4}$ 14954	$4.131 \times 10^{-4}$
<b>EXPFIT</b> 2 $1 \times 10^{-3}$ 9839 3.021 1013	$2.602 \times 10^{-4}$
FLETCBV3 100 $1 \times 10^{-3}$ 2546 $1.785 \times 10^{5}$ 198233	$-1.561 \times 10^{2}$
FLETCHBV 100 $1 \times 10^{-3}$ 107444 $-1.473 \times 10^{9}$ 239215	$3.847 \times 10^{9}$
FREUROTH 100 $1 \times 10^{-3}$ 7175 $4.683 \times 10^{-4}$ 14972	$1.047 \times 10^{-4}$
GENROSE 100 $1 \times 10^{-3}$ 54819 $2.202 \times 10^{-3}$ 124290	$9.891 \times 10^{-4}$
GULF $3 \ 1 \times 10^{-3}$ $1309 \ 2.995 \times 10^{-3}$ $1782$	$4.249 \times 10^{-3}$
HAIRY 2 $1 \times 10^{-3}$ 418 $5.279 \times 10^{-6}$ 4113	$1.608 \times 10^{-7}$
HELIX 3 $1 \times 10^{-3}$ 432 $8.522 \times 10^{2}$ 2867	$2.475 \times 10^{1}$
NCB20B 100 $1 \times 10^{-3}$ 5918 $1.353 \times 10^{-1}$ 14781	$1.608 \times 10^{-3}$
NONDIA 100 $1 \times 10^{-3}$ 6475 $1.027 \times 10^{-2}$ 10213	$1.263 \times 10^{-2}$
NONDQUAR 100 $1 \times 10^{-3}$ 7084 $1.973 \times 10^{-2}$ 28713	$3.712 \times 10^{-3}$
<b>OSBORNEA</b> 5 $1 \times 10^{-3}$ 1333 $2.881 \times 10^{-3}$ 2237	$8.790 \times 10^{-1}$
OSBORNEB 11 $1 \times 10^{-3}$ 2082 $2.869 \times 10^{-1}$ 3565	$7.271 \times 10^{-2}$
<b>PENALTY1</b> 100 $1 \times 10^{-3}$ 7210 $1.928 \times 10^{-4}$ 20312	$1.917 \times 10^{-4}$
<b>PETTILS</b> 3 $1 \times 10^{-3}$ 24764 $1.500 \times 10^{1}$ 2007	$1.183 \times 10^{-3}$
<b>PEIT2LS</b> 3 $1 \times 10^{-3}$ 28415 $1.073 \times 10^{2}$ 2021	$4.006 \times 10^{-3}$
<b>PETTALS</b> 3 1 × 10 <sup>-3</sup> 31502 1.018 × 10 <sup>3</sup> 2706	$3.066 \times 10^{-2}$
<b>DELEVANCE</b> 3 $1 \times 10^{-3}$ 31607 3 250 $\times 10^{3}$ 3001	$0.000 \times 10^{-2}$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$9.000 \times 10^{-5}$
<b>QUALITY 100</b> 1 $\wedge$ 10 11010 1.010 $\times$ 10 20201 <b>QUALITY 10</b> 2 1 $\vee$ 10 <sup>-3</sup> 500 5 207 2492	$2.420 \times 10$ $1.873 \times 10^{-5}$
<b>SINEVAL</b> 2 I X IU $088$ $0.327$ $3430$	4.0/0 × 10
<b>SINGURE</b> 100 I X 10 $0391$ 0.003 X 10 $14970$	$1.2(1 \times 10^{-5})$
<b>SISSER</b> 2 $1 \times 10^{\circ}$ 657 $9.026 \times 10^{\circ}$ 684	$5.497 \times 10^{\circ}$
<b>SPARSUUR</b> 100 $1 \times 10^{-5}$ 5542 $2.341 \times 10^{-5}$ 19979	$2.489 \times 10^{-3}$
TOINTGSS 100 $1 \times 10^{-3}$ 2896 $1.267 \times 10^{-5}$ 9505	$4.000 \times 10^{-9}$
TQUARTIC 100 $1 \times 10^{-3}$   19220 $2.492 \times 10^{-1}$   8491	$7.221 \times 10^{-1}$
<b>TRIDIA</b> 100 $1 \times 10^{-3}$   17514 $3.205 \times 10^{-5}$   92231 4	$1.283 \times 10^{-15}$
WATSON 31 $1 \times 10^{-3}$ 3173 $2.462 \times 10^{-2}$ 9834	$2.186 \times 10^{-3}$
WOODS 100 $1 \times 10^{-3}$   11220 $3.320 \times 10^{-3}$   23074	$6.287 \times 10^{-4}$
<b>ZANGWIL2</b> 2 $1 \times 10^{-3}$ 657 $9.871 \times 10^{-7}$ 598 -9	$0.999 \times 10^{-11}$

Table B.17. Total number of function evaluations used and final accuracy achieved by central-difference L-BFGS method with different choices of the finite-difference interval.

			<b>D</b> !	1 T., 4 1		
D 1 1			Fixe	d Interval	E I I I I I I I I I I I I I I I I I I I	Adaptive
Problem	<u>n</u>	$\epsilon_f$	#Evals	$\phi(x) - \phi^{\dagger}$	#Evals	$\phi(x) - \phi$
AIRCRFTB	5	$1 \times 10^{-5}$	612	$4.451 \times 10^{-4}$	2558	$3.293 \times 10^{-1}$
ALLINITU	4	$1 \times 10^{-3}$	681	$7.277 \times 10^{-6}$	2035	$7.767 \times 10^{-3}$
ARWHEAD	100	$1 \times 10^{-5}$	7487	$4.138 \times 10^{-6}$	7477	$1.108 \times 10^{-6}$
BARD	3	$1 \times 10^{-5}$	669	$1.867 \times 10^{-3}$	1500	$1.678 \times 10^{-7}$
BDQRTIC	100	$1 \times 10^{-5}$	11220	$1.539 \times 10^{-6}$	19106	$4.834 \times 10^{-6}$
BIGGS3	3	$1 \times 10^{-5}$	669	$4.019 \times 10^{-6}$	1088	$3.251 \times 10^{-7}$
BIGGS5	5	$1 \times 10^{-5}$	1756	$4.920 \times 10^{-5}$	3617	$1.465 \times 10^{-5}$
BIGGS6	6	$1 \times 10^{-5}$	2358	$-5.228 \times 10^{-3}$	4142	$-5.625 \times 10^{-3}$
BOX2	2	$1 \times 10^{-5}$	657	$2.137 \times 10^{-6}$	1327	$1.773 \times 10^{-6}$
BOX3	3	$1 \times 10^{-5}$	669	$6.746 \times 10^{-7}$	1349	$7.643 \times 10^{-7}$
BRKMCC	2	$1 \times 10^{-5}$	437	$5.850 \times 10^{-10}$	558	$1.300 \times 10^{-9}$
BROWNAL	100	$1 \times 10^{-5}$	3508	$1.384 \times 10^{-7}$	19759	$1.759 \times 10^{-16}$
BROWNDEN	4	$1 \times 10^{-5}$	681	$7.291 \times 10^{-9}$	2060	$-1.455 \times 10^{-11}$
CLIFF	2	$1 \times 10^{-5}$	657	$2.902 \times 10^{2}$	4442	$2.272 \times 10^{-4}$
CRAGGLVY	100	$1 \times 10^{-5}$	17145	$3.635 \times 10^{-5}$	32645	$5.568 \times 10^{-6}$
CUBE	2	$1 \times 10^{-5}$	657	$2.696 \times 10^{-7}$	2510	$3.013 \times 10^{-6}$
DENSCHND	3	$1 \times 10^{-5}$	1015	$3.165 \times 10^{-4}$	1878	$2.382 \times 10^{-6}$
DENSCHNE	3	$1 \times 10^{-5}$	669	$9.903 \times 10^{-1}$	1073	$7.559 \times 10^{-8}$
DIYMAANH	300	$1 \times 10^{-5}$	45747	$9.550 \times 10^{-5}$ 8.540 × 10 <sup>-5</sup>	225222	$1.565 \times 10^{-10}$
DIAMAANII	100	$1 \times 10^{-5}$	19790	$1.501 \times 10^{-4}$	200002	$1.005 \times 10^{-6}$
EDENGQU	200	$1 \times 10^{-5}$	10100	$1.001 \times 10$ 2.028 × 10 <sup>-7</sup>	20195	$1.377 \times 10^{-7}$
EDENSCH	110	$1 \times 10^{-5}$	2382	$3.038 \times 10^{-3}$	0845	$3.233 \times 10^{-4}$
EIGENALS	110	$1 \times 10^{-5}$	27094	$1.330 \times 10^{-4}$	93131	$2.888 \times 10$ 1.114 × 10 <sup>-3</sup>
EIGENBLS	110	$1 \times 10^{-5}$	86034	$9.991 \times 10^{-1}$	167257	$1.114 \times 10^{-6}$
EIGENCLS	30	$1 \times 10^{-5}$	5297	$3.243 \times 10^{-6}$	15418	$5.660 \times 10^{-5}$
ENGVAL1	100	$1 \times 10^{-6}$	5244	$1.653 \times 10^{-6}$	12432	$7.819 \times 10^{-6}$
EXPFIT	2	$1 \times 10^{-3}$	2953	$1.015 \times 10^{-2}$	1097	$3.382 \times 10^{-7}$
FLETCBV3	100	$1 \times 10^{-5}$	61879	$-7.404 \times 10^{1}$	405665	$-1.561 \times 10^{2}$
FLETCHBV	100	$1 \times 10^{-5}$	69727	$1.786 \times 10^{9}$	168055	$-3.029 \times 10^{8}$
FREUROTH	100	$1 \times 10^{-5}$	6618	$1.671 \times 10^{-6}$	17360	$1.873 \times 10^{-7}$
GENROSE	100	$1 \times 10^{-5}$	53552	$5.526 \times 10^{-6}$	119843	$3.104 \times 10^{-6}$
GULF	3	$1 \times 10^{-5}$	669	$4.155 \times 10^{-3}$	2094	$1.869 \times 10^{-6}$
HAIRY	2	$1 \times 10^{-5}$	657	$3.882 \times 10^{-9}$	2015	$2.389 \times 10^{-11}$
HELIX	3	$1 \times 10^{-5}$	445	$8.521 \times 10^{2}$	5026	$2.475 \times 10^{1}$
NCB20B	100	$1 \times 10^{-5}$	14030	$1.578 \times 10^{-3}$	37332	$7.086 \times 10^{-5}$
NONDIA	100	$1 \times 10^{-5}$	7210	$1.054 \times 10^{-5}$	13740	$1.839 \times 10^{-5}$
NONDQUAR	100	$1 \times 10^{-5}$	18545	$9.490 \times 10^{-4}$	55137	$2.686 \times 10^{-4}$
OSBORNEA	5	$1 \times 10^{-5}$	693	$9.012 \times 10^{-4}$	2295	$8.790 \times 10^{-1}$
OSBORNEB	11	$1 \times 10^{-5}$	3861	$4.164 \times 10^{-4}$	5596	$2.698 \times 10^{-5}$
PENALTY1	100	$1 \times 10^{-5}$	9496	$1.868 \times 10^{-4}$	22443	$1.875 \times 10^{-4}$
PFTT1LS	3	$1 \times 10^{-5}$	26200	5.375	2883	$8.388 \times 10^{-6}$
PETT2LS	3	$1 \times 10^{-5}$	29030	$1.010 \times 10^2$	2008	$2.499 \times 10^{-3}$
PETTSIS	3	$1 \times 10^{-5}$	30940	$6.034 \times 10^2$	2330	$4.427 \times 10^{-2}$
DEITAIS	3	$1 \times 10^{-5}$	31645	$2.140 \times 10^3$	6037	$7.936 \times 10^{-2}$
OTINETC	100	$1 \times 10^{-5}$	13738	$1.501 \times 10^{-4}$	28703	$1.930 \times 10^{-6}$
QUARIC CINEVAL	100	$1 \land 10$ $1 \lor 10^{-5}$	10100	1.001 X 10 2.050	20193	$1.377 \times 10^{-7}$
SINEVAL	100	$1 \times 10^{-5}$	10907	3.900 1.979 · 10 <sup>-6</sup>	3089	$2.070 \times 10^{-7}$
SINUUAD	100	$1 \times 10^{-5}$	0127	$1.3/3 \times 10^{\circ}$	10270	4.398 × 10 5.919 × 10-11
SISSER	2	$1 \times 10^{-5}$	657	$3.088 \times 10^{-5}$	10000	$0.212 \times 10^{-6}$
SPARSQUR	100	$1 \times 10^{-3}$	7464	$2.220 \times 10^{-3}$	19032	$5.543 \times 10^{-6}$
TOINTGSS	100	$1 \times 10^{-3}$	8221	$2.253 \times 10^{-8}$	8385	$4.000 \times 10^{-5}$
TQUARTIC	100	$1 \times 10^{-5}$	12669	$8.689 \times 10^{-4}$	15299	$3.231 \times 10^{-4}$
TRIDIA	100	$1 \times 10^{-5}$	20802	$5.305 \times 10^{-7}$	107193	$1.824 \times 10^{-17}$
WATSON	31	$1 \times 10^{-5}$	7718	$1.087 \times 10^{-3}$	16253	$1.209 \times 10^{-3}$
WOODS	100	$1 \times 10^{-5}$	9234	$2.150 \times 10^{-5}$	24007	$3.869 \times 10^{-6}$
ZANGWIL2	2	$1 \times 10^{-5}$	657	$2.019\times10^{-9}$	1347	$-1.000 \times 10^{-10}$

Table B.18. Total number of function evaluations used and final accuracy achieved by central-difference L-BFGS method with different choices of the finite-difference interval.

			<b>T</b> .	1 T / 1					
D 11			Fixe	ed Interval	Adaptive				
Problem	n	$\epsilon_f$	#Evals	$\phi(x) - \phi^*$	#Evals	$\phi(x) - \phi^*$			
AIRCRFTB	5	$1 \times 10^{-7}$	1585	$6.912 \times 10^{-12}$	2929	$2.059 \times 10^{-22}$			
ALLINITU	4	$1 \times 10^{-7}$	463	$1.644 \times 10^{-8}$	1493	$3.211 \times 10^{-10}$			
ARWHEAD	100	$1 \times 10^{-7}$	5391	$1.192 \times 10^{-8}$	8298	$2.636 \times 10^{-9}$			
BARD	3	$1 \times 10^{-7}$	669	$2.756 \times 10^{-9}$	1357	$4.989 \times 10^{-9}$			
BDQRTIC	100	$1 \times 10^{-7}$	11220	$1.885 \times 10^{-8}$	27324	$6.486 \times 10^{-9}$			
BIGGS3	3	$1 \times 10^{-7}$	669	$1.115 \times 10^{-8}$	1357	$5.031 \times 10^{-10}$			
BIGGS5	5	$1 \times 10^{-7}$	1530	$2.281 \times 10^{-8}$	3911	$4.135 \times 10^{-9}$			
BIGGS6	6	$1 \times 10^{-7}$	2022	$-5.642 \times 10^{-3}$	5582	$-5.649 \times 10^{-3}$			
BOX2	2	$1 \times 10^{-7}$	657	$3.579 \times 10^{-9}$	678	$1.419 \times 10^{-9}$			
BOX3	3	$1 \times 10^{-7}$	669	$6.337 \times 10^{-7}$	1092	$6.848 \times 10^{-10}$			
BRKMCC	2	$1 \times 10^{-7}$	657	$2.043 \times 10^{-10}$	680	$2.005 \times 10^{-10}$			
BROWNAL	100	$1 \times 10^{-7}$	3799	$2.783 \times 10^{-8}$	24944	$6.904 \times 10^{-19}$			
BROWNDEN	4	$1 \times 10^{-7}$	681	$-3.638 \times 10^{-10}$	1695	$-3.929 \times 10^{-10}$			
CLIFF	2	$1 \times 10^{-7}$	657	$5.621 \times 10^{-1}$	10094	$2.902 \times 10^2$			
CRAGGLVY	100	$1 \times 10^{-7}$	18747	$5.581 \times 10^{-7}$	39527	$3.028 \times 10^{-8}$			
CUBE	2	$1 \times 10^{-7}$	657	$1.452 \times 10^{-7}$	1905	$7.199 \times 10^{-9}$			
DENSCHND	3	$1 \times 10^{-7}$	1309	$1.440 \times 10^{-5}$	2057	$8.977 \times 10^{-8}$			
DENSCHNE	3	$1 \times 10^{-7}$	669	$3.754 \times 10^{-11}$	2034	$2.051 \times 10^{-10}$			
DIXMAANH	300	$1 \times 10^{-7}$	52791	$4.060 \times 10^{-7}$	394134	0.000			
DORTIC	100	$1 \times 10^{-7}$	14367	$2.836 \times 10^{-6}$	29677	$2.399 \times 10^{-9}$			
EDENSCH	36	$1 \times 10^{-7}$	2434	$1.106 \times 10^{-9}$	6835	$8.345 \times 10^{-10}$			
EIGENALS	110	$1 \times 10^{-7}$	86646	$7.971 \times 10^{-6}$	194737	$2.282 \times 10^{-6}$			
FIGENBLS	110	$1 \times 10^{-7}$	194875	$1.071 \times 10^{-6}$ $1.129 \times 10^{-6}$	257660	$9.230 \times 10^{-4}$			
FIGENCIS	30	$1 \times 10^{-7}$	7754	$1.120 \times 10^{-7}$	17966	$7.350 \times 10^{-8}$			
EIGLIOLD	100	$1 \times 10^{-7}$	7210	$5.117 \times 10^{-9}$	13746	$7.056 \times 10^{-9}$			
ENGVALI	200	$1 \times 10^{-7}$	3032	$6.553 \times 10^{-5}$	13/40	$1.300 \times 10^{-10}$ 8 708 × 10 <sup>-10</sup>			
EXITIT FIFTCBV3	100	$1 \times 10^{-7}$	101241	$-2.737 \times 10^{1}$	603140	$-8.437 \times 10^{1}$			
FIETCHEV	100	$1 \times 10^{-7}$	106670	$-2.757 \times 10^{-1}$	410157	$-0.437 \times 10^{-1}$ $-1.207 \times 10^{9}$			
FREIDATU	100	$1 \times 10^{-7}$	6382	$3.050 \times 10^{-9}$	14775	$-1.257 \times 10^{-9}$ 5.215 $\times 10^{-9}$			
CENDORE	100	$1 \times 10^{-7}$	59767	$-3.030 \times 10^{-8}$	112006	$-0.210 \times 10^{-8}$			
CULE	100	$1 \times 10^{-7}$	1504	$2.210 \times 10^{-7}$	2000	$5.000 \times 10^{-8}$			
GOLF	ა ი	$1 \times 10^{-7}$	1504	$3.322 \times 10^{-12}$	2090	0.296 × 10			
HAIRI	2	$1 \times 10$ $1 \times 10^{-7}$	007	$1.725 \times 10^{-0.02}$	1011	0.000			
HELIX	100	$1 \times 10$ $1 \times 10^{-7}$	415	$0.021 \times 10^{-5}$	2820	$0.700 \times 10^{-5}$			
NCB20B	100	$1 \times 10^{-7}$	51509 CC19	$2.837 \times 10^{-8}$	13074	$1.049 \times 10^{-8}$			
NUNDIA	100	$1 \times 10^{-7}$	0018	$0.417 \times 10^{-5}$	14890	$0.210 \times 10^{-5}$			
NONDQUAR	100	$1 \times 10^{-7}$	77500	$5.554 \times 10^{-5}$	239505	$1.867 \times 10^{-3}$			
OSBORNEA	5	$1 \times 10^{-7}$	1333	$2.245 \times 10^{-8}$	2162	$1.562 \times 10^{-9}$			
OSBORNEB	11	$1 \times 10^{-7}$	2933	$1.612 \times 10^{-8}$	6229	$6.560 \times 10^{-5}$			
PENALTY1	100	$1 \times 10^{-7}$	8565	$1.868 \times 10^{-4}$	76226	$2.507 \times 10^{-6}$			
PFIT1LS	3	$1 \times 10^{-7}$	26224	$4.862 \times 10^{-4}$	4269	$2.294 \times 10^{-6}$			
PFIT2LS	3	$1 \times 10^{-7}$	28717	$1.295 \times 10^{-2}$	14332	$2.756 \times 10^{-5}$			
PFIT3LS	3	$1 \times 10^{-7}$	28777	$5.848 \times 10^{-2}$	22899	$1.676 \times 10^{-5}$			
PFIT4LS	3	$1 \times 10^{-7}$	29320	$2.991 \times 10^{-1}$	25011	$1.775 \times 10^{-5}$			
QUARTC	100	$1 \times 10^{-7}$	14367	$2.836 \times 10^{-6}$	29677	$2.399 \times 10^{-9}$			
SINEVAL	2	$1 \times 10^{-7}$	1722	$1.177 \times 10^{-11}$	1879	$1.601 \times 10^{-10}$			
SINQUAD	100	$1 \times 10^{-7}$	5544	$5.534 \times 10^{-9}$	16166	$3.215 \times 10^{-9}$			
SISSER	2	$1 \times 10^{-7}$	657	$2.663 \times 10^{-8}$	703	$6.479 \times 10^{-9}$			
SPARSQUR	100	$1 \times 10^{-7}$	10891	$3.396\times 10^{-7}$	24709	$2.323\times 10^{-9}$			
TOINTGSS	100	$1 \times 10^{-7}$	4087	$4.042\times10^{-9}$	6671	$4.000 \times 10^{-9}$			
TQUARTIC	100	$1 \times 10^{-7}$	9496	$3.400 \times 10^{-7}$	19028	$3.207 \times 10^{-7}$			
TRIDIA	100	$1 \times 10^{-7}$	30055	$2.966 \times 10^{-10}$	114846	$1.242 \times 10^{-19}$			
WATSON	31	$1 \times 10^{-7}$	8296	$1.016 \times 10^{-4}$	21230	$1.002 \times 10^{-4}$			
WOODS	100	$1 \times 10^{-7}$	8594	$5.693 \times 10^{-7}$	20454	$7.795 \times 10^{-8}$			
ZANGWIL2	2	$1 \times 10^{-7}$	657	$-9.543 \times 10^{-11}$	1345	$-1.000 \times 10^{-10}$			

### APPENDIX C

# A Feasible Nonlinear Programming Approach for Constrained Derivative-Free Optimization

				FIBO						knitro					
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.29	1.000	$1.309 \times 10^{-14}$	136	191	0.28	20	104	0.09	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	2	4	0.04	1.000	0.000	10	25	0.03	1	8	0.01	1.000	0.000
HS23	2	5	4	6	0.04	2.000	0.000	12	16	0.03	6	28	0.03	2.000	0.000
HS26	3	1	53	56	12.96	$7.103 \times 10^{-10}$	$2.220 \times 10^{-16}$	5245	39137	12.85	34	244	0.1	$7.434 \times 10^{-12}$	$6.661 \times 10^{-16}$
HS32	3	2	4	7	0.04	1.000	0.000	10	16	0.03	2	15	0.01	1.000	0.000
HS34	3	2	6	9	0.1	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	70	0.08	7	44	0.02	$-8.340 \times 10^{-1}$	0.000
HS40	4	3	22	26	0.36	$-2.500 \times 10^{-1}$	$1.110\times10^{-16}$	56	225	0.26	5	44	0.01	$-2.500 \times 10^{-1}$	$1.110\times10^{-16}$
HS44	4	6	3	8	0.04	-1.000(*)	0.000	12	23	0.03	5	36	0.03	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.67	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.47	72	696	0.41	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.17	1.056	$2.398 \times 10^{-14}$	126	423	1.1	16	121	0.04	$7.669 \times 10^{-17}$	$4.441 \times 10^{-16}$
HS64	3	1	43	46	1.04	$6.300 \times 10^{3}$	0.000	300	662	0.95	12	65	0.04	$6.300 \times 10^{3}$	0.000
HS66	3	2	2	6	0.03	$5.182 \times 10^{-1}(*)$	$3.563 \times 10^{-12}$	7	12	0.02	7	43	0.03	$5.182 \times 10^{-1}$	0.000
HS67	3	14	45	49	1.7	$-1.162 \times 10^{3}(*)$	0.000	472	1408	1.59	14	75	0.03	$-1.162 \times 10^{3}$	0.000
HS72	4	2	7	11	0.26	$7.277 \times 10^{2}$	$2.819\times10^{-18}$	44	99	0.21	8	55	0.03	$7.277 \times 10^{2}$	$4.337 \times 10^{-19}$
HS75	4	5	16	20	0.49	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	180	0.43	6	43	0.04	$5.174 \times 10^{3}$	$8.527 \times 10^{-14}$
HS85	5	21	31	36	1.46	-2.216	0.000	357	903	1.31	11	85	0.06	-2.216	$1.421 \times 10^{-14}$
HS87	6	4	9	15	0.74	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	80	287	0.71	14	223	0.07	$8.997 \times 10^{3}$	$2.416 \times 10^{-13}$
HS88	2	1	5	7	0.4	1.363	0.000	92	318	0.39	27	156	0.08	1.363	$2.893 \times 10^{-17}$
HS89	3	1	10	13	1.42	1.363	0.000	246	889	1.39	5	35	0.04	$7.283 \times 10^{-8}(\dagger)$	$1.332 \times 10^{-1}$
HS90	4	1	12	16	4.22	1.363	0.000	738	3531	4.18	33	285	0.15	1.363	$6.600 \times 10^{-18}$
HS93	6	2	42	48	2.9	$1.351 \times 10^{2}$	$4.800 \times 10^{-12}$	517	2433	2.61	26	295	0.13	$1.351 \times 10^2$	0.000
HS98	6	4	3	10	0.05	3.136(*)	$1.776 \times 10^{-15}$	5	8	0.02	4	40	0.02	3.136	$1.776 \times 10^{-15}$
HS100	7	4	106	113	29.46	$6.806 \times 10^{2}$	0.000	3280	22448	28.72	51	587	0.21	$6.806 \times 10^{2}$	$7.161 \times 10^{-15}$
HS101	7	5	103	110	19.42	$1.810 \times 10^{3}$	$3.955 \times 10^{-16}$	2616	15697	18.39	35	354	0.19	$1.810 \times 10^{3}$	0.000
HS102	7	5	103	110	15.11	$9.119 \times 10^{2}$	$3.816 \times 10^{-16}$	1985	10240	13.86	21	254	0.13	$9.119 \times 10^{2}$	$2.498 \times 10^{-16}$
HS103	7	5	56	63	5.5	$5.437 \times 10^{2}$	$1.318\times10^{-16}$	825	4072	4.59	20	238	0.1	$5.437 \times 10^2$	$2.359 \times 10^{-16}$
HS104	8	5	245	253	43.5	3.951	$9.714 \times 10^{-17}$	6813	37268	41.25	30	359	0.1	3.951	$2.082 \times 10^{-16}$
LOADBAL	31	31	1	15499	0.93	1.547(***)	$1.798 \times 10^{308}$	0	1	0.0	40	1380	0.21	$4.529 \times 10^{-1}$	$2.665 \times 10^{-15}$
OPTPRLOC	30	30	9	39	1.0	$-1.642 \times 10^{1}$	$1.990 \times 10^{-13}$	54	127	0.45	16	583	0.13	$-1.642 \times 10^{1}$	$1.421 \times 10^{-13}$
CB3	3	3	1	4	0.01	2.000	0.000	2	3	0.0	2	15	0.01	2.000	0.000
CRESC50	6	100	1000	1006	215.01	$5.939 \times 10^{-1}(**)$	$5.482\times10^{-13}$	26160	96691	210.07	438	5787	1.76	$5.946 \times 10^{-1}$	$1.188 \times 10^{-14}$
DEMBO7	16	20	19	35	2.37	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	795	2.07	17	354	0.16	$1.749 \times 10^{2}$	$6.661 \times 10^{-16}$
DNIEPER	57	24	24	81	7.72	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$	273	1094	3.25	8	534	0.11	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$
EXPFITA	5	22	1	10999	0.6	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	1	0.0	19	140	0.05	$1.137 \times 10^{-3}$	$1.421 \times 10^{-14}$
HIMMELBI	100	12	114	215	72.23	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12586	67.95	72	7487	0.72	$-1.736 \times 10^{3}$	$1.066 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.21	$1.000 \times 10^{1} (***)$	$1.798\times10^{308}$	0	1	0.0	9	81	0.03	$7.593  imes 10^{-1}$	$1.110 \times 10^{-16}$
TWOBARS	2	2	7	9	0.1	1.509	0.000	24	108	0.09	9	57	0.03	1.509	$4.441 \times 10^{-16}$
DIPIGRI	7	4	116	123	85.85	$6.806 \times 10^2$	$3.581\times10^{-11}$	7461	26264	84.59	48	581	0.18	$6.806  imes 10^2$	0.000

#### C.1. Numerical Results for Feasible Initial Points

Table C.1. Noiseless Problems with Feasible  $x_0$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	2	4	0.2	1.000	$1.310 \times 10^{-14}$	75	104	0.19	8	36	0.02	1.080	0.000
HS22	2	2	1	3	0.03	1.000	$1.789 \times 10^{-12}$	9	23	0.02	0	3	0.0	1.000	$1.843 \times 10^{-14}$
HS23	2	5	3	5	0.05	2.000	0.000	11	14	0.03	3	16	0.01	2.040	0.000
HS26	3	1	4	7	0.09	$1.996 \times 10^{-2}$	$9.592 \times 10^{-14}$	32	66	0.07	14	87	0.04	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.09	1.000	0.000	10	16	0.04	2	15	0.02	1.000	0.000
HS34	3	2	5	8	0.09	$-7.820 \times 10^{-1}$	$1.066 \times 10^{-14}$	31	68	0.08	4	29	0.02	$-8.266 \times 10^{-1}$	0.000
HS40	4	3	1	5	0.02	$-2.454 \times 10^{-1}$	$2.554 \times 10^{-15}$	5	15	0.01	0	4	0.0	$-2.454 \times 10^{-1}$	$7.334 \times 10^{-9}$
HS44	4	6	3	8	0.09	-1.000(*)	0.000	12	23	0.07	4	30	0.02	$-1.500 \times 10^{1}$	0.000
HS47	5	3	6	11	0.14	$5.639 \times 10^{-2}$	$1.466 \times 10^{-12}$	46	118	0.12	17	134	0.04	$4.638 \times 10^{-8}$	$3.167 \times 10^{-9}$
HS50	5	3	19	24	1.26	1.056	$2.398 \times 10^{-14}$	126	423	1.19	9	71	0.03	$4.941 \times 10^{-2}$	$4.441 \times 10^{-16}$
HS64	3	1	1	4	0.05	$6.813 \times 10^{3}$	$4.742 \times 10^{-11}$	14	72	0.05	0	3	0.0	$6.813 \times 10^{3}$	0.000
HS66	3	2	1	4	0.02	$5.800 \times 10^{-1}$	$3.563 \times 10^{-12}$	6	10	0.01	0	3	0.0	$5.800 \times 10^{-1}$	0.000
HS67	3	14	31	34	1.91	$-1.053 \times 10^{3}$	0.000	402	1319	1.82	8	45	0.02	$-1.157 \times 10^{3}$	0.000
HS72	4	2	1	5	0.01	$7.407 \times 10^{2}$	0.000	2	5	0.01	0	4	0.0	$7.407 \times 10^{2}$	0.000
HS75	4	5	1	5	0.02	$5.350 \times 10^{3}$	$1.259 \times 10^{-8}$	4	8	0.01	5	37	0.03	$5.174 \times 10^{3}$	$7.739 \times 10^{-8}$
HS85	5	21	29	34	1.22	-2.061	$1.421 \times 10^{-14}$	348	880	1.04	10	78	0.05	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	1	7	0.07	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	9	24	0.06	0	5	0.0	$8.997 \times 10^{3}$	$2.956 \times 10^{-12}$
HS88	2	1	1	3	0.07	1.385	$2.158 \times 10^{-14}$	19	41	0.07	0	3	0.0	1.385	0.000
HS89	3	1	2	5	0.29	1.365	$1.123 \times 10^{-16}$	57	205	0.28	5	35	0.03	$7.283 \times 10^{-8}(\dagger)$	$1.332 \times 10^{-1}$
HS90	4	1	2	6	0.53	1.446	$6.661 \times 10^{-16}$	86	454	0.52	24	228	0.11	1.362	$5.472 \times 10^{-7}$
HS93	6	2	1	7	0.11	$1.371 \times 10^{2}$	$2.946 \times 10^{-11}$	10	24	0.11	0	5	0.0	$1.371 \times 10^{2}$	0.000
HS98	6	4	2	8	0.03	3.136	0.000	4	6	0.02	3	32	0.02	3.136	0.000
HS100	7	4	1	8	0.12	$7.140 \times 10^{2}$	0.000	22	79	0.12	0	5	0.0	$7.140 \times 10^{2}$	0.000
HS101	7	5	26	33	5.24	$1.982 \times 10^3$	$3.053\times10^{-16}$	740	4383	5.04	9	106	0.04	$1.848 \times 10^3$	0.000
HS102	7	5	34	41	4.55	$9.948 \times 10^2$	$2.776 \times 10^{-17}$	641	3320	4.29	11	159	0.05	$9.191 \times 10^2$	0.000
HS103	7	5	15	22	1.69	$5.623 \times 10^{2}$	$8.913 \times 10^{-12}$	276	1263	1.63	16	202	0.06	$5.437 \times 10^{2}$	$1.040 \times 10^{-7}$
HS104	8	5	1	9	0.11	4.200	$1.459 \times 10^{-11}$	14	41	0.11	0	6	0.0	4.200	0.000
LOADBAL	31	31	1	15499	0.97	1.547(***)	$1.798 \times 10^{308}$	0	1	0.0	6	231	0.04	$5.252 \times 10^{-1}$	$1.710 \times 10^{-14}$
OPTPRLOC	30	30	5	35	0.46	$-1.567 \times 10^{1}$	$2.220 \times 10^{-16}$	44	104	0.4	2	103	0.04	$-1.587 \times 10^{1}$	$6.661 \times 10^{-16}$
CB3	3	3	1	4	0.01	2.000	0.000	2	3	0.01	1	10	0.01	2.000	$7.270 \times 10^{-13}$
CRESC50	6	100	16	22	2.44	$6.057 \times 10^{-1}$	$9.032 \times 10^{-11}$	344	1528	2.38	58	700	0.32	$6.155 \times 10^{-1}$	0.000
DEMBO7	16	20	13	29	1.39	$1.879 \times 10^{2}$	$5.917 \times 10^{-8}$	198	507	1.19	15	318	0.1	$1.749 \times 10^{2}$	$1.869 \times 10^{-13}$
DNIEPER	57	24	4	61	1.82	$1.941 \times 10^{4}$	$4.264 \times 10^{-11}$	96	457	1.35	2	177	0.03	$1.874 \times 10^{4}$	$1.800 \times 10^{-7}$
EXPFITA	5	22	1	10999	0.61	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	1	0.0	7	56	0.02	$9.889 \times 10^{-2}$	$1.354 \times 10^{-14}$
HIMMELBI	100	12	114	215	70.47	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12586	68.05	3	408	0.09	$-1.643 \times 10^{3}$	0.000
SYNTHES1	6	6	1	2999	0.22	$1.000 \times 10^{1}$ (***)	$1.798 \times 10^{308}$	0	1	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	2	4	0.05	1.523	$7.885 \times 10^{-9}$	13	38	0.04	7	49	0.05	1.509	$2.179 \times 10^{-9}$
DIPIGRI	7	4	1	8	0.14	$7.140 \times 10^2$	0.000	22	79	0.13	0	5	0.0	$7.140 \times 10^2$	0.000

Table C.2. Noiseless Problems with Feasible  $x_0$ .  $\tau = 10^{-1}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.
							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	2	4	0.18	1.000	$1.310 \times 10^{-14}$	75	104	0.17	18	88	0.08	1.001	0.000
HS22	2	2	1	3	0.03	1.000	$1.789 \times 10^{-12}$	9	23	0.02	0	3	0.0	1.000	$1.843 \times 10^{-14}$
HS23	2	5	3	5	0.04	2.000	0.000	11	14	0.03	4	20	0.02	2.000	0.000
HS26	3	1	36	39	1.85	$1.737 \times 10^{-4}$	$1.040 \times 10^{-12}$	754	4338	1.77	14	87	0.03	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.05	1.000	0.000	10	16	0.03	2	15	0.02	1.000	0.000
HS34	3	2	6	9	0.11	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	70	0.08	6	39	0.01	$-8.340 \times 10^{-1}$	$5.446 \times 10^{-9}$
HS40	4	3	2	6	0.03	$-2.491 \times 10^{-1}$	$2.189 \times 10^{-9}$	8	21	0.02	3	32	0.02	$-2.500 \times 10^{-1}$	$2.040 \times 10^{-9}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	23	0.03	4	30	0.03	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.41	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.33	72	696	0.35	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.02	1.056	$2.398 \times 10^{-14}$	126	423	0.96	11	85	0.03	$1.339 \times 10^{-6}$	$8.882 \times 10^{-16}$
HS64	3	1	13	16	0.46	$6.303 \times 10^{3}$	$1.443 \times 10^{-15}$	126	399	0.44	7	40	0.02	$6.300 \times 10^{3}$	$7.051 \times 10^{-10}$
HS66	3	2	2	5	0.02	$5.182 \times 10^{-1}$	$3.563 \times 10^{-12}$	7	12	0.02	3	21	0.02	$5.186 \times 10^{-1}$	0.000
HS67	3	14	38	41	1.51	$-1.162 \times 10^{3}$	0.000	448	1377	1.42	9	50	0.02	$-1.161 \times 10^{3}$	0.000
HS72	4	2	5	9	0.17	$7.284 \times 10^{2}$	0.000	38	90	0.16	3	24	0.02	$7.277 \times 10^{2}$	$1.663 \times 10^{-7}$
HS75	4	5	16	20	0.32	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	180	0.28	5	37	0.04	$5.174 \times 10^{3}$	$7.739 \times 10^{-8}$
HS85	5	21	31	36	0.95	-2.216	0.000	357	903	0.87	10	78	0.05	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	1	7	0.07	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	9	24	0.06	0	5	0.0	$8.997 \times 10^3$	$2.956 \times 10^{-12}$
HS88	2	1	2	4	0.2	1.363	$9.216 \times 10^{-16}$	41	139	0.19	23	140	0.06	1.362	$4.623 \times 10^{-7}$
HS89	3	1	4	7	0.49	1.363	$3.663 \times 10^{-16}$	101	374	0.48	5	35	0.04	$7.283 \times 10^{-8}(\dagger)$	$1.332 \times 10^{-1}$
HS90	4	1	6	10	1.54	1.363	$5.260 \times 10^{-17}$	277	1541	1.52	24	228	0.14	1.362	$5.472 \times 10^{-7}$
HS93	6	2	21	27	1.5	$1.352 \times 10^{2}$	$4.695 \times 10^{-11}$	333	1706	1.43	8	87	0.04	$1.351 \times 10^{2}$	0.000
HS98	6	4	2	8	0.04	3.136	0.000	4	6	0.01	3	32	0.02	3.136	0.000
HS100	7	4	51	58	15.18	$6.813 \times 10^2$	$1.388 \times 10^{-17}$	1609	11182	15.01	3	55	0.02	$6.811 \times 10^2$	0.000
HS101	7	5	44	51	7.81	$1.812 \times 10^3$	$3.006\times10^{-10}$	1239	7244	7.65	18	198	0.1	$1.810 \times 10^3$	$2.139 \times 10^{-8}$
HS102	7	5	69	76	10.13	$9.125 \times 10^2$	$4.272\times10^{-13}$	1483	7663	9.89	15	200	0.06	$9.119 \times 10^2$	$5.354 \times 10^{-7}$
HS103	7	5	37	44	3.94	$5.437 \times 10^{2}$	$7.980 \times 10^{-16}$	671	3465	3.8	16	202	0.08	$5.437 \times 10^2$	$1.040 \times 10^{-7}$
HS104	8	5	47	55	6.18	3.952	$5.218 \times 10^{-15}$	1210	6047	5.99	20	259	0.05	3.951	$1.537 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.9	1.547(***)	$1.798 \times 10^{308}$	0	1	0.0	9	330	0.05	$4.531 \times 10^{-1}$	$1.421 \times 10^{-14}$
OPTPRLOC	30	30	6	36	0.44	$-1.642 \times 10^{1}$	0.000	47	111	0.39	3	136	0.06	$-1.641 \times 10^{1}$	0.000
CB3	3	3	1	4	0.01	2.000	0.000	2	3	0.01	1	10	0.0	2.000	$7.270 \times 10^{-13}$
CRESC50	6	100	68	74	18.0	$5.942 \times 10^{-1}$	0.000	2416	11257	17.76	438	5787	1.81	$5.946 \times 10^{-1}$	$1.188 \times 10^{-14}$
DEMBO7	16	20	18	34	2.11	$1.749 \times 10^{2}$	$1.067 \times 10^{-10}$	301	793	2.03	15	318	0.1	$1.749 \times 10^{2}$	$1.869 \times 10^{-13}$
DNIEPER	57	24	5	62	1.42	$1.874 \times 10^{4}$	$1.405 \times 10^{-9}$	100	465	1.35	2	177	0.06	$1.874 \times 10^{4}$	$1.800 \times 10^{-7}$
EXPFITA	5	22	1	10999	0.58	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	1	0.0	12	91	0.03	$1.355 \times 10^{-3}$	$6.750 \times 10^{-14}$
HIMMELBI	100	12	114	215	68.58	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12586	67.1	12	1326	0.26	$-1.734 \times 10^{3}$	$5.684 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.21	$1.000 \times 10^{1}$ (***)	$1.798 \times 10^{308}$	0	1	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	3	5	0.06	1.509	$4.742 \times 10^{-11}$	16	55	0.05	7	49	0.04	1.509	$2.179 \times 10^{-9}$
DIPIGRI	7	4	51	58	11.33	$6.812 \times 10^2$	$6.939\times10^{-18}$	1491	9509	11.16	3	55	0.01	$6.811 \times 10^2$	0.000

Table C.3. Noiseless Problems with Feasible  $x_0$ .  $\tau = 10^{-3}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.29	1.000	$1.309 \times 10^{-14}$	136	191	0.28	19	100	0.1	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	1	3	0.03	1.000	$1.789 \times 10^{-12}$	9	23	0.02	0	3	0.0	1.000	$1.843 \times 10^{-14}$
HS23	2	5	3	5	0.03	2.000	0.000	11	14	0.02	5	24	0.02	2.000	0.000
HS26	3	1	38	41	8.14	$8.872 \times 10^{-7}$	$4.441 \times 10^{-16}$	3356	26976	8.07	14	87	0.03	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.04	1.000	0.000	10	16	0.03	2	15	0.02	1.000	0.000
HS34	3	2	6	9	0.11	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	70	0.08	6	39	0.01	$-8.340 \times 10^{-1}$	$5.446 \times 10^{-9}$
HS40	4	3	4	8	0.06	$-2.500 \times 10^{-1}$	$1.058 \times 10^{-12}$	13	32	0.04	3	32	0.03	$-2.500 \times 10^{-1}$	$2.040 \times 10^{-9}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	23	0.04	4	30	0.02	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.42	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.35	72	696	0.38	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.11	1.056	$2.398 \times 10^{-14}$	126	423	1.06	11	85	0.04	$1.339 \times 10^{-6}$	$8.882 \times 10^{-16}$
HS64	3	1	17	20	0.55	$6.300 \times 10^{3}$	0.000	156	435	0.52	7	40	0.02	$6.300 \times 10^3$	$7.051 \times 10^{-10}$
HS66	3	2	2	5	0.03	$5.182 \times 10^{-1}$	$3.563 \times 10^{-12}$	7	12	0.02	4	27	0.02	$5.182 \times 10^{-1}$	0.000
HS67	3	14	39	42	1.45	$-1.162 \times 10^{3}$	0.000	450	1380	1.38	10	55	0.03	$-1.162 \times 10^{3}$	0.000
HS72	4	2	6	10	0.24	$7.277 \times 10^{2}$	$6.263 \times 10^{-14}$	43	97	0.22	4	31	0.02	$7.277 \times 10^{2}$	$1.572 \times 10^{-10}$
HS75	4	5	16	20	0.31	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	180	0.27	5	37	0.03	$5.174 \times 10^{3}$	$7.739 \times 10^{-8}$
HS85	5	21	31	36	0.92	-2.216	0.000	357	903	0.83	10	78	0.05	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	2	8	0.11	$8.997 \times 10^{3}$	$7.844 \times 10^{-12}$	15	43	0.1	3	33	0.03	$8.997 \times 10^{3}$	$2.396 \times 10^{-8}$
HS88	2	1	4	6	0.31	1.363	$1.301 \times 10^{-14}$	69	218	0.3	23	140	0.08	1.362	$4.623 \times 10^{-7}$
HS89	3	1	5	8	0.55	1.363	0.000	113	432	0.54	5	35	0.04	$7.283 \times 10^{-8}(\dagger)$	$1.332 \times 10^{-1}$
HS90	4	1	11	15	3.6	1.363	$8.118\times 10^{-18}$	718	3450	3.57	24	228	0.1	1.362	$5.472 \times 10^{-7}$
HS93	6	2	25	31	1.57	$1.351 \times 10^2$	$3.331\times10^{-16}$	371	1842	1.5	8	87	0.05	$1.351 \times 10^2$	0.000
HS98	6	4	2	8	0.03	3.136	0.000	4	6	0.02	3	32	0.02	3.136	0.000
HS100	7	4	89	96	24.86	$6.806 \times 10^{2}$	0.000	2860	19399	24.54	14	212	0.07	$6.806 \times 10^{2}$	0.000
HS101	7	5	71	78	12.74	$1.810 \times 10^{3}$	$2.776 \times 10^{-16}$	2004	11919	12.48	18	198	0.08	$1.810 \times 10^{3}$	$2.139 \times 10^{-8}$
HS102	7	5	78	85	10.45	$9.119 \times 10^{2}$	0.000	1601	8225	10.11	15	200	0.07	$9.119 \times 10^{2}$	$5.354 \times 10^{-7}$
HS103	7	5	39	46	4.22	$5.437 \times 10^2$	$6.106\times10^{-16}$	685	3515	4.09	16	202	0.09	$5.437 \times 10^2$	$1.040 \times 10^{-7}$
HS104	8	5	80	88	13.31	3.951	0.000	2262	12381	13.0	20	259	0.04	3.951	$1.537 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.88	1.547(***)	$1.798 \times 10^{308}$	0	1	0.0	17	594	0.09	$4.529 \times 10^{-1}$	$1.483 \times 10^{-14}$
OPTPRLOC	30	30	8	38	0.52	$-1.642 \times 10^{1}$	$4.504 \times 10^{-10}$	53	125	0.44	13	487	0.11	$-1.642 \times 10^{1}$	$4.092 \times 10^{-8}$
CB3	3	3	1	4	0.01	2.000	0.000	2	3	0.01	1	10	0.0	2.000	$7.270 \times 10^{-13}$
CRESC50	6	100	1000	1006	214.64	$5.939 \times 10^{-1}(**)$	$5.482 \times 10^{-13}$	26160	96691	211.17	438	5787	1.92	$5.946 \times 10^{-1}$	$1.188 \times 10^{-14}$
DEMBO7	16	20	19	35	2.06	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	795	1.97	17	354	0.14	$1.749 \times 10^{2}$	$6.661 \times 10^{-16}$
DNIEPER	57	24	5	62	1.4	$1.874 \times 10^{4}$	$1.405 \times 10^{-9}$	100	465	1.36	2	177	0.06	$1.874 \times 10^{4}$	$1.800 \times 10^{-7}$
EXPFITA	5	22	1	10999	0.63	$2.999 \times 10^{1} (***)$	$1.798 \times 10^{308}$	0	1	0.0	13	98	0.03	$1.137 \times 10^{-3}$	$6.047 \times 10^{-12}$
HIMMELBI	100	12	114	215	68.51	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12586	67.29	31	3273	0.41	$-1.736 \times 10^{3}$	$1.155 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.21	$1.000 \times 10^{1} (***)$	$1.798\times10^{308}$	0	1	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	3	5	0.06	1.509	$4.742\times 10^{-11}$	16	55	0.05	7	49	0.03	1.509	$2.179 \times 10^{-9}$
DIPIGRI	7	4	87	94	78.71	$6.806 \times 10^2$	$1.265\times 10^{-13}$	7019	23030	78.4	14	212	0.06	$6.806\times 10^2$	0.000

Table C.4. Noiseless Problems with Feasible  $x_0$ .  $\tau = 10^{-5}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.29	1.000	$1.309 \times 10^{-14}$	136	191	0.28	19	100	0.1	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	1	3	0.03	1.000	$1.789 \times 10^{-12}$	9	23	0.02	0	3	0.0	1.000	$1.843 \times 10^{-14}$
HS23	2	5	3	5	0.04	2.000	0.000	11	14	0.03	5	24	0.03	2.000	0.000
HS26	3	1	43	46	12.4	$2.878 \times 10^{-8}$	$1.110 \times 10^{-16}$	5194	38850	12.32	28	203	0.08	$2.719 \times 10^{-11}$	$8.518 \times 10^{-7}$
HS32	3	2	4	7	0.05	1.000	0.000	10	16	0.03	2	15	0.02	1.000	0.000
HS34	3	2	6	9	0.09	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	70	0.07	6	39	0.02	$-8.340 \times 10^{-1}$	$5.446 \times 10^{-9}$
HS40	4	3	4	8	0.05	$-2.500 \times 10^{-1}$	$1.058 \times 10^{-12}$	13	32	0.04	3	32	0.05	$-2.500 \times 10^{-1}$	$2.040 \times 10^{-9}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	23	0.03	4	30	0.02	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.44	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.36	72	696	0.36	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.12	1.056	$2.398 \times 10^{-14}$	126	423	1.06	13	100	0.03	$2.358 \times 10^{-8}$	$4.441 \times 10^{-16}$
HS64	3	1	17	20	0.55	$6.300 \times 10^3$	0.000	156	435	0.52	7	40	0.02	$6.300 \times 10^3$	$7.051 \times 10^{-10}$
HS66	3	2	2	5	0.03	$5.182 \times 10^{-1}$	$3.563 \times 10^{-12}$	7	12	0.02	5	33	0.02	$5.182 \times 10^{-1}$	$5.154 \times 10^{-11}$
HS67	3	14	40	43	1.46	$-1.162 \times 10^{3}$	$2.274 \times 10^{-12}$	458	1389	1.36	11	60	0.03	$-1.162 \times 10^{3}$	$1.364 \times 10^{-12}$
HS72	4	2	7	11	0.21	$7.277 \times 10^{2}$	$2.819\times10^{-18}$	44	99	0.2	4	31	0.02	$7.277 \times 10^{2}$	$1.572 \times 10^{-10}$
HS75	4	5	16	20	0.31	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	180	0.28	5	37	0.04	$5.174 \times 10^{3}$	$7.739 \times 10^{-8}$
HS85	5	21	31	36	1.41	-2.216	0.000	357	903	1.3	10	78	0.03	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	9	15	0.7	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	80	287	0.67	6	73	0.03	$8.997 \times 10^{3}$	$1.152 \times 10^{-8}$
HS88	2	1	4	6	0.31	1.363	$1.301 \times 10^{-14}$	69	218	0.3	23	140	0.07	1.362	$4.623 \times 10^{-7}$
HS89	3	1	6	9	0.68	1.363	$5.535 \times 10^{-16}$	142	532	0.66	5	35	0.04	$7.283 \times 10^{-8}(\dagger)$	$1.332 \times 10^{-1}$
HS90	4	1	11	15	3.6	1.363	$8.118\times10^{-18}$	718	3450	3.57	24	228	0.1	1.362	$5.472 \times 10^{-7}$
HS93	6	2	36	42	2.18	$1.351 \times 10^{2}$	0.000	483	2291	2.06	9	96	0.05	$1.351 \times 10^{2}$	0.000
HS98	6	4	2	8	0.02	3.136	0.000	4	6	0.02	3	32	0.01	3.136	0.000
HS100	7	4	106	113	28.46	$6.806 \times 10^{2}$	0.000	3280	22448	28.08	28	362	0.19	$6.806 \times 10^{2}$	$5.093 \times 10^{-11}$
HS101	7	5	90	97	15.48	$1.810 \times 10^{3}$	$2.742 \times 10^{-12}$	2377	14118	15.15	20	216	0.1	$1.810 \times 10^{3}$	$1.032 \times 10^{-8}$
HS102	7	5	89	96	12.09	$9.119 \times 10^{2}$	$2.012 \times 10^{-16}$	1847	9679	11.74	15	200	0.08	$9.119 \times 10^{2}$	$5.354 \times 10^{-7}$
HS103	7	5	42	49	4.08	$5.437 \times 10^{2}$	$4.302 \times 10^{-16}$	713	3618	3.92	16	202	0.08	$5.437 \times 10^2$	$1.040 \times 10^{-7}$
HS104	8	5	88	96	13.71	3.951	$4.441 \times 10^{-16}$	2400	13125	13.38	20	259	0.05	3.951	$1.537 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.89	1.547(***)	$1.798 \times 10^{308}$	0	1	0.0	18	627	0.09	$4.529 \times 10^{-1}$	$8.549 \times 10^{-15}$
OPTPRLOC	30	30	8	38	0.51	$-1.642 \times 10^{1}$	$4.504 \times 10^{-10}$	53	125	0.43	13	487	0.11	$-1.642 \times 10^{1}$	$4.092 \times 10^{-8}$
CB3	3	3	1	4	0.01	2.000	0.000	2	3	0.0	1	10	0.0	2.000	$7.270 \times 10^{-13}$
CRESC50	6	100	1000	1006	214.17	$5.939 \times 10^{-1}(**)$	$5.482 \times 10^{-13}$	26160	96691	210.8	438	5787	1.87	$5.946 \times 10^{-1}$	$1.188 \times 10^{-14}$
DEMBO7	16	20	19	35	2.11	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	795	2.04	17	354	0.15	$1.749 \times 10^{2}$	$6.661 \times 10^{-16}$
DNIEPER	57	24	23	80	3.43	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$	272	1088	3.27	6	416	0.08	$1.874 \times 10^{4}$	$5.623 \times 10^{-7}$
EXPFITA	5	22	1	10999	0.6	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	1	0.0	14	105	0.04	$1.137 \times 10^{-3}$	$1.776 \times 10^{-14}$
HIMMELBI	100	12	114	215	69.11	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12586	68.01	41	4300	0.52	$-1.736 \times 10^{3}$	$7.105 \times 10^{-15}$
SYNTHES1	6	6	1	2999	0.23	$1.000 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	1	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	5	7	0.1	1.509	0.000	21	84	0.08	7	49	0.03	1.509	$2.179 \times 10^{-9}$
DIPIGRI	7	4	96	103	80.88	$6.806 \times 10^2$	$4.547\times10^{-13}$	7235	24605	80.56	28	365	0.1	$6.806\times 10^2$	$7.237\times10^{-11}$

Table C.5. Noiseless Problems with Feasible  $x_0$ .  $\tau = 10^{-7}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

### C.2. Numerical Results for Infeasible Initial Point

We now include the cost of obtaining an initial feasible point for FIBO. In other words, the CPU time and number of constraint evaluations associated with obtaining a feasible starting point are counted. Simple comparison with the previous table indicates that only a few constraint evaluations are needed in order to obtain a feasible solution. As for KNITRO, we instead run from the given initial point  $x_0$ , which is potentially infeasible.



Figure C.1. Infeasible  $x_0$ . Log-ratio Plot for Comparing the Final Accuracy (4.3.1) of FIBO and FD.

As indicated by Figures C.1-C.3, the conclusions remain the same as in Section 4.3 and the cost of obtaining a feasible starting point is quite negligible.



Figure C.2. Infeasible  $x_0$ . Log-ratio plot comparing FIBO and FD in terms of the number of function evaluations to satisfy (4.3.2) for  $\tau = 10^{-1}$  (upper left),  $10^{-3}$  (upper right),  $10^{-5}$  (bottom left),  $10^{-7}$  (bottom right).



Figure C.3. Infeasible  $x_0$ . Log-ratio plot comparing FIBO and FD in terms of the number of constraint evaluations to satisfy (4.3.2) for  $\tau = 10^{-1}$  (upper left),  $10^{-3}$  (upper right),  $10^{-5}$  (bottom left),  $10^{-7}$  (bottom right).

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.32	1.000	$1.309 \times 10^{-14}$	136	193	0.28	20	106	0.09	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	2	4	0.05	1.000	0.000	10	27	0.03	5	24	0.01	1.000	0.000
HS23	2	5	4	6	0.05	2.000	0.000	12	18	0.03	7	32	0.01	2.000	0.000
HS26	3	1	53	56	13.0	$7.103 \times 10^{-10}$	$2.220 \times 10^{-16}$	5245	39138	12.85	34	244	0.11	$7.434 \times 10^{-12}$	$6.661 \times 10^{-16}$
HS32	3	2	4	7	0.05	1.000	0.000	10	17	0.03	2	15	0.01	1.000	0.000
HS34	3	2	6	9	0.12	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	71	0.08	7	44	0.02	$-8.340\times10^{-1}$	0.000
HS40	4	3	22	26	0.37	$-2.500 \times 10^{-1}$	$1.110 \times 10^{-16}$	56	229	0.26	7	51	0.02	$-2.500 \times 10^{-1}$	$1.665 \times 10^{-16}$
HS44	4	6	3	8	0.06	-1.000(*)	0.000	12	24	0.03	5	36	0.01	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.67	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7879	3.47	72	696	0.35	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.18	1.056	$2.398 \times 10^{-14}$	126	424	1.1	16	121	0.04	$7.669 \times 10^{-17}$	$4.441 \times 10^{-16}$
HS64	3	1	43	46	1.07	$6.300 \times 10^{3}$	0.000	300	674	0.95	16	89	0.04	$6.300 \times 10^{3}$	$1.110 \times 10^{-16}$
HS66	3	2	2	6	0.04	$5.182 \times 10^{-1}(*)$	$3.563 \times 10^{-12}$	7	13	0.02	7	43	0.01	$5.182 \times 10^{-1}$	0.000
HS67	3	14	45	49	1.71	$-1.162 \times 10^{3}(*)$	0.000	472	1409	1.59	14	75	0.03	$-1.162 \times 10^{3}$	0.000
HS72	4	2	7	11	0.32	$7.277 \times 10^{2}$	$2.819\times10^{-18}$	44	111	0.21	15	96	0.05	$7.277 \times 10^2$	$1.301 \times 10^{-18}$
HS75	4	5	16	20	0.54	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	185	0.43	7	48	0.03	$5.174 \times 10^3$	$2.842 \times 10^{-14}$
HS85	5	21	31	36	1.47	-2.216	0.000	357	904	1.31	11	85	0.03	-2.216	$1.421 \times 10^{-14}$
HS87	6	4	9	15	0.76	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	80	290	0.71	14	216	0.07	$8.997 \times 10^3$	$2.274 \times 10^{-13}$
HS88	2	1	5	7	0.43	1.363	0.000	92	329	0.39	19	106	0.06	1.363	0.000
HS89	3	1	10	13	1.52	1.363	0.000	246	1033	1.39	28	172	0.09	1.363	$9.636 \times 10^{-18}$
HS90	4	1	12	16	4.26	1.363	0.000	738	3541	4.18	48	426	0.24	1.363	0.000
HS93	6	2	42	48	2.9	$1.351 \times 10^{2}$	$4.800 \times 10^{-12}$	517	2434	2.61	26	295	0.08	$1.351 \times 10^2$	0.000
HS98	6	4	3	10	0.09	3.136(*)	$1.776 \times 10^{-15}$	5	13	0.02	6	56	0.01	3.136	0.000
HS100	7	4	106	113	29.47	$6.806 \times 10^{2}$	0.000	3280	22449	28.72	51	587	0.16	$6.806 \times 10^2$	$7.161 \times 10^{-15}$
HS101	7	5	103	110	19.45	$1.810 \times 10^{3}$	$3.955 \times 10^{-16}$	2616	15705	18.39	56	582	0.22	$1.810 \times 10^3$	0.000
HS102	7	5	103	110	15.14	$9.119 \times 10^{2}$	$3.816\times10^{-16}$	1985	10248	13.86	37	386	0.13	$9.119 \times 10^2$	$2.498 \times 10^{-16}$
HS103	7	5	56	63	5.53	$5.437 \times 10^2$	$1.318\times10^{-16}$	825	4079	4.59	28	284	0.08	$5.437 \times 10^2$	$2.220 \times 10^{-16}$
HS104	8	5	245	253	43.52	3.951	$9.714 \times 10^{-17}$	6813	37274	41.25	17	190	0.04	3.951	$1.943 \times 10^{-16}$
LOADBAL	31	31	1	15499	0.94	1.547(***)	$1.798 \times 10^{308}$	0	2	0.0	40	1380	0.19	$4.529 \times 10^{-1}$	$2.665 \times 10^{-15}$
OPTPRLOC	30	30	9	39	1.01	$-1.642 \times 10^{1}$	$1.990 \times 10^{-13}$	54	130	0.45	12	456	0.06	$-1.642 \times 10^{1}$	$1.279 \times 10^{-13}$
CB3	3	3	1	4	0.03	2.000	0.000	2	9	0.0	7	40	0.01	2.000	0.000
CRESC50	6	100	1000	1006	215.05	$5.939 \times 10^{-1}(**)$	$5.482 \times 10^{-13}$	26160	96697	210.07	456	6023	2.46	$5.948 \times 10^{-1}$	$4.547 \times 10^{-13}$
DEMBO7	16	20	19	35	2.42	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	800	2.07	45	893	0.2	$1.748 \times 10^{2}$	$3.608 \times 10^{-16}$
DNIEPER	57	24	24	81	7.75	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$	273	1097	3.25	6	418	0.07	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$
EXPFITA	5	22	1	10999	0.62	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	19	140	0.05	$1.137 \times 10^{-3}$	$1.421 \times 10^{-14}$
HIMMELBI	100	12	114	215	72.25	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12588	67.95	96	9950	1.01	$-1.736 \times 10^{3}$	$2.842 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.24	$1.000 \times 10^{1}$ (***)	$1.798\times10^{308}$	0	2	0.0	9	81	0.02	$7.593 \times 10^{-1}$	$1.110\times 10^{-16}$
TWOBARS	2	2	7	9	0.13	1.509	0.000	24	113	0.09	11	57	0.02	1.509	$2.220\times 10^{-16}$
DIPIGRI	7	4	116	123	85.87	$6.806  imes 10^2$	$3.581\times 10^{-11}$	7461	26265	84.59	48	581	0.16	$6.806  imes 10^2$	0.000

Table C.6. Noiseless Problems with Infeasible  $x_0$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	2	4	0.2	1.000	$1.310 \times 10^{-14}$	75	106	0.19	8	38	0.02	1.080	0.000
HS22	2	2	1	3	0.04	1.000	$1.789 \times 10^{-12}$	9	25	0.02	3	16	0.01	1.000	$1.231 \times 10^{-6}$
HS23	2	5	3	5	0.07	2.000	0.000	11	16	0.03	3	16	0.01	2.075	0.000
HS26	3	1	4	7	0.09	$1.996 \times 10^{-2}$	$9.592 \times 10^{-14}$	32	67	0.07	14	87	0.03	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.1	1.000	0.000	10	17	0.04	2	15	0.01	1.000	0.000
HS34	3	2	5	8	0.11	$-7.820 \times 10^{-1}$	$1.066 \times 10^{-14}$	31	69	0.08	4	29	0.01	$-8.266 \times 10^{-1}$	0.000
HS40	4	3	1	5	0.03	$-2.454 \times 10^{-1}$	$2.554 \times 10^{-15}$	5	19	0.01	4	33	0.01	$-2.500 \times 10^{-1}$	$5.162 \times 10^{-8}$
HS44	4	6	3	8	0.1	-1.000(*)	0.000	12	24	0.07	4	30	0.01	$-1.500 \times 10^{1}$	0.000
HS47	5	3	6	11	0.15	$5.639 \times 10^{-2}$	$1.466 \times 10^{-12}$	46	119	0.12	17	134	0.04	$4.638 \times 10^{-8}$	$3.167 \times 10^{-9}$
HS50	5	3	19	24	1.26	1.056	$2.398 \times 10^{-14}$	126	423	1.19	9	71	0.03	$4.941 \times 10^{-2}$	$4.441 \times 10^{-16}$
HS64	3	1	1	4	0.09	$6.813 \times 10^{3}$	$4.742 \times 10^{-11}$	14	84	0.05	5	30	0.02	$6.705 \times 10^{3}$	$3.942 \times 10^{-5}$
HS66	3	2	1	4	0.05	$5.800 \times 10^{-1}$	$3.563 \times 10^{-12}$	6	11	0.01	0	3	0.0	$5.800 \times 10^{-1}$	0.000
HS67	3	14	31	34	1.92	$-1.053 \times 10^3$	0.000	402	1320	1.82	8	45	0.02	$-1.157 \times 10^{3}$	0.000
HS72	4	2	1	5	0.03	$7.407 \times 10^{2}$	0.000	2	17	0.01	11	72	0.04	$7.277 \times 10^{2}$	$8.766 \times 10^{-7}$
HS75	4	5	1	5	0.06	$5.350 \times 10^{3}$	$1.259 \times 10^{-8}$	4	13	0.01	5	36	0.03	$5.174 \times 10^{3}$	$1.880 \times 10^{-4}$
HS85	5	21	29	34	1.23	-2.061	$1.421 \times 10^{-14}$	348	881	1.04	10	78	0.03	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	1	7	0.1	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	9	27	0.06	3	33	0.02	$8.997 \times 10^{3}$	$3.775 \times 10^{-8}$
HS88	2	1	1	3	0.13	1.385	$2.158 \times 10^{-14}$	19	52	0.07	16	94	0.05	1.363	$1.074 \times 10^{-8}$
HS89	3	1	2	5	0.41	1.365	$1.123 \times 10^{-16}$	57	349	0.28	23	147	0.07	1.362	$7.400 \times 10^{-7}$
HS90	4	1	2	6	0.56	1.446	$6.661 \times 10^{-16}$	86	464	0.52	36	334	0.15	1.363	$1.142 \times 10^{-7}$
HS93	6	2	1	7	0.13	$1.371 \times 10^{2}$	$2.946 \times 10^{-11}$	10	25	0.11	0	5	0.0	$1.371 \times 10^{2}$	0.000
HS98	6	4	2	8	0.06	3.136	0.000	4	11	0.02	4	40	0.01	3.214	0.000
HS100	7	4	1	8	0.14	$7.140 \times 10^{2}$	0.000	22	80	0.12	0	5	0.0	$7.140 \times 10^{2}$	0.000
HS101	7	5	26	33	5.28	$1.982 \times 10^{3}$	$3.053 \times 10^{-16}$	740	4391	5.04	26	269	0.1	$1.810 \times 10^{3}$	$2.249 \times 10^{-5}$
HS102	7	5	34	41	4.59	$9.948 \times 10^{2}$	$2.776 \times 10^{-17}$	641	3328	4.29	25	272	0.09	$9.129 \times 10^{2}$	$5.866 \times 10^{-5}$
HS103	7	5	15	22	1.72	$5.623 \times 10^2$	$8.913\times10^{-12}$	276	1270	1.63	20	212	0.07	$5.434 \times 10^2$	$3.104 \times 10^{-4}$
HS104	8	5	1	9	0.15	4.200	$1.459 \times 10^{-11}$	14	47	0.11	11	130	0.03	3.951	$5.769 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.99	1.547(***)	$1.798 \times 10^{308}$	0	2	0.0	6	231	0.04	$5.252 \times 10^{-1}$	$1.710 \times 10^{-14}$
OPTPRLOC	30	30	5	35	0.48	$-1.567 \times 10^{1}$	$2.220 \times 10^{-16}$	44	107	0.4	12	456	0.06	$-1.642 \times 10^{1}$	$1.279 \times 10^{-13}$
CB3	3	3	1	4	0.02	2.000	0.000	2	9	0.01	5	30	0.01	2.000	$3.017 \times 10^{-6}$
CRESC50	6	100	16	22	2.49	$6.057 \times 10^{-1}$	$9.032 \times 10^{-11}$	344	1534	2.38	10	91	0.06	$6.567 \times 10^{-1}$	$9.586 \times 10^{-4}$
DEMBO7	16	20	13	29	1.44	$1.879 \times 10^{2}$	$5.917 \times 10^{-8}$	198	512	1.19	16	358	0.12	$1.748 \times 10^{2}$	$1.074 \times 10^{-5}$
DNIEPER	57	24	4	61	1.86	$1.941 \times 10^{4}$	$4.264 \times 10^{-11}$	96	460	1.35	5	359	0.07	$1.874 \times 10^4$	$3.740 \times 10^{-10}$
EXPFITA	5	22	1	10999	0.62	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	7	56	0.02	$9.889 \times 10^{-2}$	$1.354 \times 10^{-14}$
HIMMELBI	100	12	114	215	70.51	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12588	68.05	3	408	0.05	$-1.625 \times 10^{3}$	0.000
SYNTHES1	6	6	1	2999	0.23	$1.000 \times 10^{1} (***)$	$1.798 \times 10^{308}$	0	2	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	2	4	0.07	1.523	$7.885 \times 10^{-9}$	13	43	0.04	4	26	0.01	1.558	0.000
DIPIGRI	7	4	1	8	0.15	$7.140 \times 10^2$	0.000	22	80	0.13	0	5	0.0	$7.140 \times 10^2$	0.000

Table C.7. Noiseless Problems with Infeasible  $x_0$ .  $\tau = 10^{-1}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	2	4	0.19	1.000	$1.310 \times 10^{-14}$	75	106	0.17	18	90	0.08	1.001	0.000
HS22	2	2	1	3	0.04	1.000	$1.789 \times 10^{-12}$	9	25	0.02	3	16	0.01	1.000	$1.231 \times 10^{-6}$
HS23	2	5	3	5	0.05	2.000	0.000	11	16	0.03	4	20	0.01	2.001	0.000
HS26	3	1	36	39	1.86	$1.737 \times 10^{-4}$	$1.040 \times 10^{-12}$	754	4339	1.77	14	87	0.03	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.06	1.000	0.000	10	17	0.03	2	15	0.01	1.000	0.000
HS34	3	2	6	9	0.11	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	71	0.08	6	39	0.01	$-8.340 \times 10^{-1}$	$5.446 \times 10^{-9}$
HS40	4	3	2	6	0.05	$-2.491 \times 10^{-1}$	$2.189 \times 10^{-9}$	8	25	0.02	4	33	0.01	$-2.500 \times 10^{-1}$	$5.162 \times 10^{-8}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	24	0.03	4	30	0.01	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.41	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.33	72	696	0.36	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.02	1.056	$2.398 \times 10^{-14}$	126	423	0.96	11	85	0.03	$1.339 \times 10^{-6}$	$8.882 \times 10^{-16}$
HS64	3	1	13	16	0.5	$6.303 \times 10^{3}$	$1.443 \times 10^{-15}$	126	411	0.44	9	54	0.02	$6.306 \times 10^{3}$	0.000
HS66	3	2	2	5	0.03	$5.182 \times 10^{-1}$	$3.563 \times 10^{-12}$	7	13	0.02	3	21	0.01	$5.186 \times 10^{-1}$	0.000
HS67	3	14	38	41	1.51	$-1.162 \times 10^{3}$	0.000	448	1378	1.42	9	50	0.02	$-1.161 \times 10^{3}$	0.000
HS72	4	2	5	9	0.2	$7.284 \times 10^{2}$	0.000	38	102	0.16	11	72	0.05	$7.277 \times 10^{2}$	$8.766 \times 10^{-7}$
HS75	4	5	16	20	0.35	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	185	0.28	5	36	0.03	$5.174 \times 10^{3}$	$1.880 \times 10^{-4}$
HS85	5	21	31	36	0.96	-2.216	0.000	357	904	0.87	10	78	0.03	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	1	7	0.09	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	9	27	0.06	3	33	0.01	$8.997 \times 10^{3}$	$3.775 \times 10^{-8}$
HS88	2	1	2	4	0.23	1.363	$9.216 \times 10^{-16}$	41	150	0.19	16	94	0.05	1.363	$1.074 \times 10^{-8}$
HS89	3	1	4	7	0.6	1.363	$3.663 \times 10^{-16}$	101	518	0.48	23	147	0.07	1.362	$7.400 \times 10^{-7}$
HS90	4	1	6	10	1.57	1.363	$5.260 \times 10^{-17}$	277	1551	1.52	36	334	0.16	1.363	$1.142 \times 10^{-7}$
HS93	6	2	21	27	1.51	$1.352 \times 10^{2}$	$4.695 \times 10^{-11}$	333	1707	1.43	8	87	0.03	$1.351 \times 10^{2}$	0.000
HS98	6	4	2	8	0.06	3.136	0.000	4	11	0.01	6	56	0.01	3.136	0.000
HS100	7	4	51	58	15.19	$6.813 \times 10^{2}$	$1.388 \times 10^{-17}$	1609	11183	15.01	3	55	0.02	$6.811 \times 10^2$	0.000
HS101	7	5	44	51	7.84	$1.812 \times 10^{3}$	$3.006 \times 10^{-10}$	1239	7252	7.65	26	269	0.1	$1.810 \times 10^3$	$2.249 \times 10^{-5}$
HS102	7	5	69	76	10.15	$9.125 \times 10^{2}$	$4.272 \times 10^{-13}$	1483	7671	9.89	26	282	0.09	$9.120 \times 10^{2}$	$3.705 \times 10^{-6}$
HS103	7	5	37	44	3.97	$5.437 \times 10^2$	$7.980 \times 10^{-16}$	671	3472	3.8	20	212	0.07	$5.434 \times 10^2$	$3.104 \times 10^{-4}$
HS104	8	5	47	55	6.2	3.952	$5.218 \times 10^{-15}$	1210	6053	5.99	11	130	0.03	3.951	$5.769 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.92	1.547(***)	$1.798 \times 10^{308}$	0	2	0.0	9	330	0.05	$4.531 \times 10^{-1}$	$1.421 \times 10^{-14}$
OPTPRLOC	30	30	6	36	0.45	$-1.642 \times 10^{1}$	0.000	47	114	0.39	12	456	0.06	$-1.642 \times 10^{1}$	$1.279 \times 10^{-13}$
CB3	3	3	1	4	0.03	2.000	0.000	2	9	0.01	5	30	0.01	2.000	$3.017 \times 10^{-6}$
CRESC50	6	100	68	74	18.03	$5.942 \times 10^{-1}$	0.000	2416	11263	17.76	456	6023	2.47	$5.948 \times 10^{-1}$	$4.547 \times 10^{-13}$
DEMBO7	16	20	18	34	2.17	$1.749 \times 10^{2}$	$1.067 \times 10^{-10}$	301	798	2.03	16	358	0.11	$1.748 \times 10^{2}$	$1.074 \times 10^{-5}$
DNIEPER	57	24	5	62	1.45	$1.874 \times 10^{4}$	$1.405 \times 10^{-9}$	100	468	1.35	5	359	0.07	$1.874 \times 10^{4}$	$3.740 \times 10^{-10}$
EXPFITA	5	22	1	10999	0.59	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	12	91	0.03	$1.355 \times 10^{-3}$	$6.750 \times 10^{-14}$
HIMMELBI	100	12	114	215	68.61	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12588	67.1	16	1734	0.19	$-1.734 \times 10^{3}$	$3.553 \times 10^{-15}$
SYNTHES1	6	6	1	2999	0.22	$1.000 \times 10^{1}$ (***)	$1.798 \times 10^{308}$	0	2	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	3	5	0.08	1.509	$4.742\times10^{-11}$	16	60	0.05	9	49	0.02	1.509	$8.844\times 10^{-8}$
DIPIGRI	7	4	51	58	11.34	$6.812 \times 10^2$	$6.939\times10^{-18}$	1491	9510	11.16	3	55	0.01	$6.811 \times 10^2$	0.000

Table C.8. Noiseless Problems with Infeasible  $x_0$ .  $\tau = 10^{-3}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.29	1.000	$1.309 \times 10^{-14}$	136	191	0.28	19	102	0.09	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	1	3	0.04	1.000	$1.789 \times 10^{-12}$	9	25	0.02	3	16	0.01	1.000	$1.231 \times 10^{-6}$
HS23	2	5	3	5	0.04	2.000	0.000	11	16	0.02	5	24	0.01	2.000	0.000
HS26	3	1	38	41	8.15	$8.872 \times 10^{-7}$	$4.441 \times 10^{-16}$	3356	26977	8.07	14	87	0.03	$4.252 \times 10^{-7}$	$9.633 \times 10^{-7}$
HS32	3	2	4	7	0.05	1.000	0.000	10	17	0.03	2	15	0.01	1.000	0.000
HS34	3	2	6	9	0.11	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	71	0.08	6	39	0.01	$-8.340\times10^{-1}$	$5.446  imes 10^{-9}$
HS40	4	3	4	8	0.08	$-2.500 \times 10^{-1}$	$1.058 \times 10^{-12}$	13	36	0.04	4	33	0.01	$-2.500 \times 10^{-1}$	$5.162 \times 10^{-8}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	24	0.04	4	30	0.01	$-1.500 \times 10^{1}$	0.000
HS47	5	3	23	28	3.42	$1.989 \times 10^{-7}$	$3.781 \times 10^{-11}$	1391	7878	3.35	72	696	0.36	$1.007 \times 10^{1}$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.11	1.056	$2.398 \times 10^{-14}$	126	423	1.06	11	85	0.03	$1.339 \times 10^{-6}$	$8.882 \times 10^{-16}$
HS64	3	1	17	20	0.58	$6.300 \times 10^{3}$	0.000	156	447	0.52	11	64	0.03	$6.300 \times 10^{3}$	$6.916 \times 10^{-6}$
HS66	3	2	2	5	0.03	$5.182 \times 10^{-1}$	$3.563\times10^{-12}$	7	13	0.02	4	27	0.01	$5.182 \times 10^{-1}$	0.000
HS67	3	14	39	42	1.46	$-1.162 \times 10^{3}$	0.000	450	1381	1.38	10	55	0.03	$-1.162 \times 10^{3}$	0.000
HS72	4	2	6	10	0.27	$7.277 \times 10^{2}$	$6.263\times10^{-14}$	43	109	0.22	11	72	0.04	$7.277 \times 10^{2}$	$8.766 \times 10^{-7}$
HS75	4	5	16	20	0.34	$5.174 \times 10^{3}$	$5.684 \times 10^{-14}$	97	185	0.27	5	36	0.03	$5.174 \times 10^{3}$	$1.880 \times 10^{-4}$
HS85	5	21	31	36	0.94	-2.216	0.000	357	904	0.83	10	78	0.03	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	2	8	0.13	$8.997 \times 10^{3}$	$7.844 \times 10^{-12}$	15	46	0.1	3	33	0.02	$8.997 \times 10^3$	$3.775 \times 10^{-8}$
HS88	2	1	4	6	0.34	1.363	$1.301\times10^{-14}$	69	229	0.3	16	94	0.05	1.363	$1.074 \times 10^{-8}$
HS89	3	1	5	8	0.64	1.363	0.000	113	576	0.54	23	147	0.08	1.362	$7.400 \times 10^{-7}$
HS90	4	1	11	15	3.63	1.363	$8.118 \times 10^{-18}$	718	3460	3.57	36	334	0.18	1.363	$1.142 \times 10^{-7}$
HS93	6	2	25	31	1.58	$1.351 \times 10^{2}$	$3.331 \times 10^{-16}$	371	1843	1.5	8	87	0.03	$1.351 \times 10^{2}$	0.000
HS98	6	4	2	8	0.05	3.136	0.000	4	11	0.02	6	56	0.01	3.136	0.000
HS100	7	4	89	96	24.87	$6.806 \times 10^{2}$	0.000	2860	19400	24.54	14	212	0.07	$6.806 \times 10^{2}$	0.000
HS101	7	5	71	78	12.77	$1.810 \times 10^{3}$	$2.776 \times 10^{-16}$	2004	11927	12.48	27	280	0.1	$1.810 \times 10^{3}$	$2.398 \times 10^{-5}$
HS102	7	5	78	85	10.48	$9.119 \times 10^{2}$	0.000	1601	8233	10.11	28	305	0.09	$9.119 \times 10^{2}$	$2.176 \times 10^{-6}$
HS103	7	5	39	46	4.26	$5.437 \times 10^{2}$	$6.106 \times 10^{-16}$	685	3522	4.09	20	212	0.07	$5.434 \times 10^{2}$	$3.104 \times 10^{-4}$
HS104	8	5	80	88	13.33	3.951	0.000	2262	12387	13.0	11	130	0.03	3.951	$5.769 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.89	1.547(***)	$1.798 \times 10^{308}$	0	2	0.0	17	594	0.09	$4.529 \times 10^{-1}$	$1.483 \times 10^{-14}$
OPTPRLOC	30	30	8	38	0.54	$-1.642 \times 10^{1}$	$4.504 \times 10^{-10}$	53	128	0.44	12	456	0.06	$-1.642 \times 10^{1}$	$1.279 \times 10^{-13}$
CB3	3	3	1	4	0.03	2.000	0.000	2	9	0.01	5	30	0.01	2.000	$3.017 \times 10^{-6}$
CRESC50	6	100	1000	1006	214.67	$5.939 \times 10^{-1}(**)$	$5.482 \times 10^{-13}$	26160	96697	211.17	456	6023	2.37	$5.948 \times 10^{-1}$	$4.547 \times 10^{-13}$
DEMB07	16	20	19	35	2.11	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	800	1.97	16	358	0.1	$1.748 \times 10^{2}$	$1.074 \times 10^{-5}$
DNIEPER	57	24	5	62	1.43	$1.874 \times 10^{4}$	$1.405 \times 10^{-9}$	100	468	1.36	5	359	0.07	$1.874 \times 10^{4}$	$3.740 \times 10^{-10}$
EXPEITA	5	22	1	10999	0.63	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	13	98	0.03	$1.137 \times 10^{-3}$	$6.047 \times 10^{-12}$
HIMMELBI	100	12	114	215	68.54	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12588	67.29	41	4291	0.45	$-1.736 \times 10^{3}$	$1.421 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.21	$1.000 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	7	65	0.01	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	3	5	0.08	1.509	$4.742 \times 10^{-11}$	16	60	0.05	9	49	0.02	1.509	$8.844 \times 10^{-8}$
DIPIGRI	7	4	87	94	78.71	$6.806 \times 10^{2}$	$1.265 \times 10^{-13}$	7019	23031	78.4	14	212	0.06	$6.806 \times 10^{2}$	0.000

Table C.9. Noiseless Problems with Infeasible  $x_0 \tau = 10^{-5}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

							FIBO							knitro	
Problem	n	m	#iter	#feval	time	f	feas err	#iter(sub)	#ceval	time(sub)	#iter	#feval	time	f	feas err
HS13	2	1	3	5	0.29	1.000	$1.309 \times 10^{-14}$	136	191	0.28	19	102	0.09	$9.983 \times 10^{-1}$	$6.579 \times 10^{-10}$
HS22	2	2	1	3	0.04	1.000	$1.789 \times 10^{-12}$	9	25	0.02	3	16	0.01	1.000	$1.231 \times 10^{-6}$
HS23	2	5	3	5	0.05	2.000	0.000	11	16	0.03	6	28	0.01	2.000	0.000
HS26	3	1	43	46	12.41	$2.878 \times 10^{-8}$	$1.110 \times 10^{-16}$	5194	38851	12.32	28	203	0.08	$2.719 \times 10^{-11}$	$8.518 \times 10^{-7}$
HS32	3	2	4	7	0.05	1.000	0.000	10	17	0.03	2	15	0.01	1.000	0.000
HS34	3	2	6	9	0.09	$-8.340 \times 10^{-1}$	$1.776 \times 10^{-15}$	32	71	0.07	6	39	0.01	$-8.340 \times 10^{-1}$	$5.446 \times 10^{-9}$
HS40	4	3	4	8	0.07	$-2.500 \times 10^{-1}$	$1.058 \times 10^{-12}$	13	36	0.04	4	33	0.01	$-2.500 \times 10^{-1}$	$5.162 \times 10^{-8}$
HS44	4	6	3	8	0.05	-1.000(*)	0.000	12	24	0.03	4	30	0.01	$-1.500 imes10^1$	0.000
HS47	5	3	23	28	3.44	$1.989 \times 10^{-7}$	$3.781\times10^{-11}$	1391	7878	3.36	72	696	0.33	$1.007 \times 10^1$	$6.661 \times 10^{-16}$
HS50	5	3	19	24	1.12	1.056	$2.398\times10^{-14}$	126	423	1.06	13	100	0.03	$2.358 \times 10^{-8}$	$4.441 \times 10^{-16}$
HS64	3	1	17	20	0.59	$6.300 \times 10^{3}$	0.000	156	447	0.52	11	64	0.03	$6.300 \times 10^{3}$	$6.916 \times 10^{-6}$
HS66	3	2	2	5	0.03	$5.182 \times 10^{-1}$	$3.563 \times 10^{-12}$	7	13	0.02	5	33	0.01	$5.182 \times 10^{-1}$	$5.154 \times 10^{-11}$
HS67	3	14	40	43	1.46	$-1.162 \times 10^{3}$	$2.274 \times 10^{-12}$	458	1390	1.36	11	60	0.03	$-1.162 \times 10^{3}$	$1.364 \times 10^{-12}$
HS72	4	2	7	11	0.25	$7.277 \times 10^{2}$	$2.819\times10^{-18}$	44	111	0.2	11	72	0.04	$7.277 \times 10^{2}$	$8.766 \times 10^{-7}$
HS75	4	5	16	20	0.35	$5.174 \times 10^{3}$	$5.684\times10^{-14}$	97	185	0.28	5	36	0.03	$5.174 \times 10^{3}$	$1.880 \times 10^{-4}$
HS85	5	21	31	36	1.42	-2.216	0.000	357	904	1.3	10	78	0.03	-2.216	$2.389 \times 10^{-8}$
HS87	6	4	9	15	0.7	$8.997 \times 10^{3}$	$2.274 \times 10^{-13}$	80	287	0.67	6	75	0.02	$8.997 \times 10^3$	$7.329 \times 10^{-11}$
HS88	2	1	4	6	0.34	1.363	$1.301\times10^{-14}$	69	229	0.3	16	94	0.05	1.363	$1.074 \times 10^{-8}$
HS89	3	1	6	9	0.78	1.363	$5.535 \times 10^{-16}$	142	676	0.66	23	147	0.07	1.362	$7.400 \times 10^{-7}$
HS90	4	1	11	15	3.63	1.363	$8.118\times 10^{-18}$	718	3460	3.57	36	334	0.18	1.363	$1.142 \times 10^{-7}$
HS93	6	2	36	42	2.19	$1.351 \times 10^{2}$	0.000	483	2292	2.06	9	96	0.03	$1.351 \times 10^{2}$	0.000
HS98	6	4	2	8	0.05	3.136	0.000	4	11	0.02	6	56	0.01	3.136	0.000
HS100	7	4	106	113	28.46	$6.806 \times 10^{2}$	0.000	3280	22448	28.08	28	362	0.11	$6.806 \times 10^{2}$	$5.093 \times 10^{-11}$
HS101	7	5	90	97	15.51	$1.810 \times 10^{3}$	$2.742 \times 10^{-12}$	2377	14126	15.15	28	289	0.11	$1.810 \times 10^3$	$1.302 \times 10^{-5}$
HS102	7	5	89	96	12.12	$9.119 \times 10^{2}$	$2.012 \times 10^{-16}$	1847	9687	11.74	29	314	0.1	$9.119 \times 10^{2}$	$7.385 \times 10^{-7}$
HS103	7	5	42	49	4.11	$5.437 \times 10^{2}$	$4.302 \times 10^{-16}$	713	3625	3.92	20	212	0.07	$5.434 \times 10^{2}$	$3.104 \times 10^{-4}$
HS104	8	5	88	96	13.73	3.951	$4.441 \times 10^{-16}$	2400	13131	13.38	11	130	0.03	3.951	$5.769 \times 10^{-7}$
LOADBAL	31	31	1	15499	0.9	1.547(***)	$1.798 \times 10^{308}$	0	2	0.0	18	627	0.09	$4.529 \times 10^{-1}$	$8.549 \times 10^{-15}$
OPTPRLOC	30	30	8	38	0.53	$-1.642 \times 10^{1}$	$4.504 \times 10^{-10}$	53	128	0.43	12	456	0.06	$-1.642 \times 10^{1}$	$1.279 \times 10^{-13}$
CB3	3	3	1	4	0.02	2.000	0.000	2	9	0.0	5	30	0.01	2.000	$3.017 \times 10^{-6}$
CRESC50	6	100	1000	1006	214.2	$5.939 \times 10^{-1}(**)$	$5.482 \times 10^{-13}$	26160	96697	210.8	456	6023	2.34	$5.948 \times 10^{-1}$	$4.547 \times 10^{-13}$
DEMBO7	16	20	19	35	2.11	$1.748 \times 10^{2}$	$4.073 \times 10^{-15}$	302	795	2.04	16	358	0.11	$1.748 \times 10^{2}$	$1.074 \times 10^{-5}$
DNIEPER	57	24	23	80	3.47	$1.874 \times 10^{4}$	$5.684 \times 10^{-14}$	272	1091	3.27	5	359	0.06	$1.874 \times 10^{4}$	$3.740 \times 10^{-10}$
EXPFITA	5	22	1	10999	0.6	$2.999 \times 10^{1}(***)$	$1.798 \times 10^{308}$	0	2	0.0	14	105	0.04	$1.137 \times 10^{-3}$	$1.776 \times 10^{-14}$
HIMMELBI	100	12	114	215	69.13	$-1.467 \times 10^{3}(*)$	$1.789 \times 10^{-9}$	2618	12588	68.01	46	4805	0.48	$-1.736 \times 10^{3}$	$2.132 \times 10^{-14}$
SYNTHES1	6	6	1	2999	0.24	$1.000 \times 10^{1}$ (***)	$1.798 \times 10^{308}$	0	2	0.0	7	65	0.02	$7.593 \times 10^{-1}$	$3.261 \times 10^{-7}$
TWOBARS	2	2	5	7	0.11	1.509	0.000	21	89	0.08	9	49	0.02	1.509	$8.844 \times 10^{-8}$
DIPIGRI	7	4	96	103	80.89	$6.806 \times 10^2$	$4.547\times 10^{-13}$	7235	24606	80.56	28	365	0.11	$6.806  imes 10^2$	$7.237\times10^{-11}$

Table C.10. Noiseless Problems with Infeasible  $x_0 \tau = 10^{-7}$ ; n: number of variables, m: number of constraints, #iter: number of (outer) iterations, #feval: number of function evaluations, time: total CPU time passed, f: final objective value, feas err: final feasibility error, #iter(sub): total number of iterations for solving TR subproblem, #ceval: total number of constraint evaluations(note that the number of constraint evaluations and function evaluations are same for KNITRO), time(sub): total CPU time elapsed for solving subproblem. \* indicates that FIBO terminates with singular interpolation system error, \*\* indicates that FIBO terminates with maximum number of iterations, \*\*\* indicates that FIBO terminates with maximum number of function evaluations.

### APPENDIX D

# Analyzing the Performance of DFO Methods on a Wider Class of Problems

## D.1. Proof for Theorem 6.3.1

**Proof.** By definition of the quadratic model (4.2.1), we have that

$$f(y^{i}) - m(x) = m(y^{i}) - m(x)$$
  
=  $(y^{i} - x)^{T}(g + Hx) + \frac{1}{2}(y^{i} - x)^{T}H(y^{i} - x)$   
=  $(y^{i} - x)^{T}\nabla m(x) + \frac{1}{2}(y^{i} - x)^{T}H(y^{i} - x)$ 

Therefore, using Taylor expansion, it holds that

(D.1.1)  

$$(y^{i} - x)^{T} (\nabla m(x) - \nabla \phi(x))$$
(D.1.2)  

$$= \phi(y^{i}) - (y^{i} - x)^{T} \nabla \phi(x) - \phi(x) - \frac{1}{2} (y^{i} - x)^{T} H(y^{i} - x) + \phi(x) - m(x) + \epsilon(y^{i})$$
(D.1.3)  

$$= \frac{1}{2} (y^{i} - x)^{T} \nabla^{2} f(x + t_{i}(y^{i} - x))(y^{i} - x) - \frac{1}{2} (y^{i} - x)^{T} H(y^{i} - x) + \phi(x) - m(x) + \epsilon(y^{i})$$

for some  $t_i \in (0, 1)$ 

Subtracting the (D.1.3) associated with  $y^0$  from i = 1, ..., p,

$$\begin{aligned} &(y^{i} - y^{0})(\nabla m(x) - \nabla \phi(x)) \\ &= \frac{1}{2}(y^{i} - x)^{T} \nabla^{2} f(x + t_{i}(y^{i} - x))(y^{i} - x) - \frac{1}{2}(y^{i} - x)^{T} H(y^{i} - x) \\ &- \frac{1}{2}(y^{0} - x)^{T} \nabla^{2} f(x + t_{i}(y^{0} - x))(y^{0} - x) + \frac{1}{2}(y^{0} - x)^{T} H(y^{0} - x) + \epsilon(y^{i}) - \epsilon(y^{0}) \\ &\leq \frac{5}{2}(\|H\| + L_{2})\Delta^{2} + 2\epsilon_{f} \end{aligned}$$

where the last inequality is derived using  $||y^i - x|| \leq 2\Delta$  and  $||y^0 - x|| \leq \Delta$  and the bounded noise assumption.

Combining p equations, we have that

$$\|M(\nabla m(x) - \nabla \phi(x))\| \le \frac{5\sqrt{p}}{2}(\|H\| + L_2)\Delta^2 + 2\sqrt{p}\epsilon_f$$

where

$$M = \begin{bmatrix} y^1 - y^0 & y^2 - y^0 & \dots & y^p - y^0 \end{bmatrix}$$

Defining the scaled matrix  $\hat{M} = M/\Delta$ , we have that

$$\|\nabla m(x) - \nabla \phi(x)\| \le \frac{5\sqrt{p} \|\hat{M}^{\dagger}\|}{2} (\|H\| + L_2)\Delta + \frac{2\sqrt{p} \|\hat{M}^{\dagger}\|\epsilon_f}{\Delta}$$

In addition, substituting  $x = y^0$  in (D.1.3), we have that

$$(y^{i} - y^{0})^{T} (\nabla m(x) - \nabla \phi(x)) \leq \frac{1}{2} (L_{2} + ||H||) \Delta^{2} + \phi(y^{0}) - m(y^{0}) + \epsilon(y^{i})$$
  

$$\implies m(y^{0}) - \phi(y^{0}) \leq \frac{1}{2} (L_{2} + ||H||) \Delta^{2} + ||\nabla m(x) - \nabla \phi(x)||\Delta + \epsilon_{f}$$
  

$$\leq \frac{5\sqrt{p} ||\hat{M}^{\dagger}|| + 1}{2} (||H|| + L_{2}) \Delta^{2} + (2\sqrt{p} ||\hat{M}^{\dagger}|| + 1) \epsilon_{f}$$

#### D.2. Numerical Results for Comparing DFO-TR against CMA-ES

An alternative IBO solver is DFO-TR, which has higher linear algebra cost than NEWUOA at every iteration. Therefore, we consider problems with less than 50 variables and validate our findings with respect to NEWUOA. We set final trust region radius to  $10^{-8}$  and the stop\_predict convergence test to 0 to avoid early termination. As demonstrated in Figure D.1 - D.2, similar conclusions can be made regarding the relative performance of DFOTR and CMA-ES. However, we note in Figure D.3 that relaxing the ratio test improves the final accuracy of DFO-TR, especially when the noise level is large.



Figure D.1. (Selected) Accuracy Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.1) comparing CMA-ES and DFO-TR with restarts within a budget of 500n.



Figure D.2. (Selected) Efficiency Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.3) comparing CMA-ES and DFO-TR with restarts for  $\tau = 10^{-8}$ .



Figure D.3. (Selected) Accuracy Log-Ratio Profiles for Noisy Problems with noise level 0.001 (Left) and  $10^{-7}$  (Right). Plots of (6.2.1) comparing DFO-TR with and without ratio relaxation using a budget of 500n.